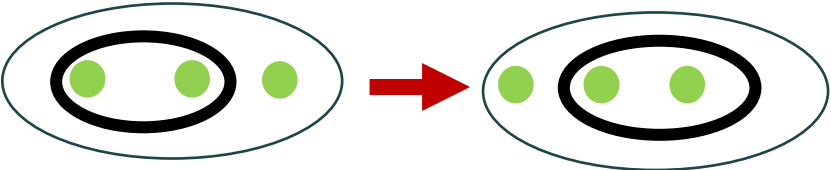
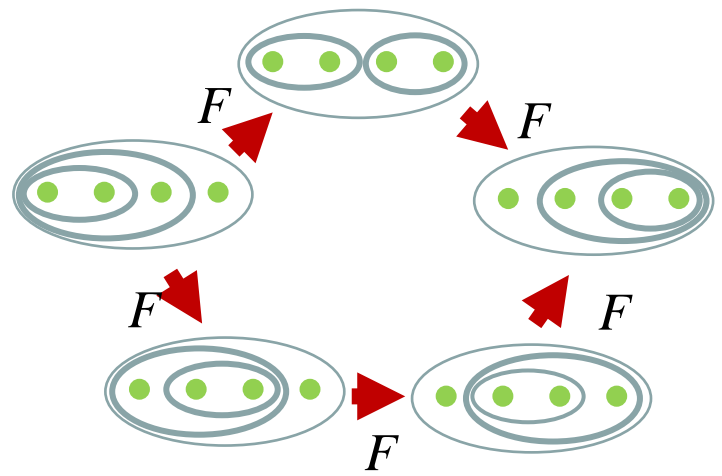


● ← Fibonacci anyon with topological “charge” 1

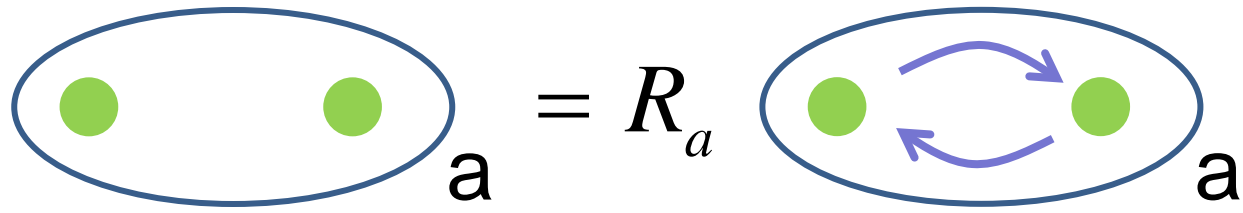
Fusion rules: $1 \times 1 = 0 + 1$; $1 \times 0 = 0 \times 1 = 1$; $0 \times 0 = 0$

F Matrix:  $F = \begin{pmatrix} \varphi^{-1} & \varphi^{-1/2} \\ \varphi^{-1/2} & -\varphi^{-1} \end{pmatrix}$

Pentagon

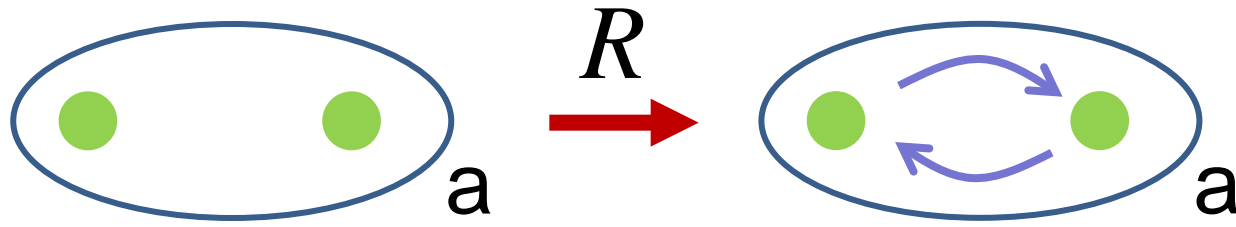


The R Matrix



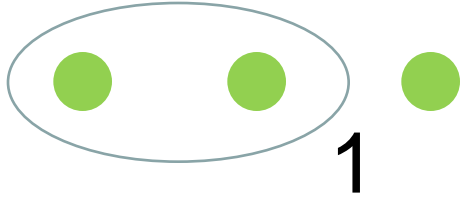
$$R = \begin{pmatrix} R_0 & 0 \\ 0 & R_1 \end{pmatrix}$$

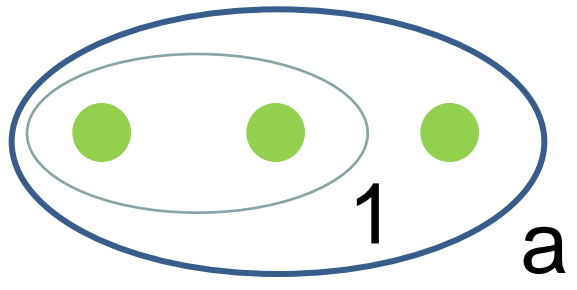
The R Matrix

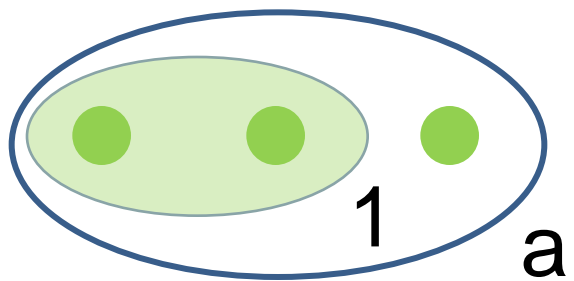


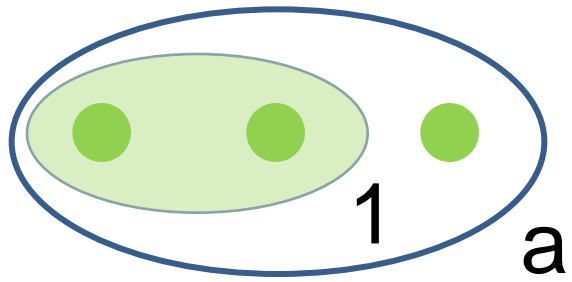
$$R = \begin{pmatrix} R_0 & 0 \\ 0 & R_1 \end{pmatrix}$$



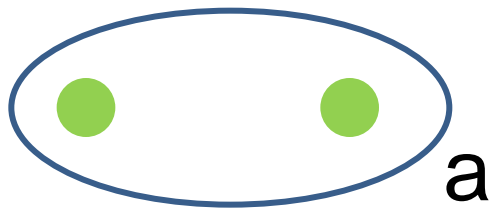


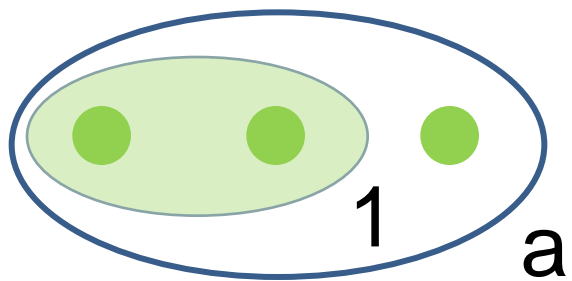




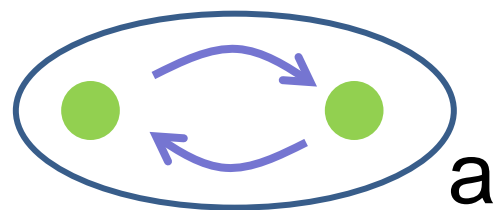
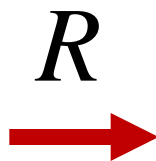
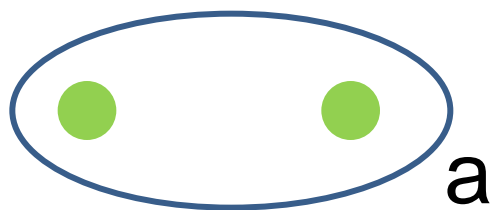


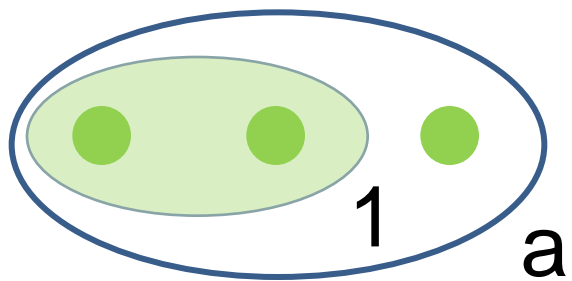
Equivalent to



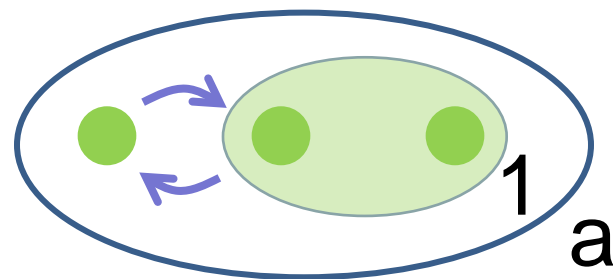


Equivalent to

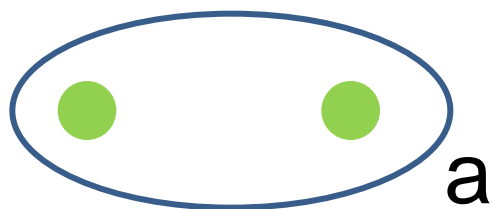




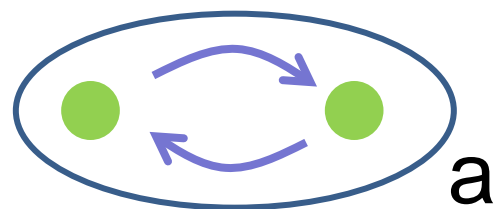
R
→

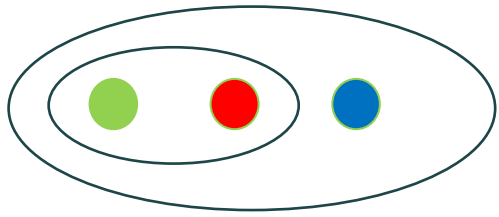


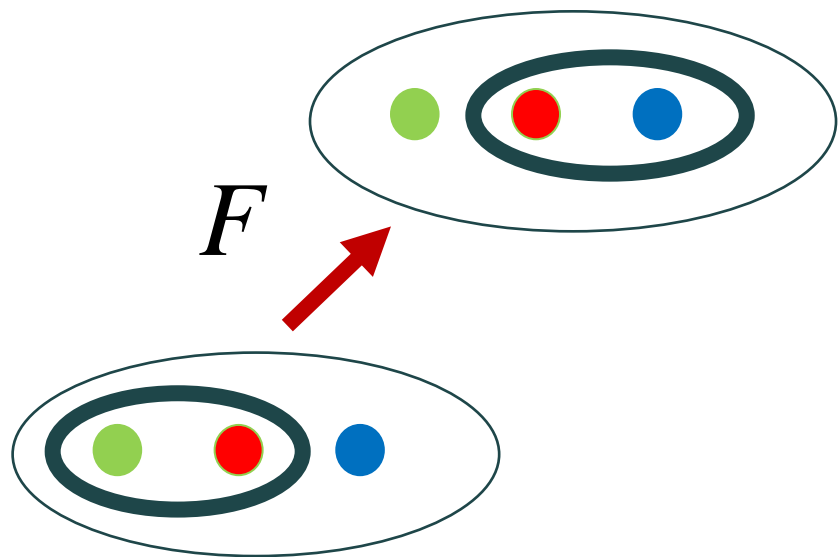
Equivalent to

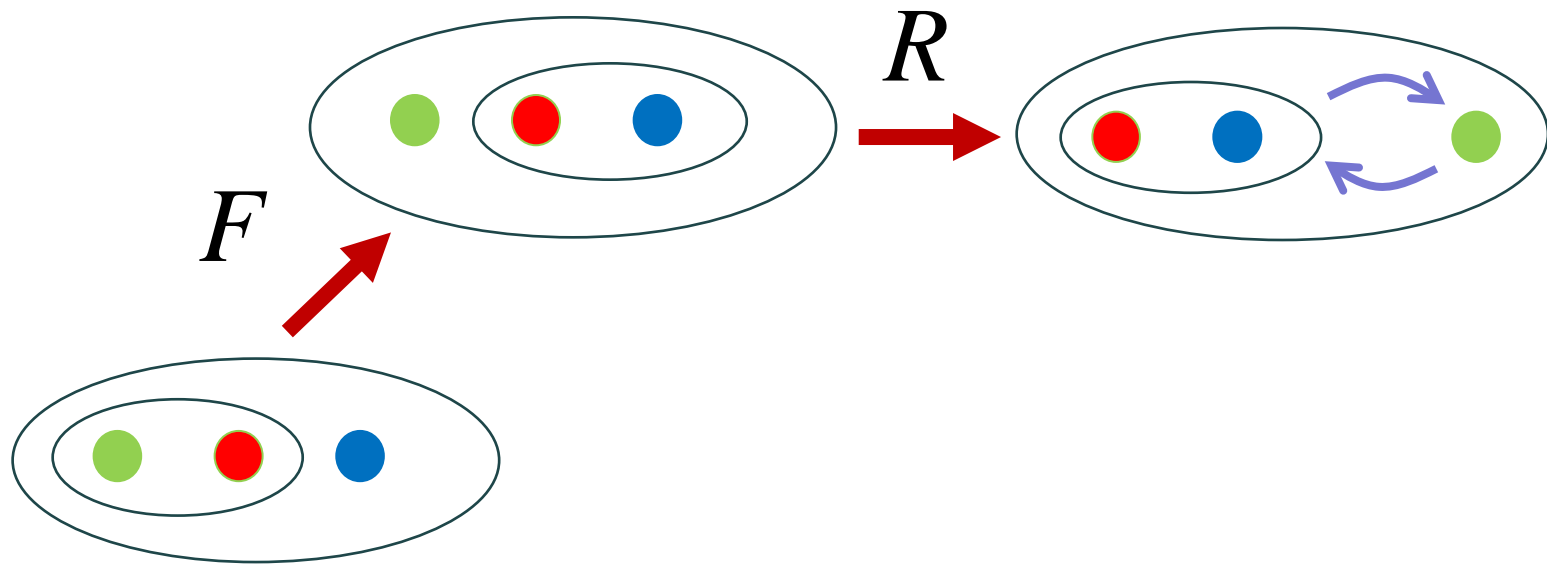


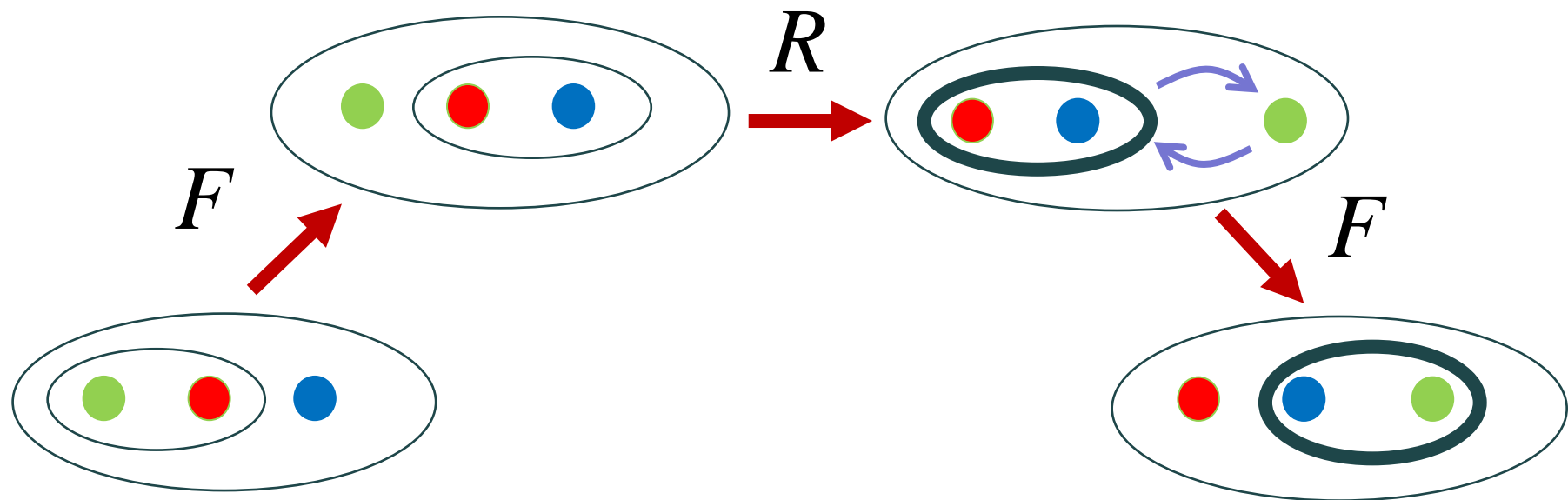
R
→

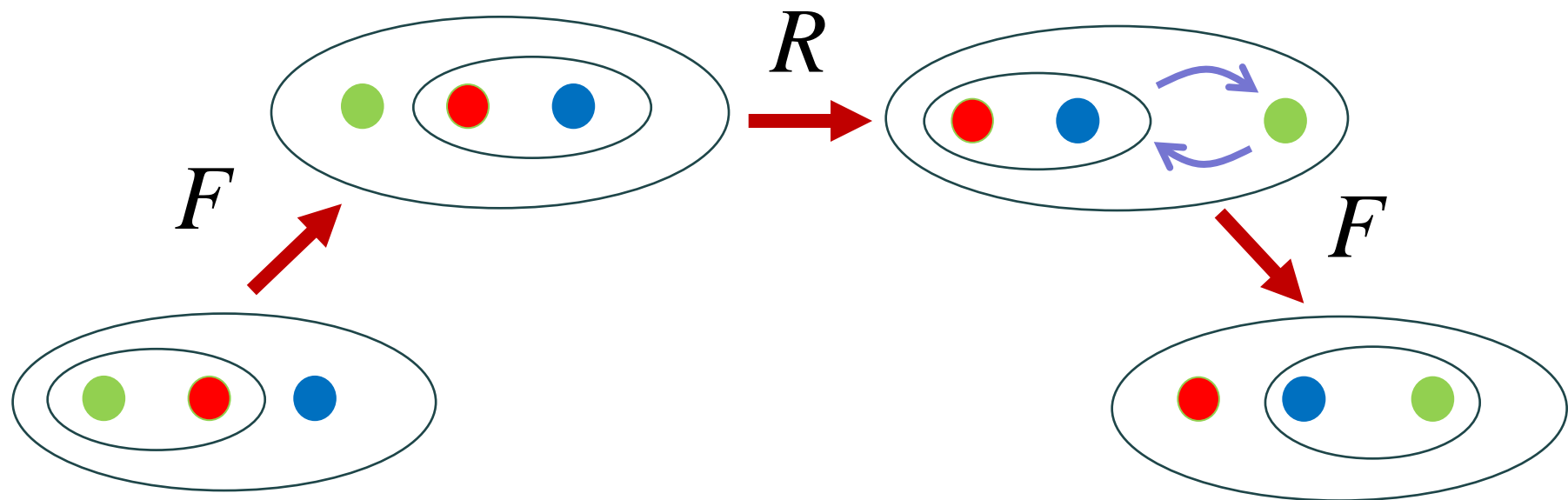


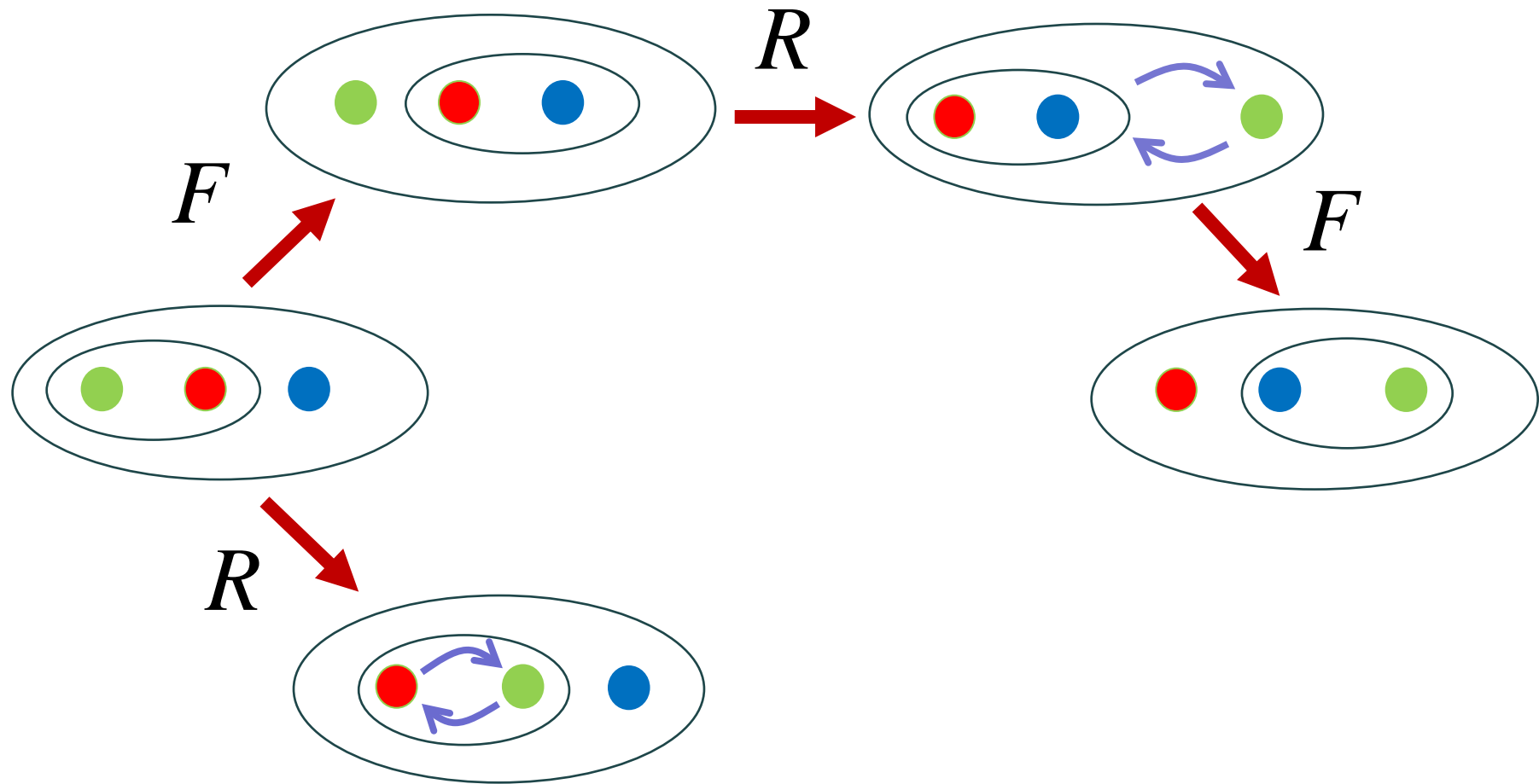


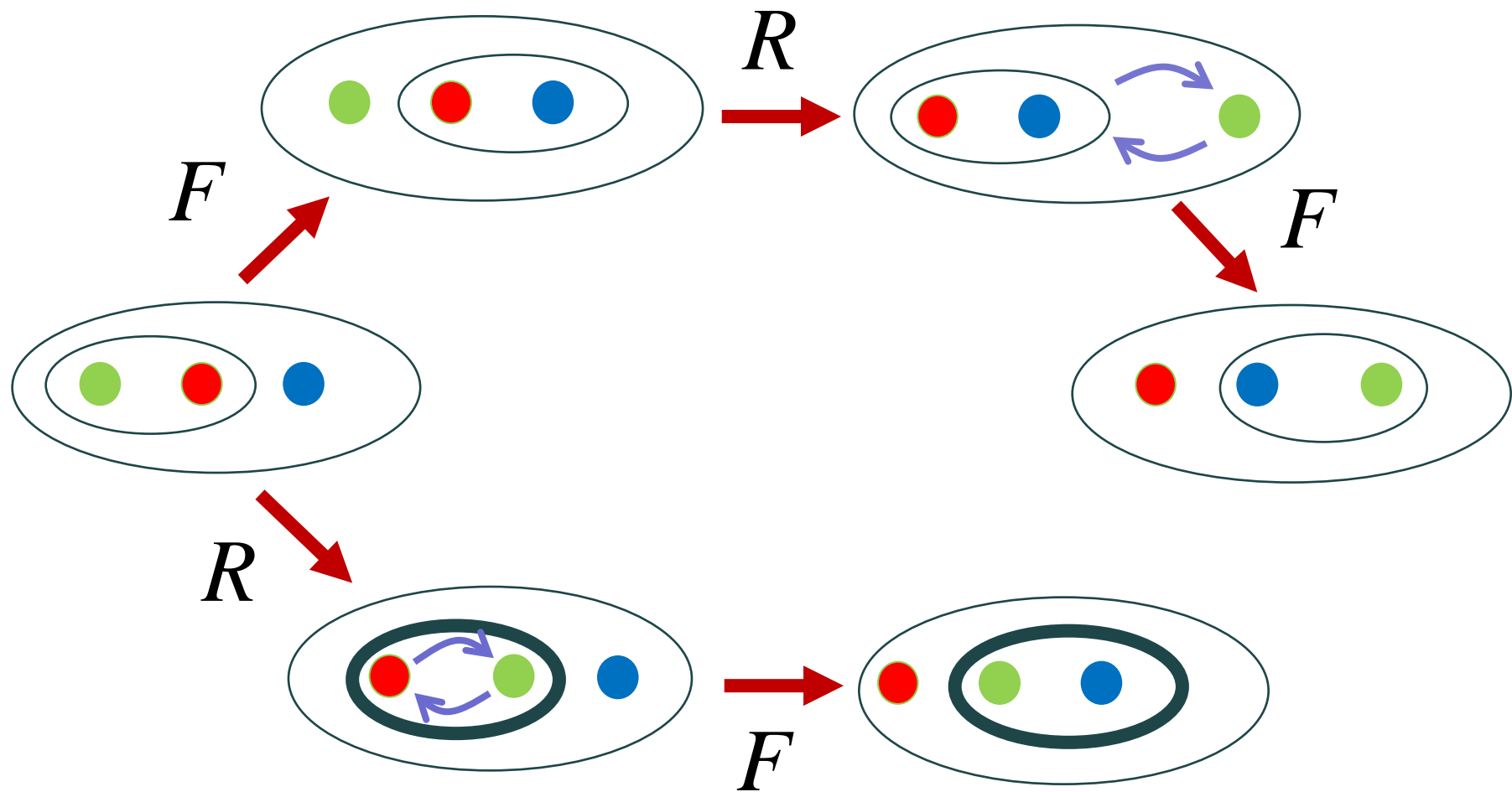


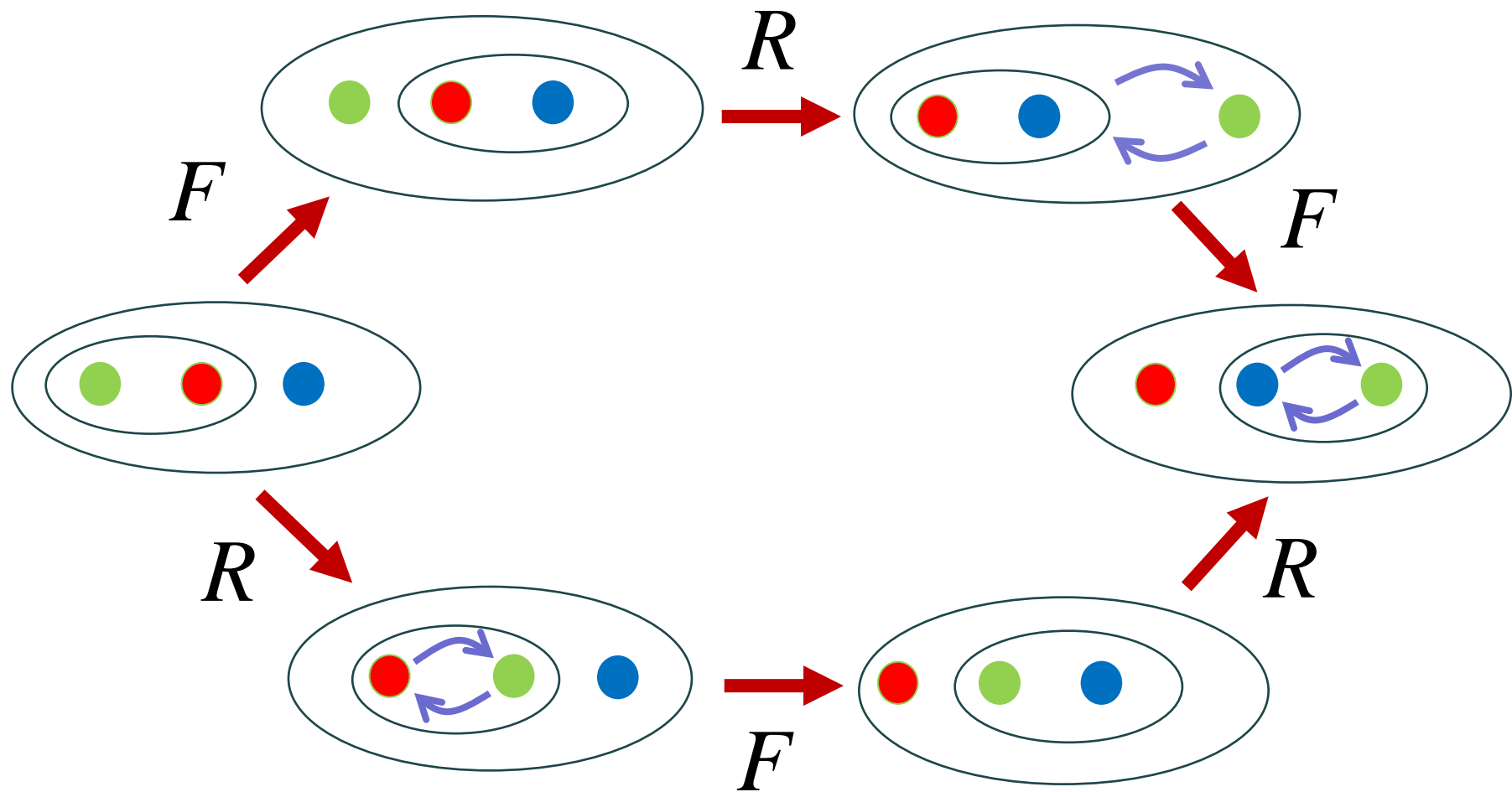




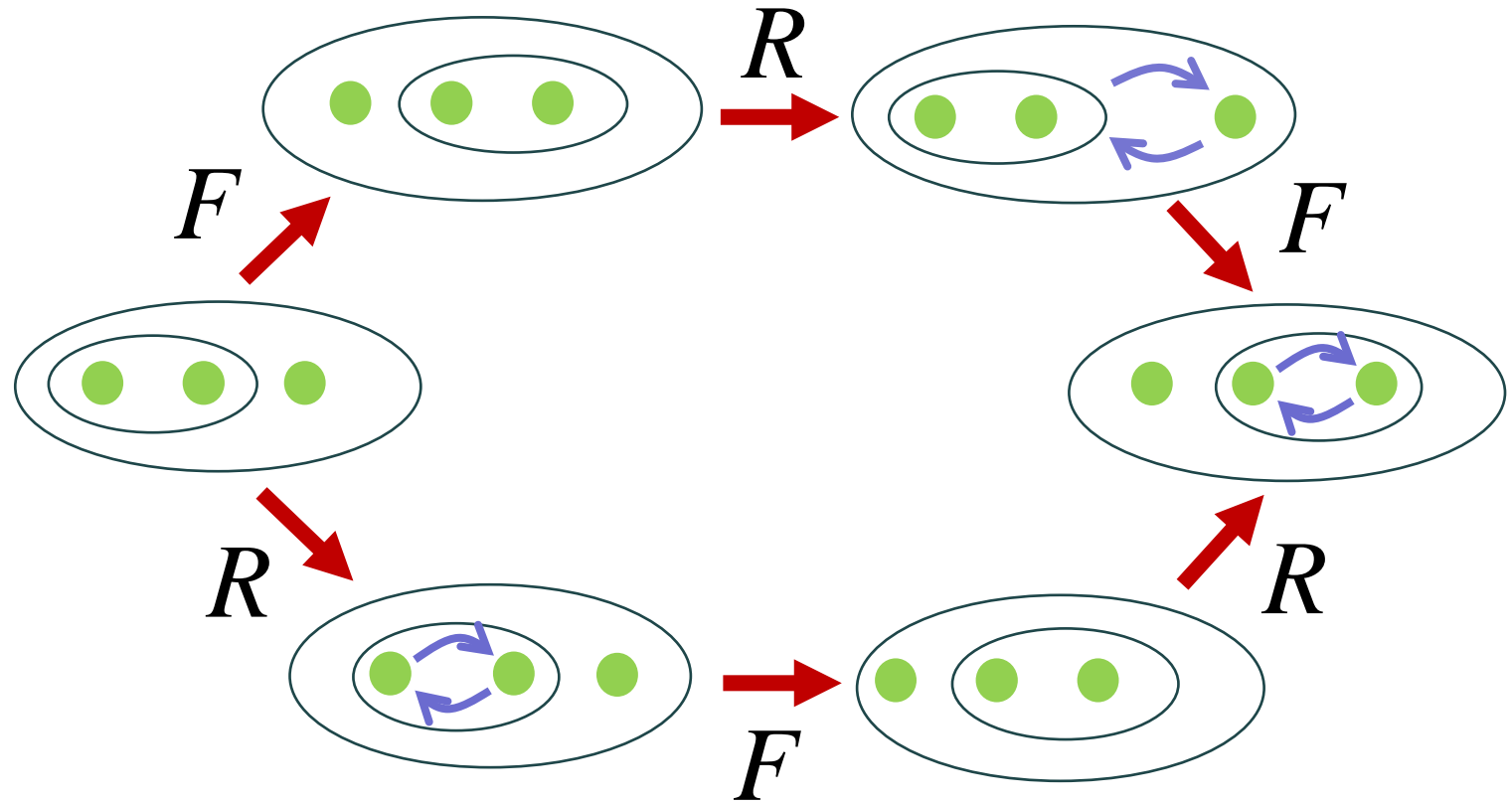




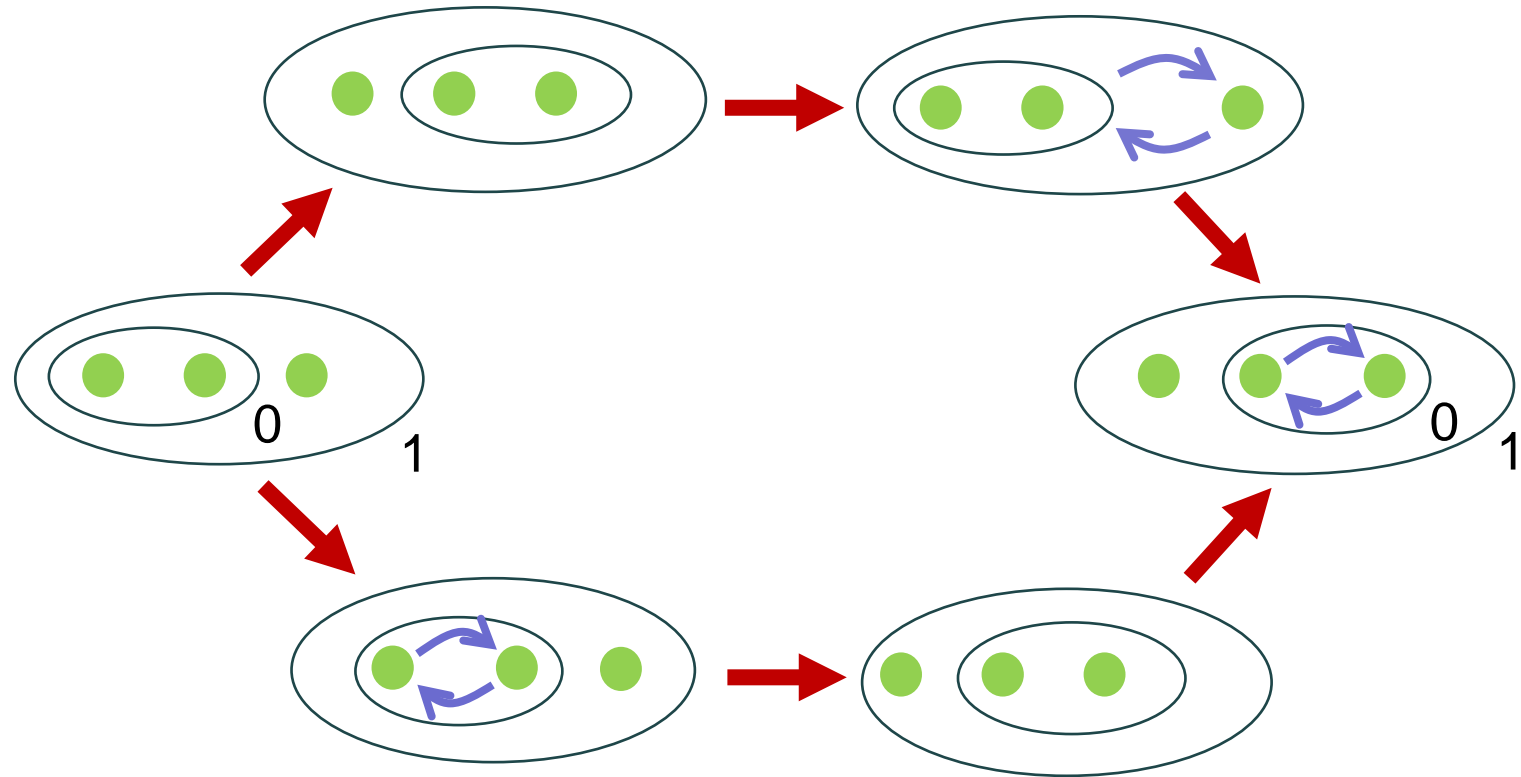




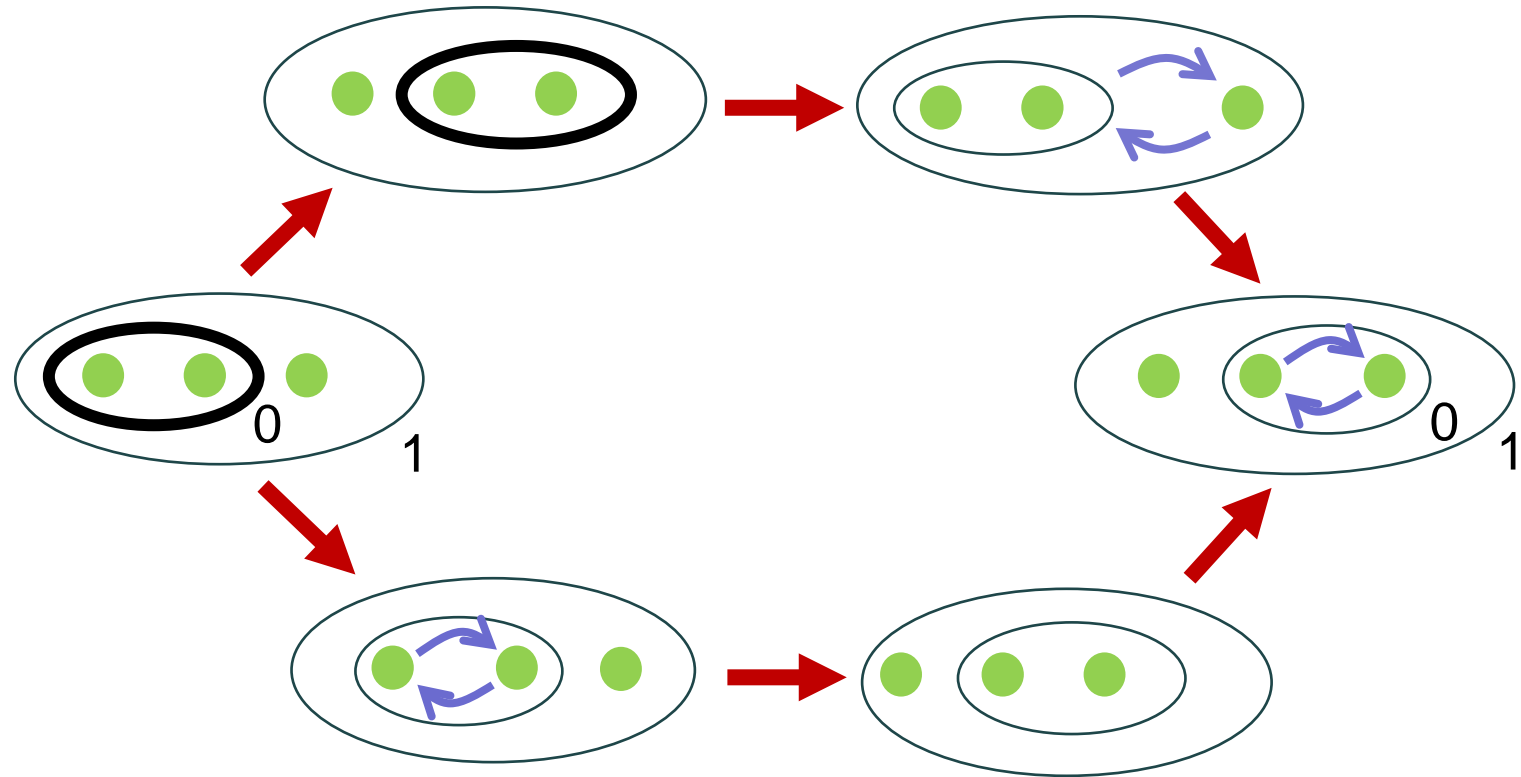
The Hexagon Equation



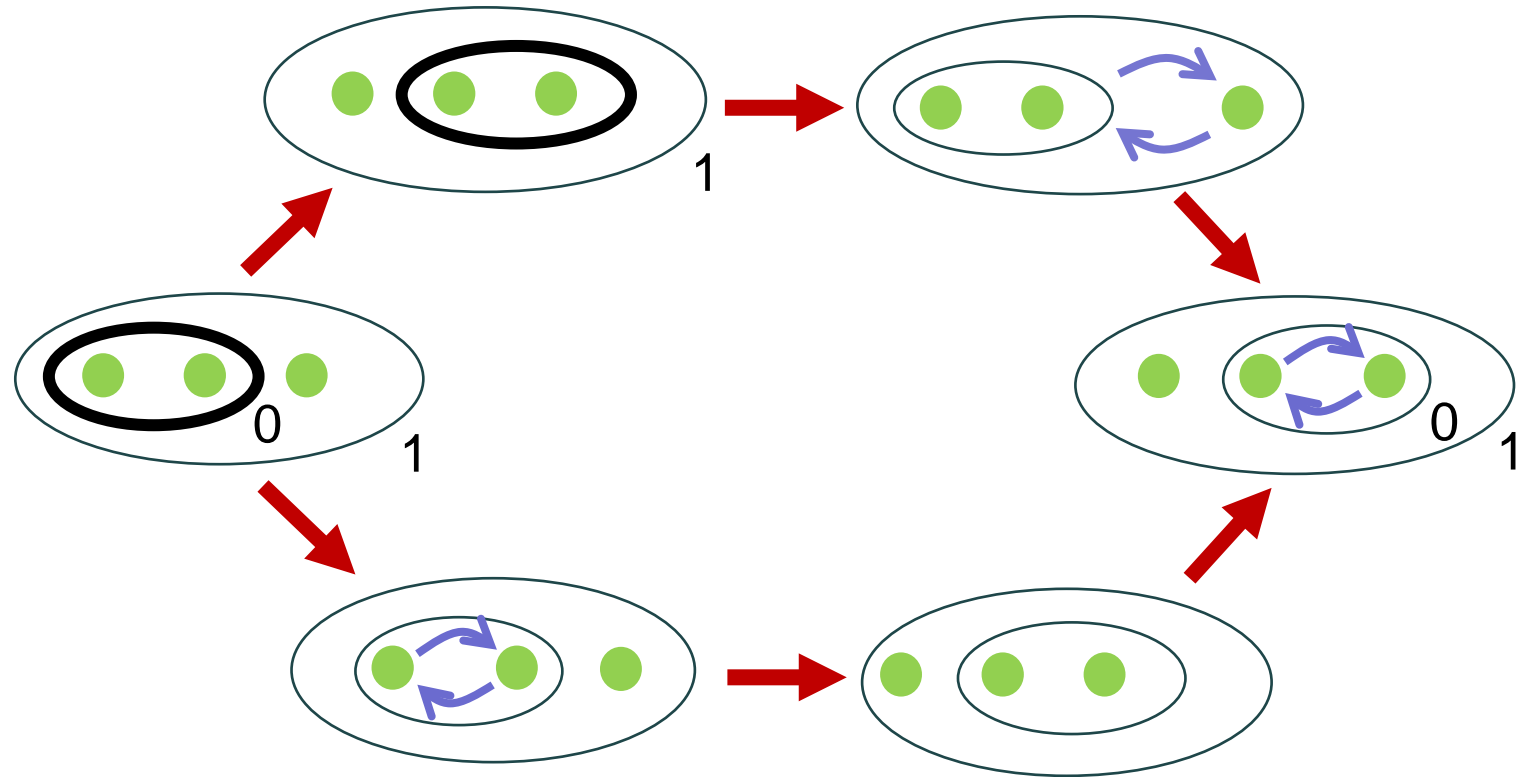
The Hexagon Equation



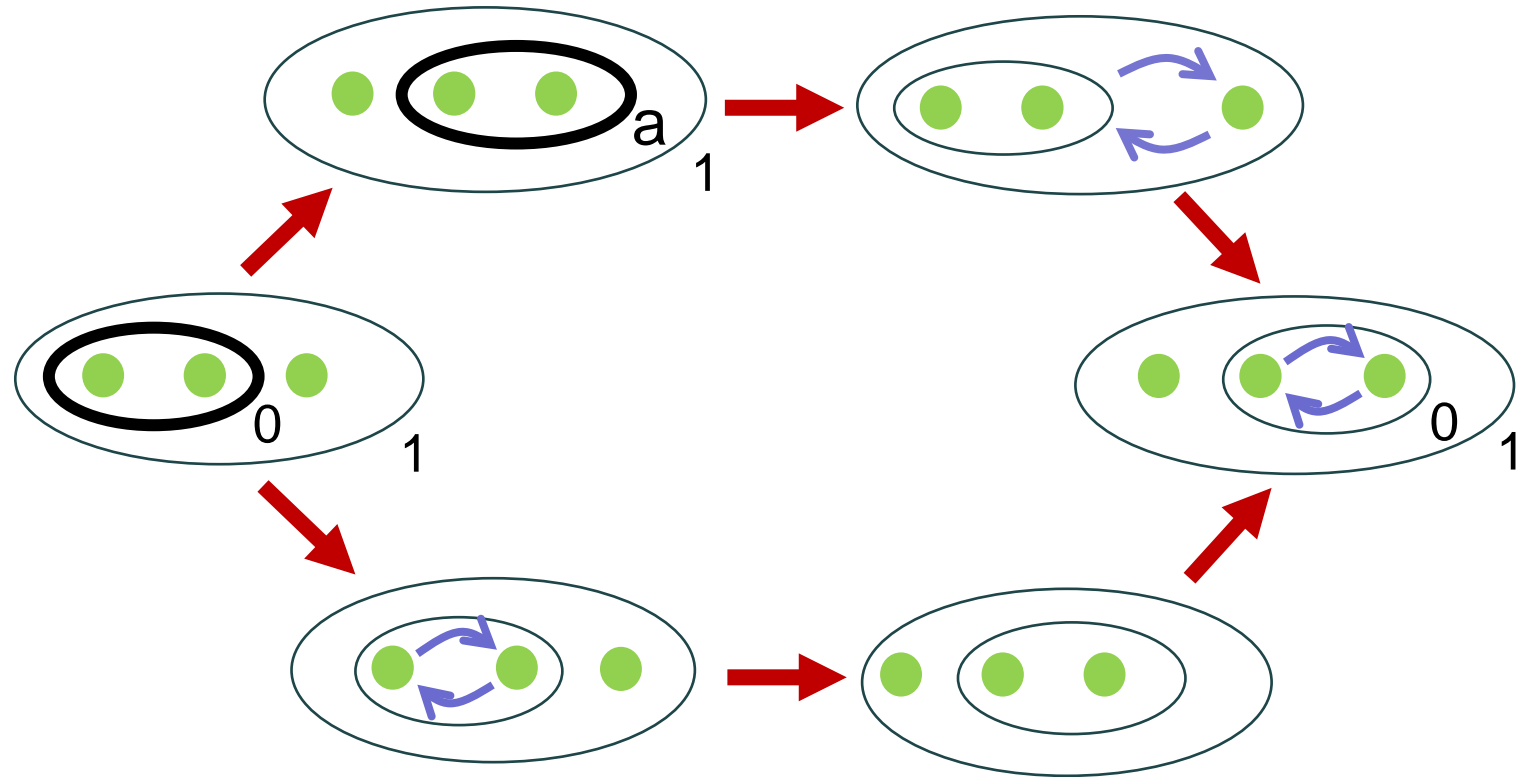
The Hexagon Equation



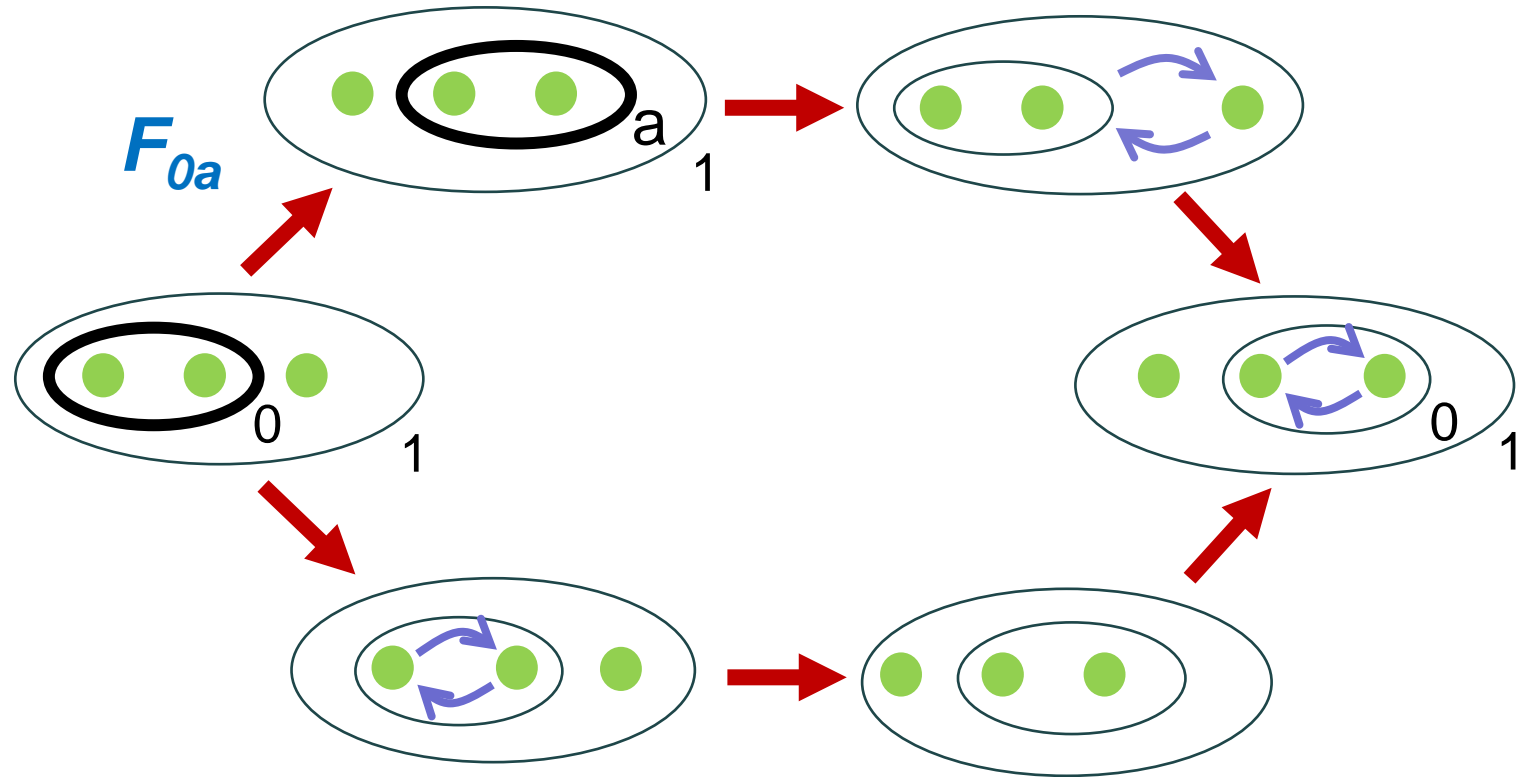
The Hexagon Equation



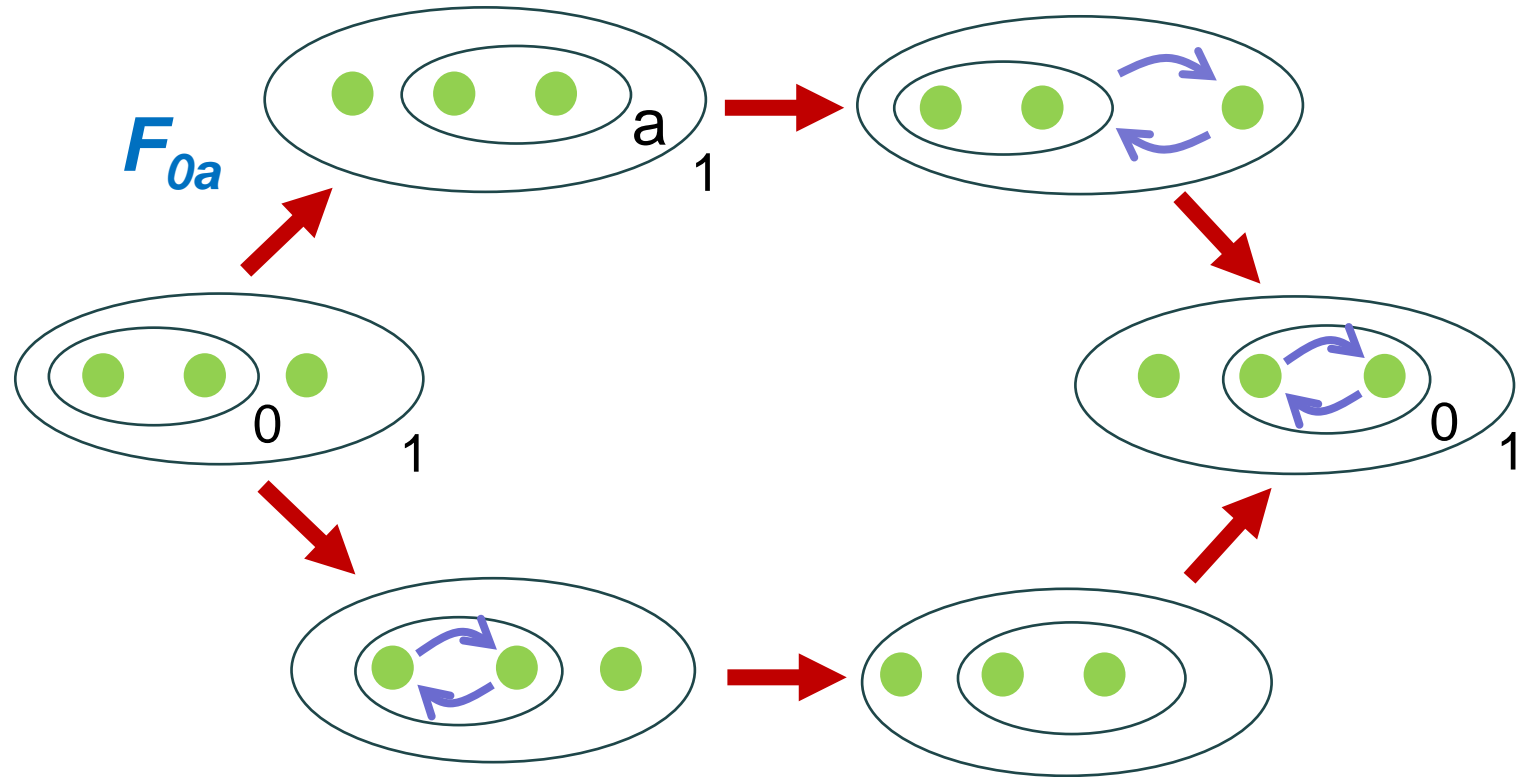
The Hexagon Equation



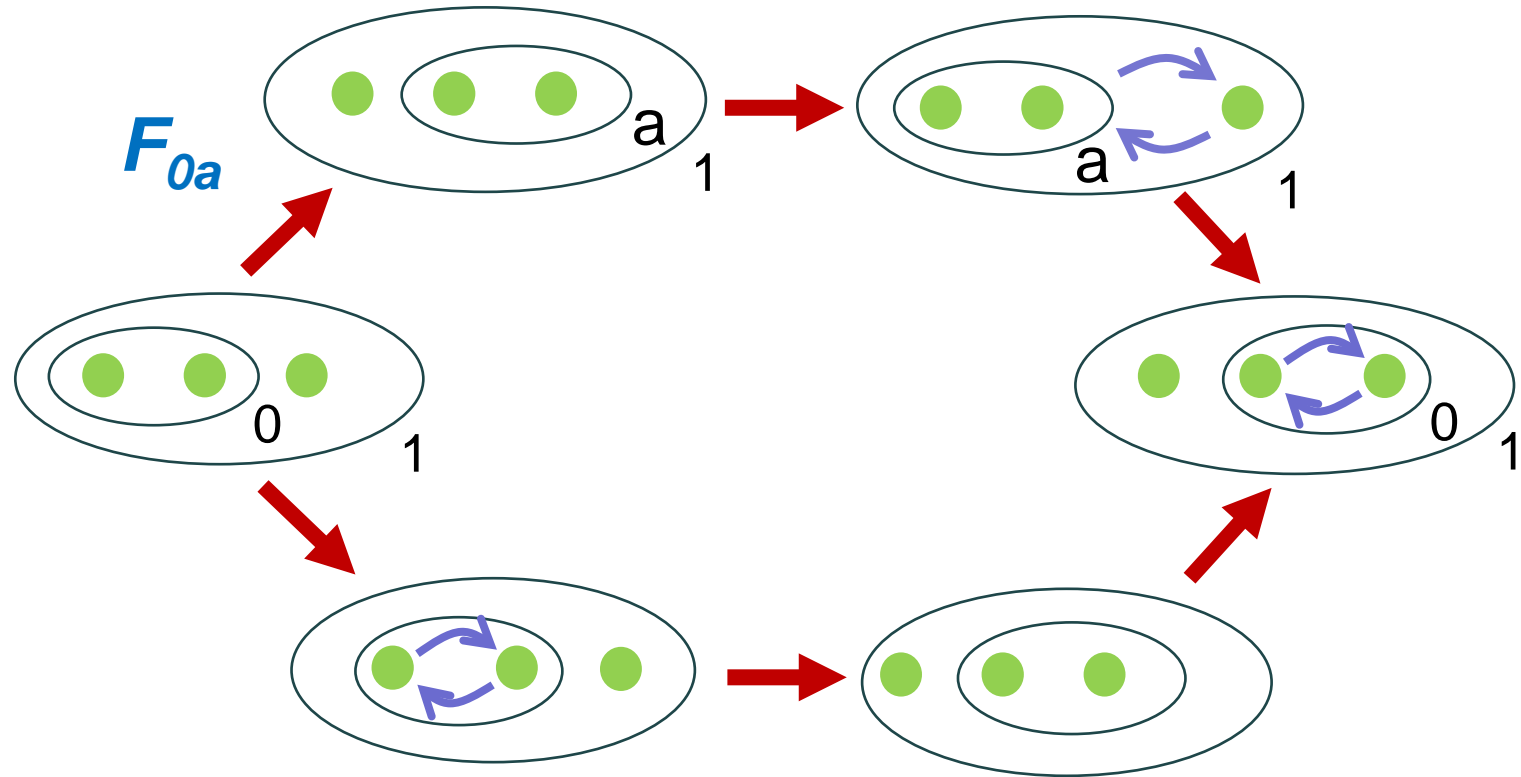
The Hexagon Equation



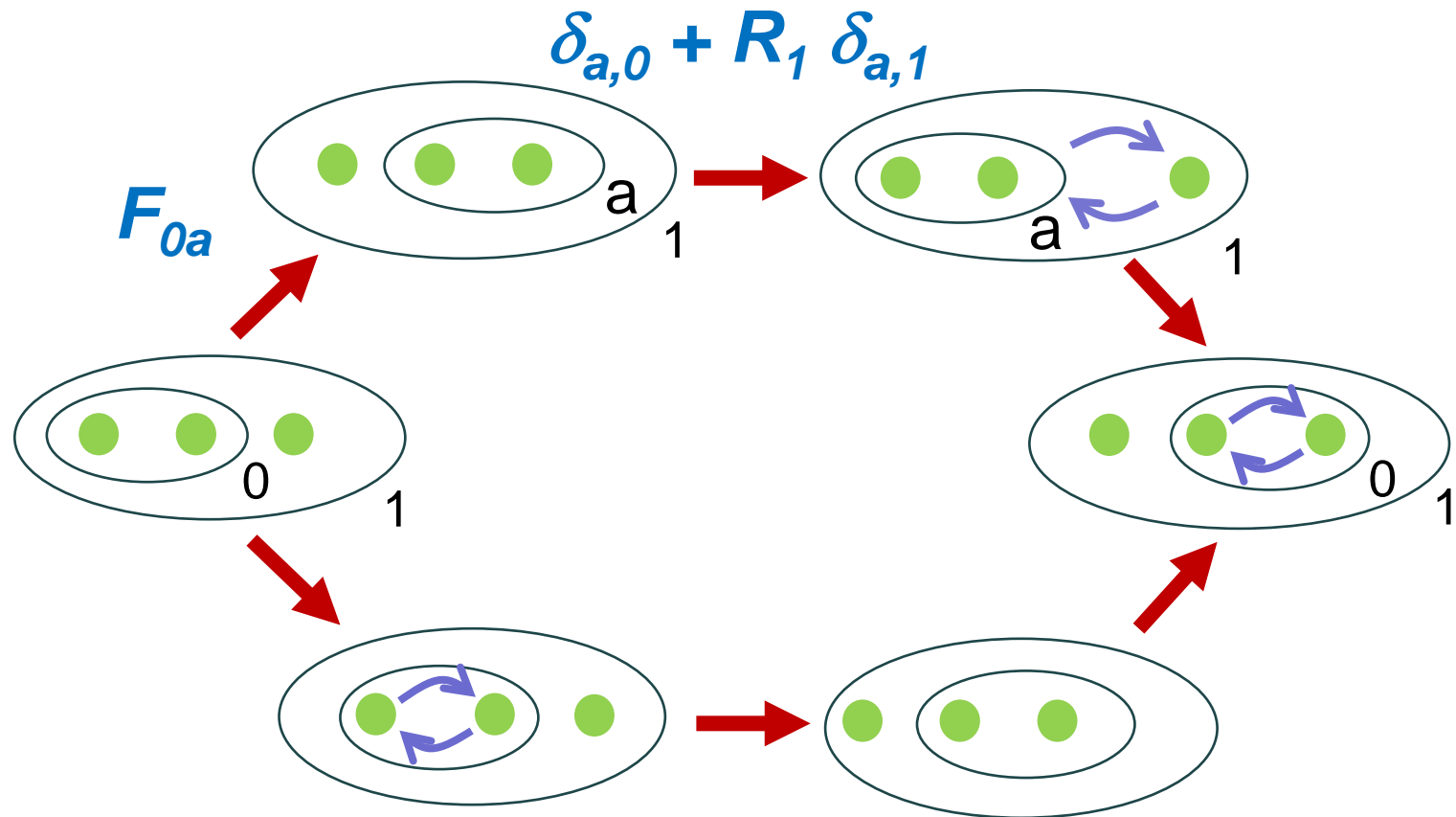
The Hexagon Equation



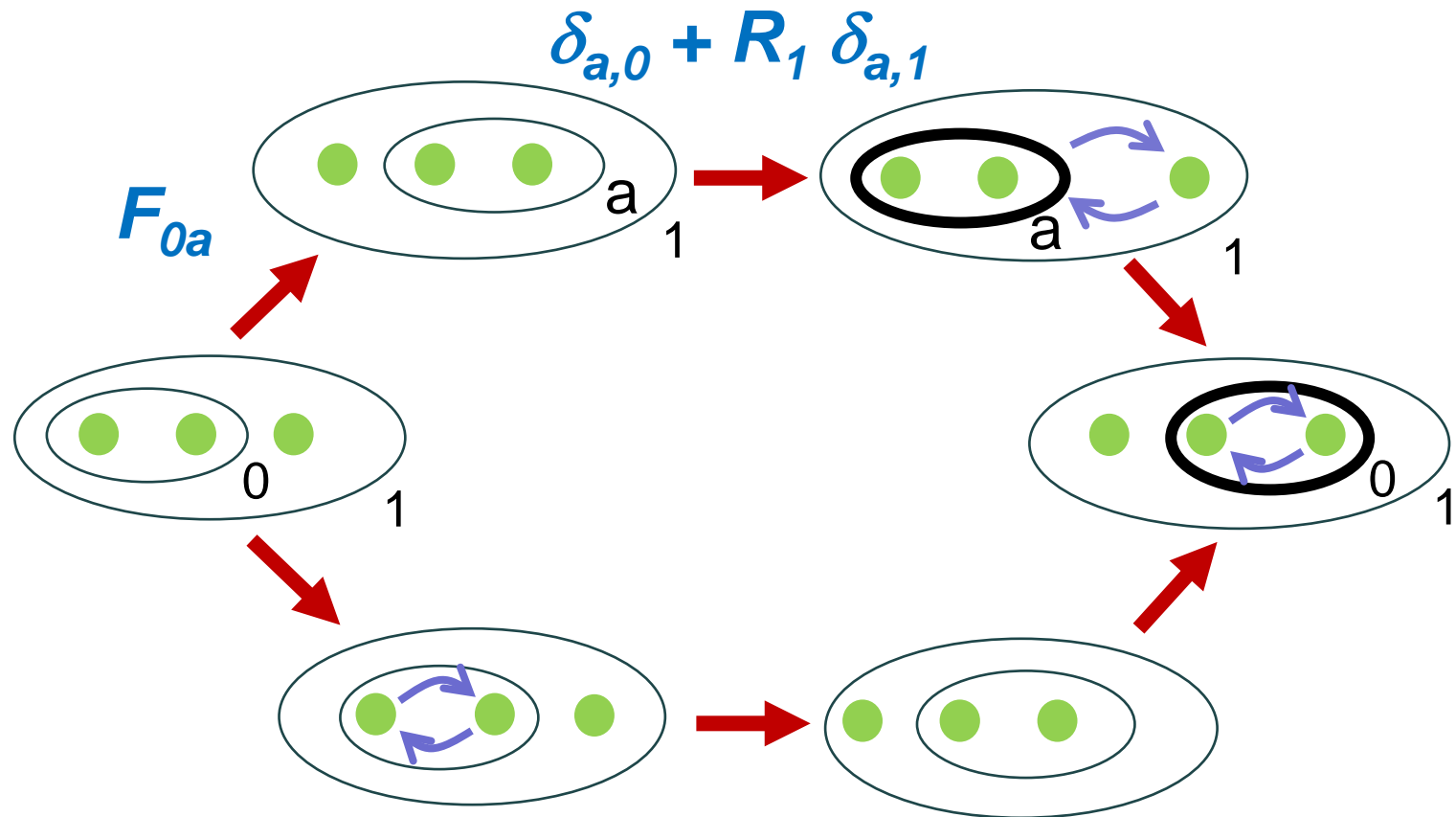
The Hexagon Equation



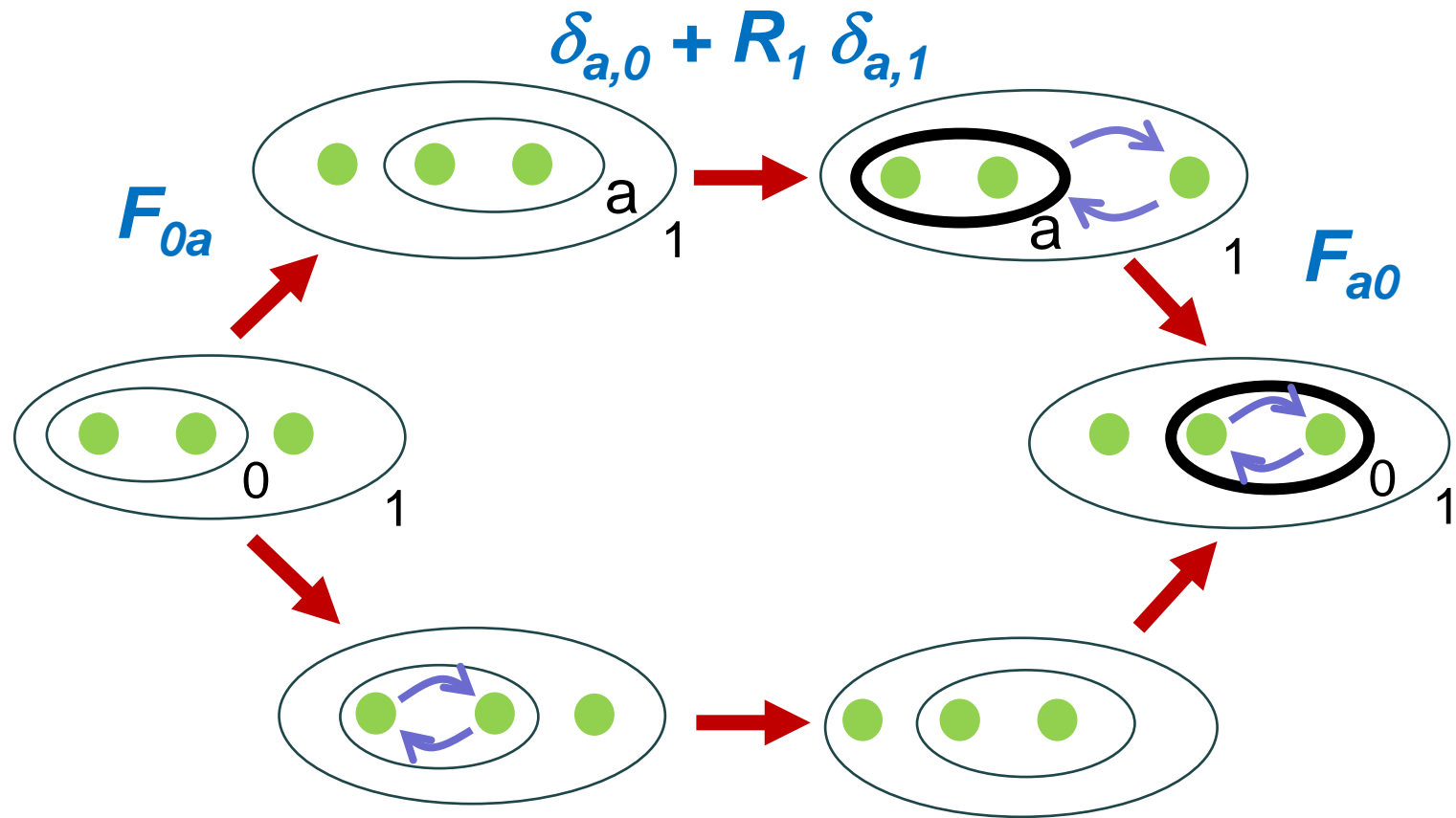
The Hexagon Equation



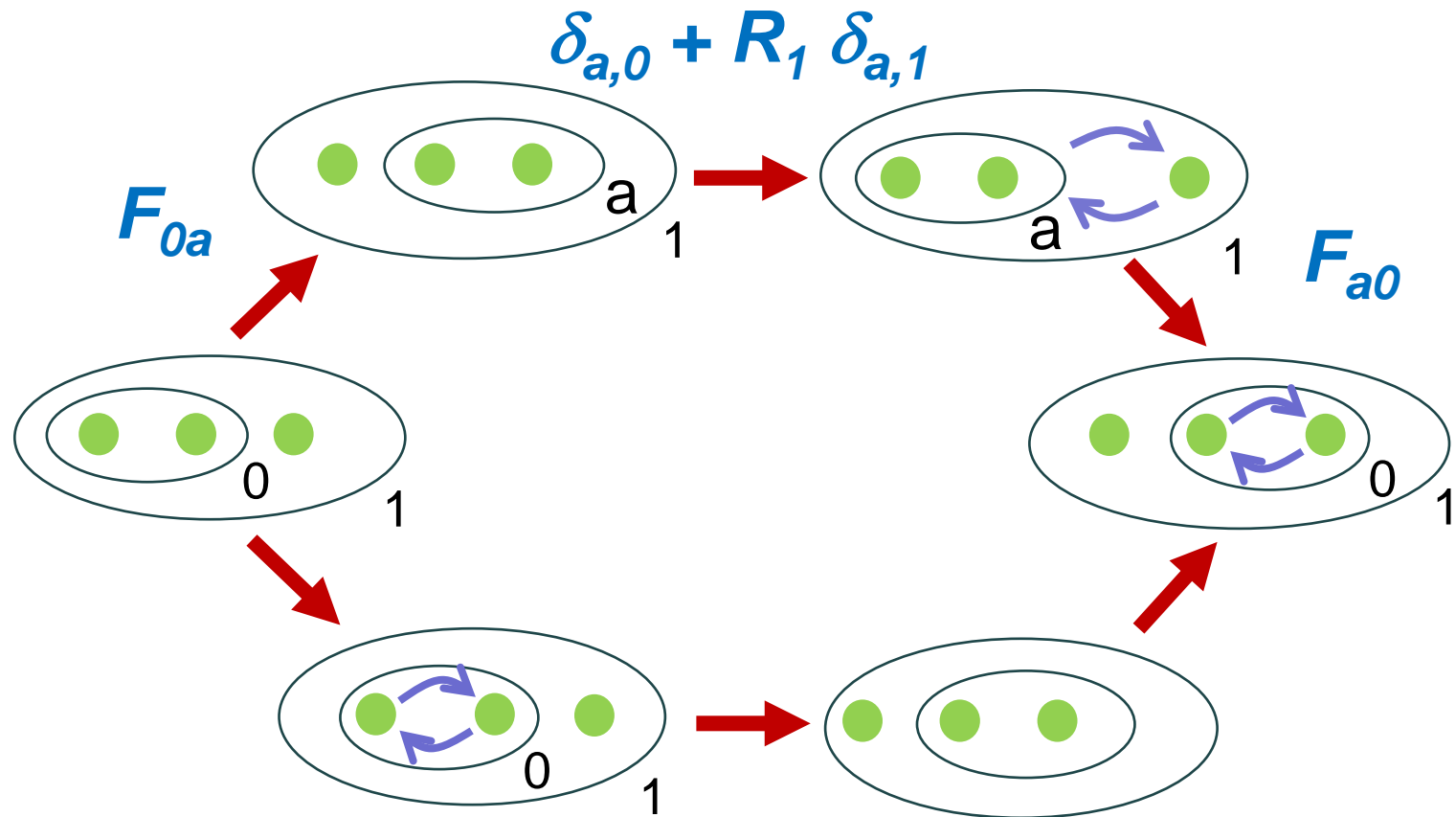
The Hexagon Equation



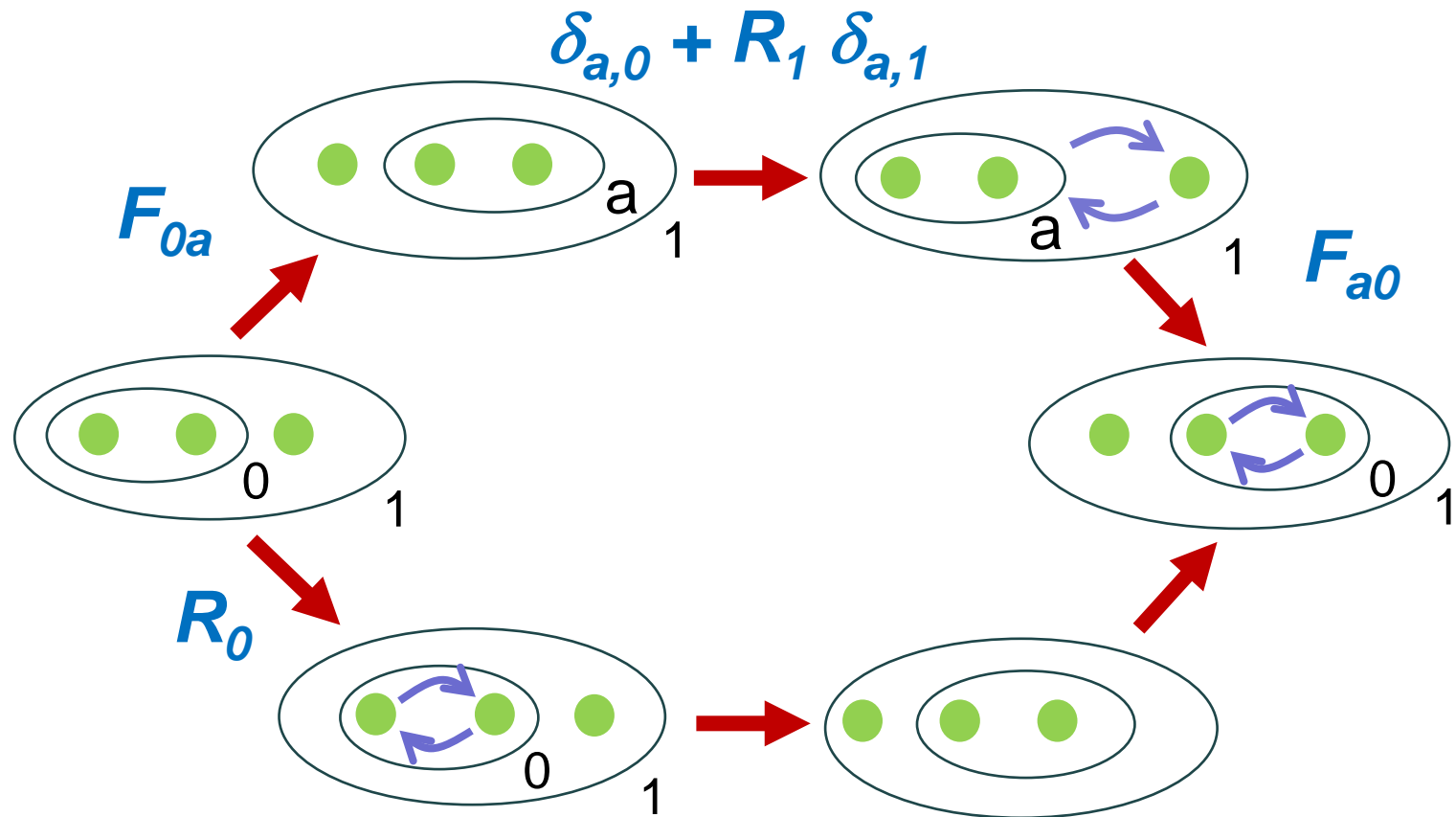
The Hexagon Equation



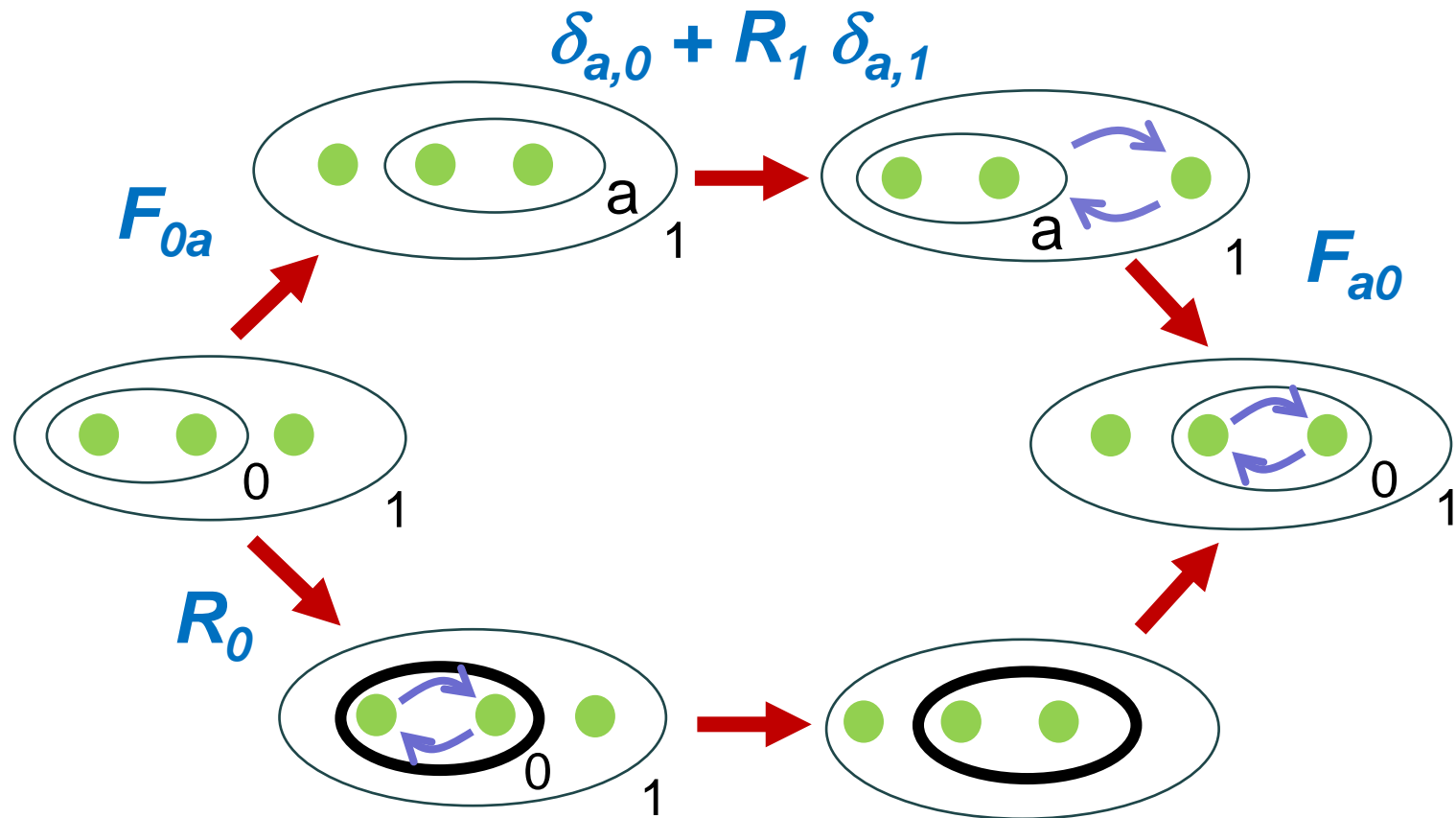
The Hexagon Equation



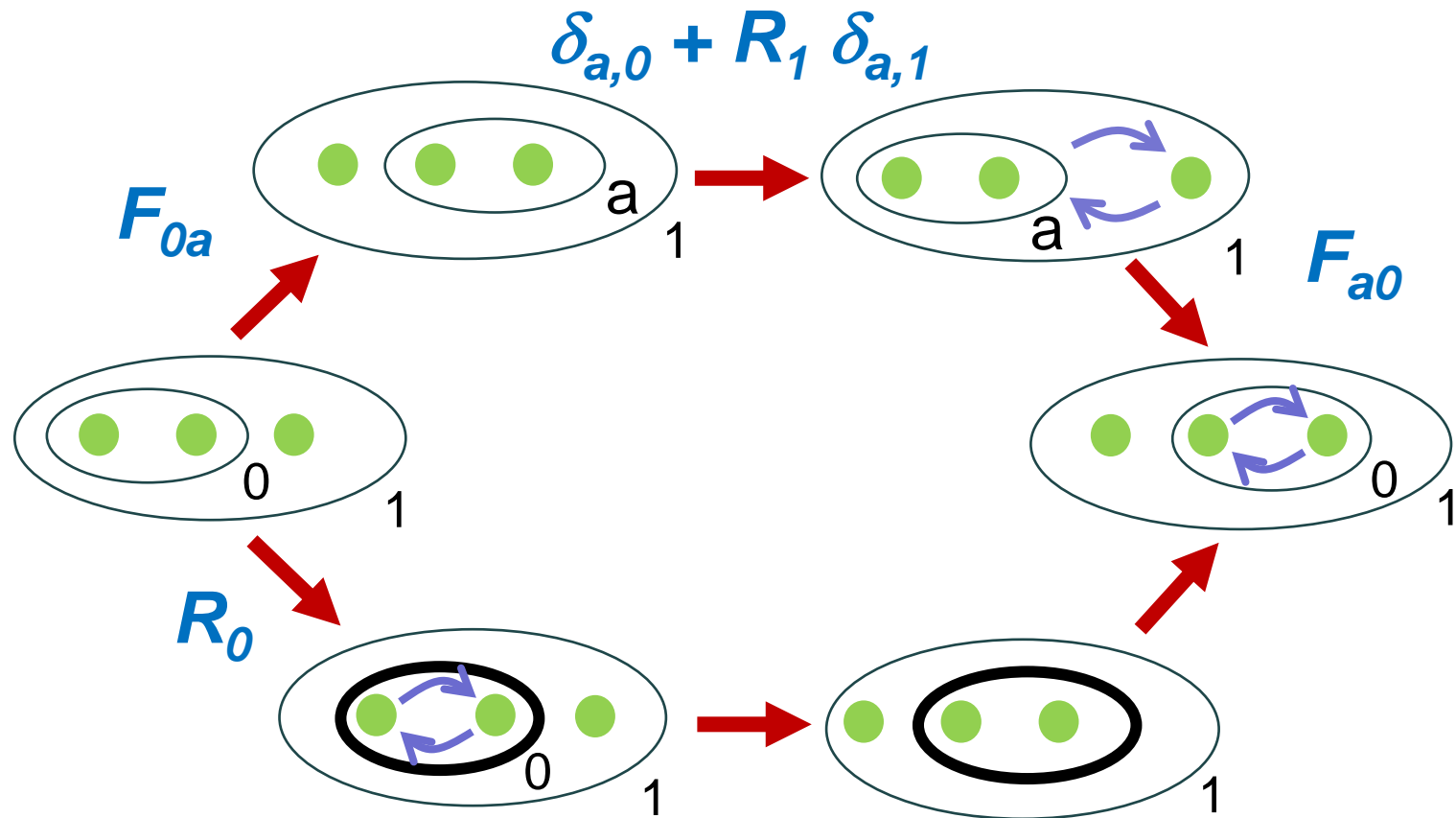
The Hexagon Equation



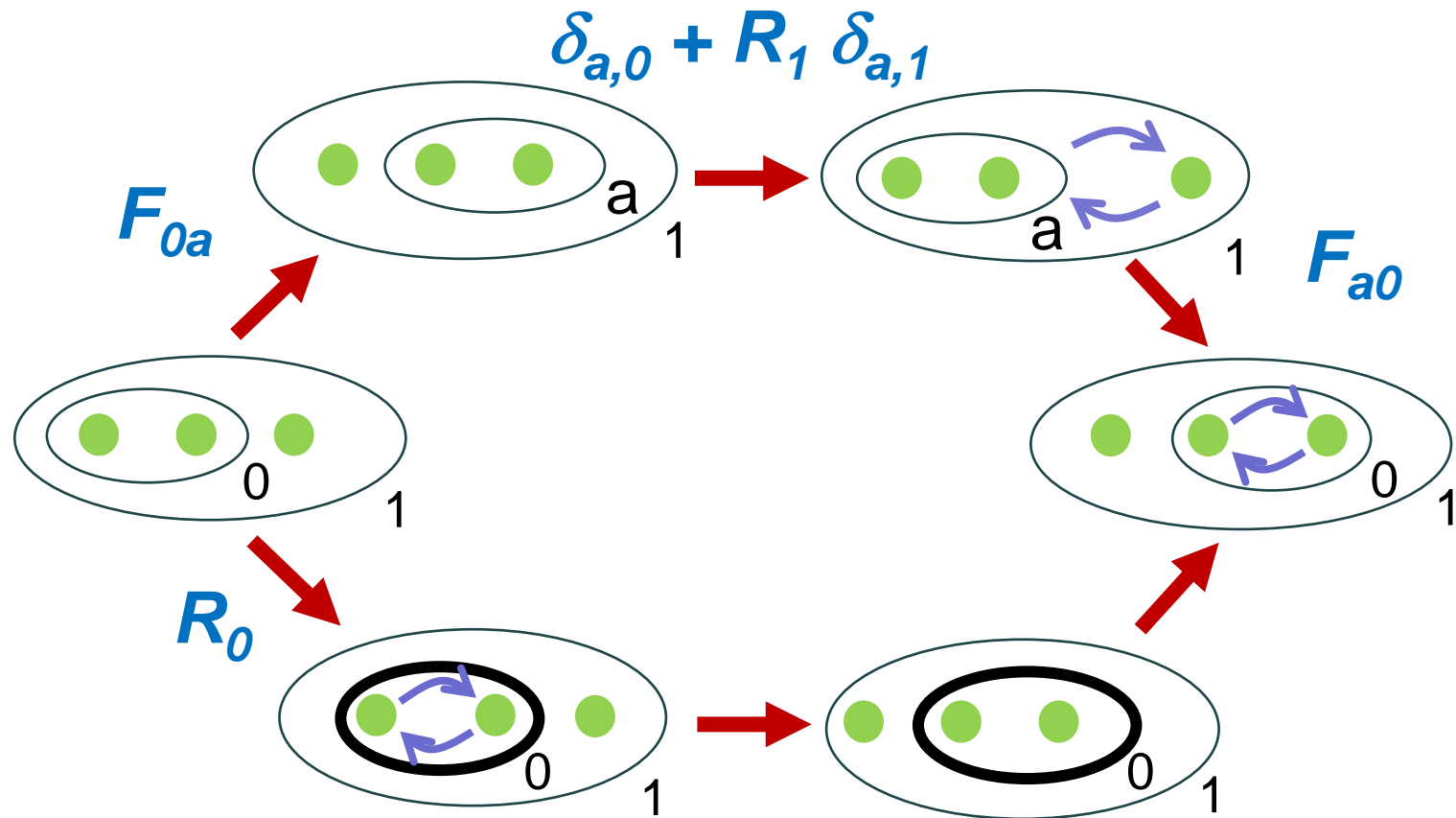
The Hexagon Equation



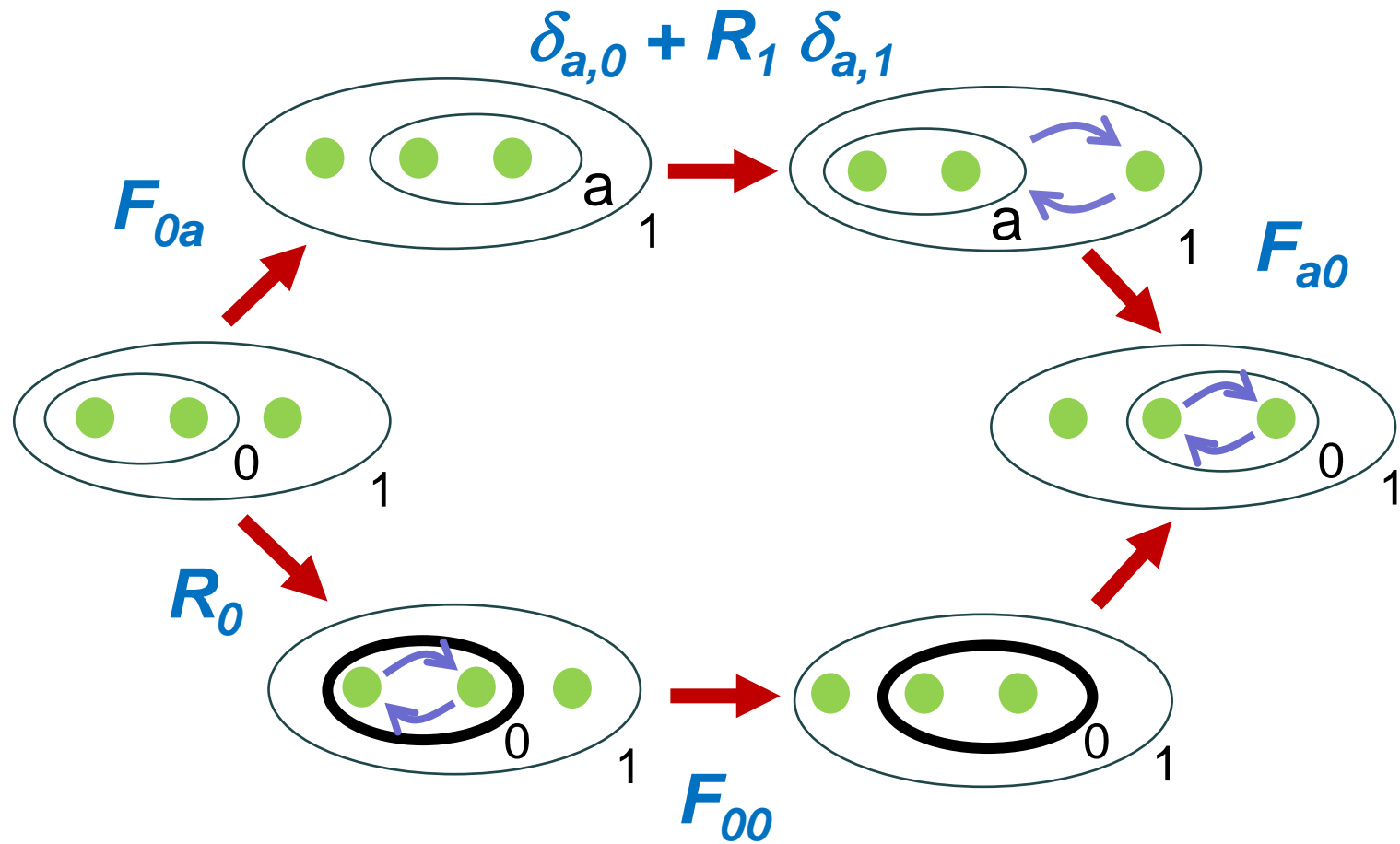
The Hexagon Equation



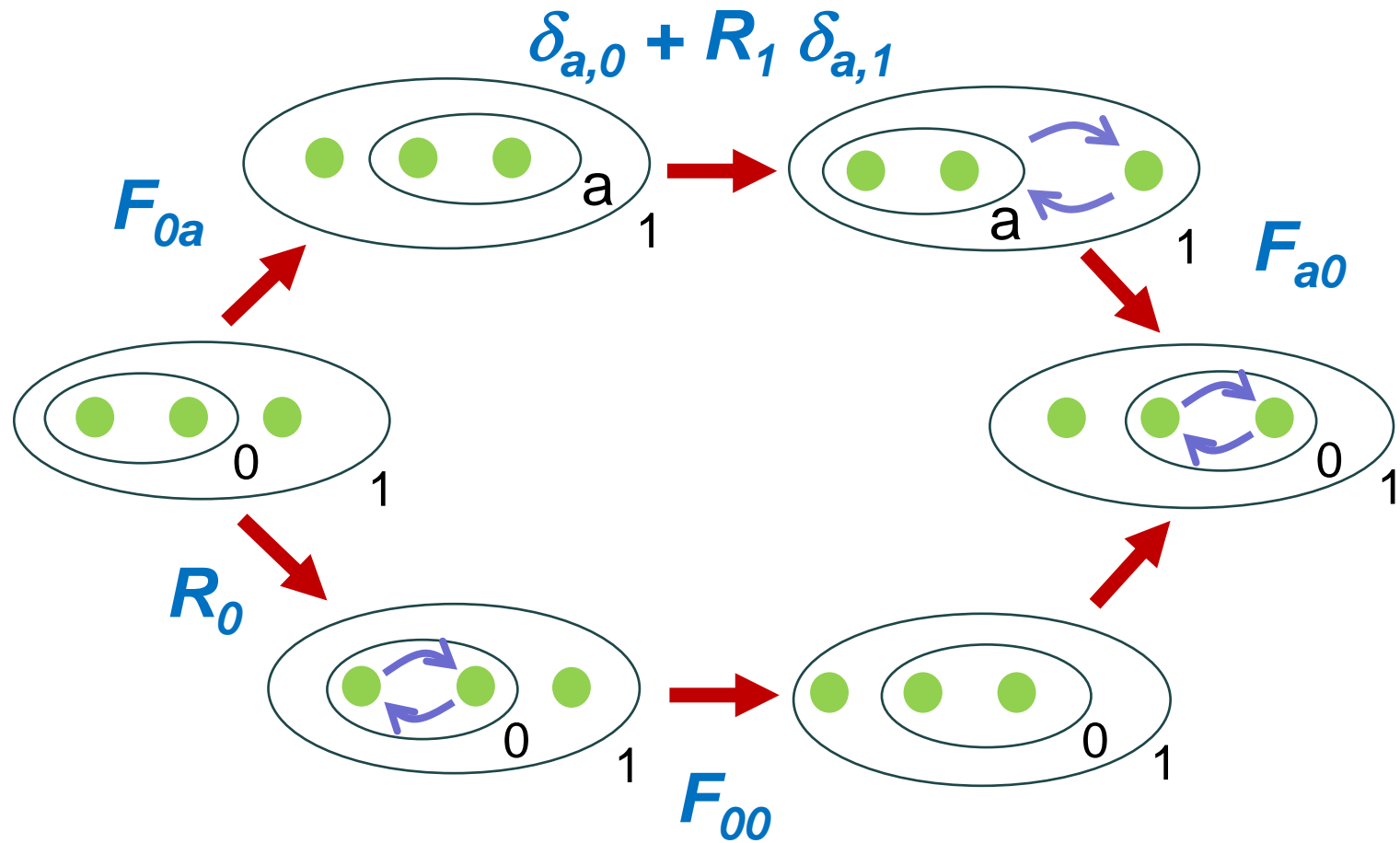
The Hexagon Equation



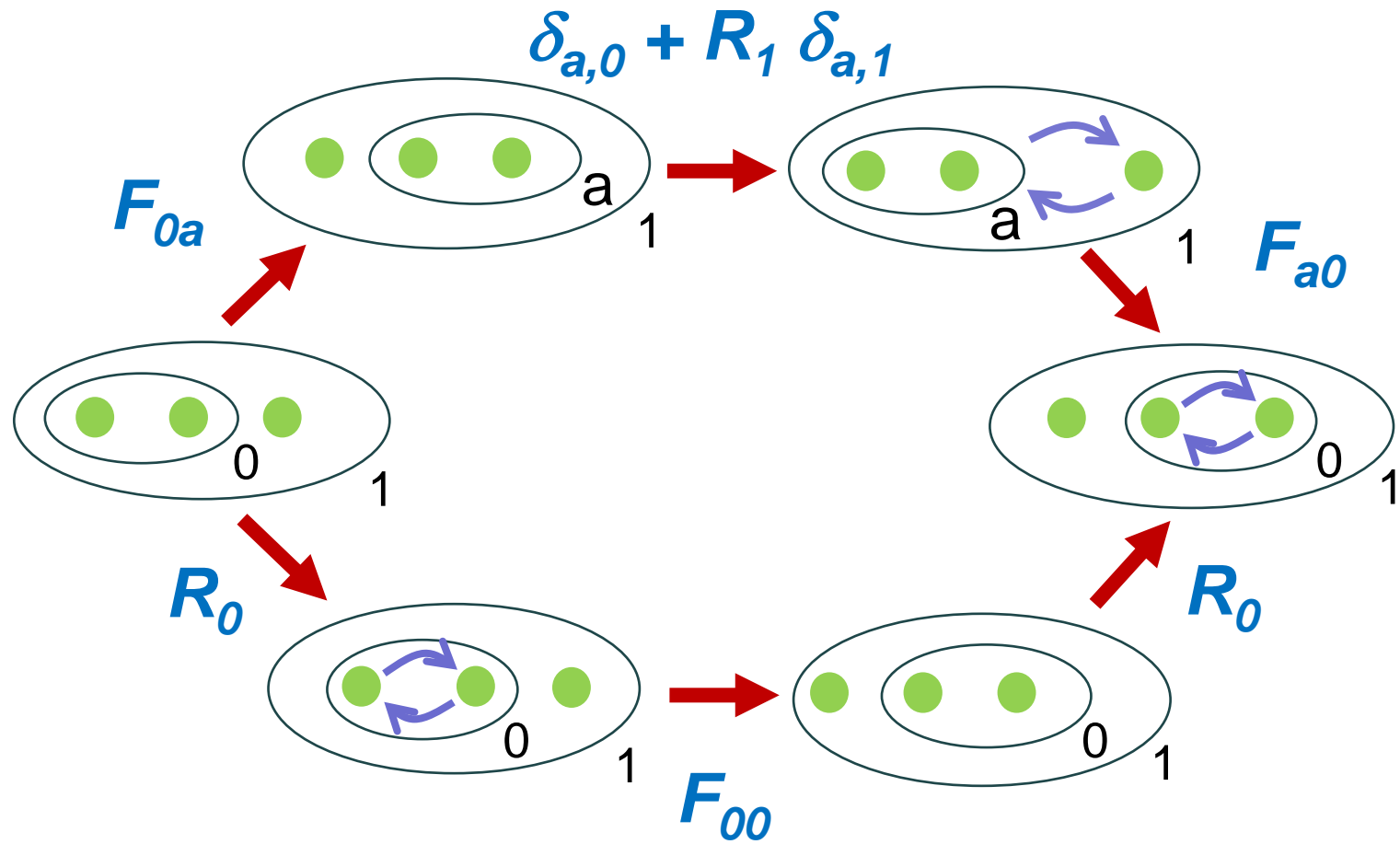
The Hexagon Equation



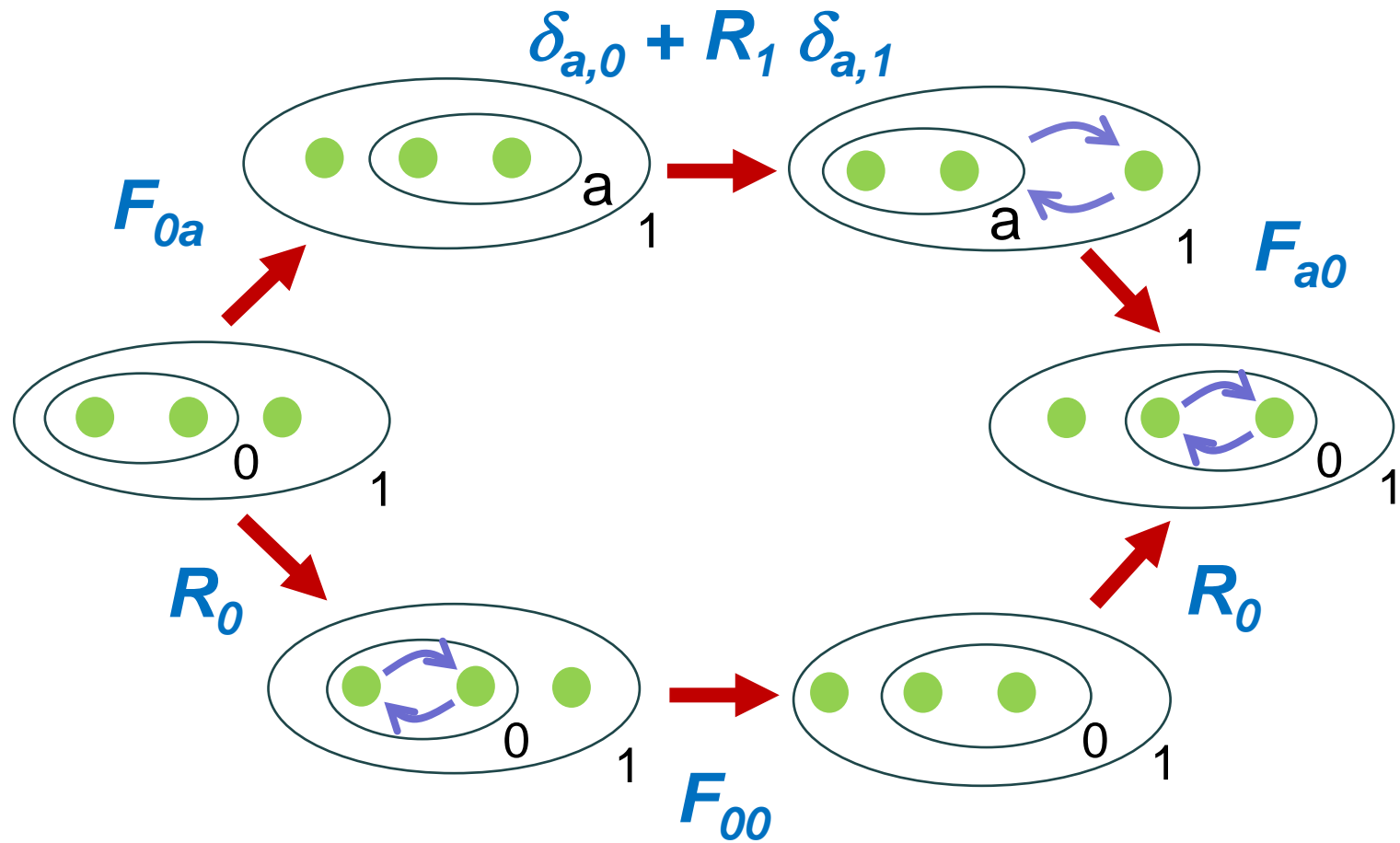
The Hexagon Equation



The Hexagon Equation

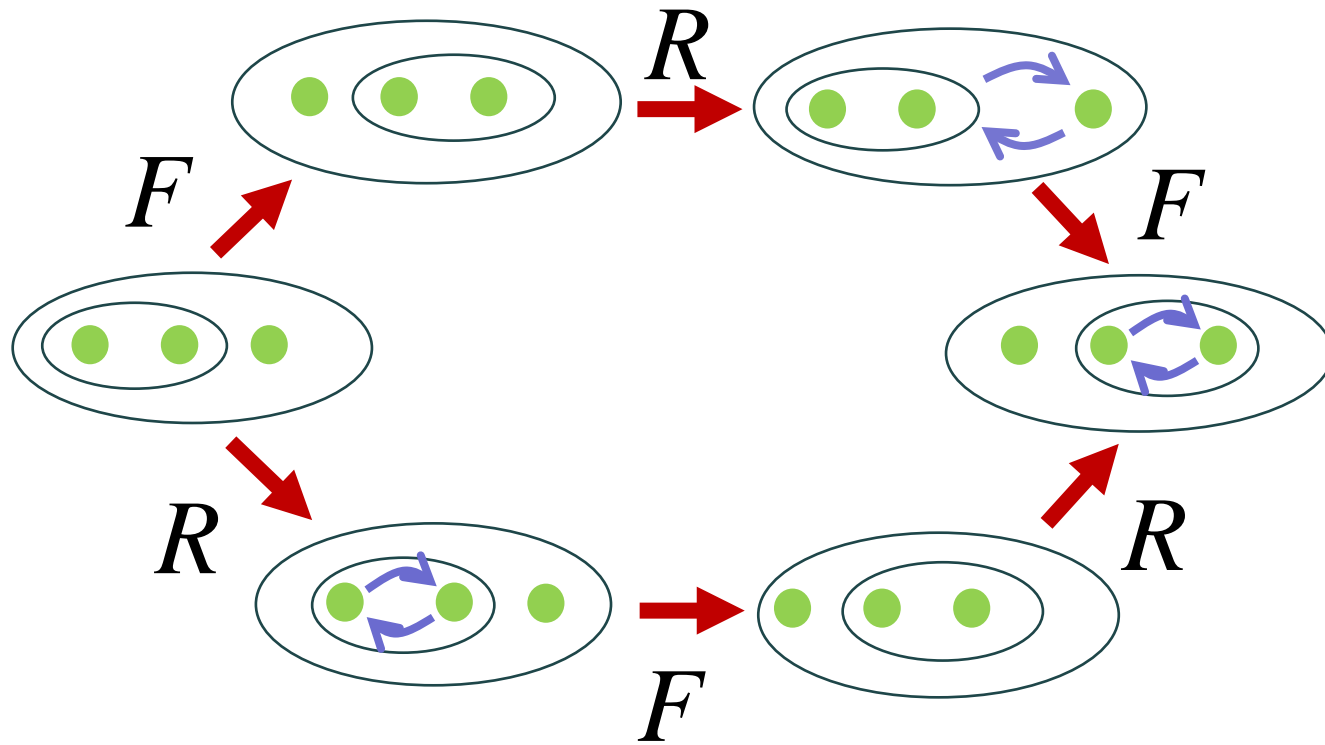


The Hexagon Equation



Top path $\rightarrow \sum_a F_{0a} (\delta_{a,0} + R_1 \delta_{a,1}) F_{a0} = R_0 F_{00} R_0 \leftarrow$ Bottom path

The Hexagon Equation



Only two solutions:

$$R = \begin{pmatrix} e^{-i4\pi/5} & 0 \\ 0 & e^{i3\pi/5} \end{pmatrix}$$

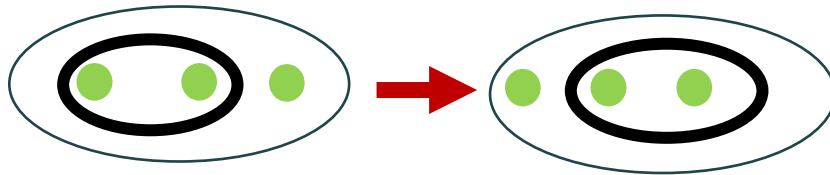
$$R = \begin{pmatrix} e^{i4\pi/5} & 0 \\ 0 & e^{-i3\pi/5} \end{pmatrix}$$

Basic Content of an Anyon Theory

“Charge” values: 0, 1

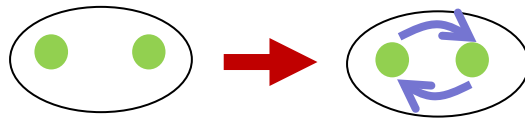
Fusion rules: $1 \times 1 = 0 + 1$; $1 \times 0 = 0 \times 1 = 1$; $0 \times 0 = 0$

F Matrix:



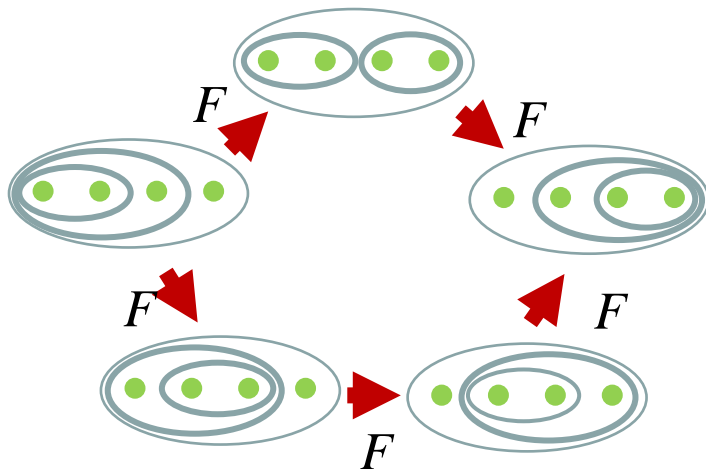
$$F = \begin{pmatrix} \varphi^{-1} & \varphi^{-1/2} \\ \varphi^{-1/2} & -\varphi^{-1} \end{pmatrix}$$

R Matrix:

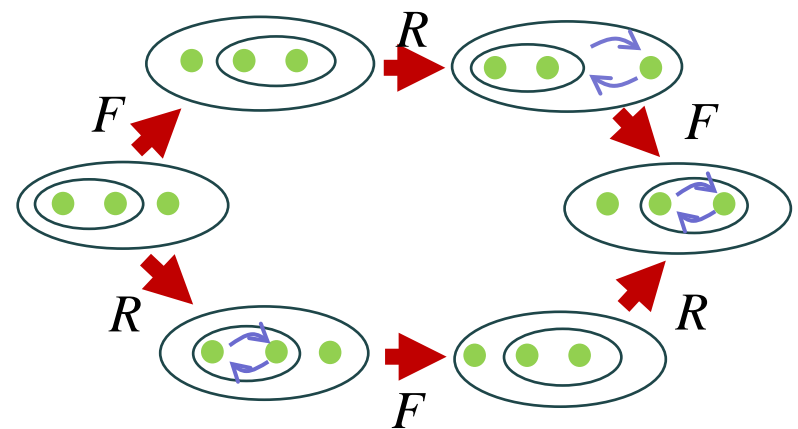


$$R = \begin{pmatrix} e^{-i4\pi/5} & 0 \\ 0 & e^{i3\pi/5} \end{pmatrix}$$

Pentagon

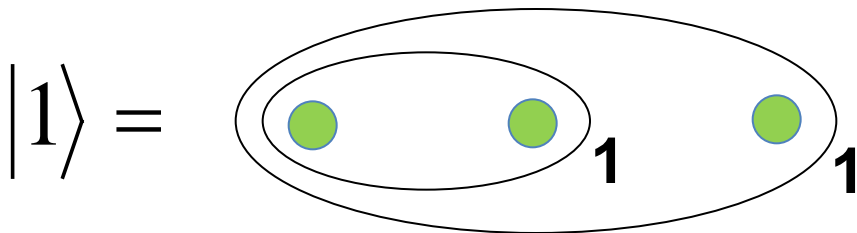
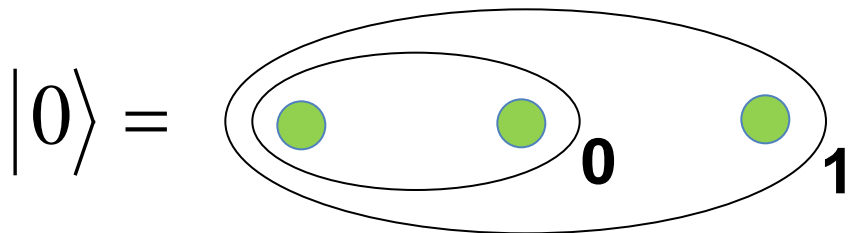


Hexagon

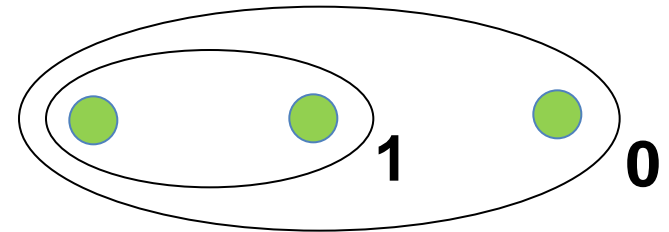


Qubit Encoding

Qubit States



Non-Computational State

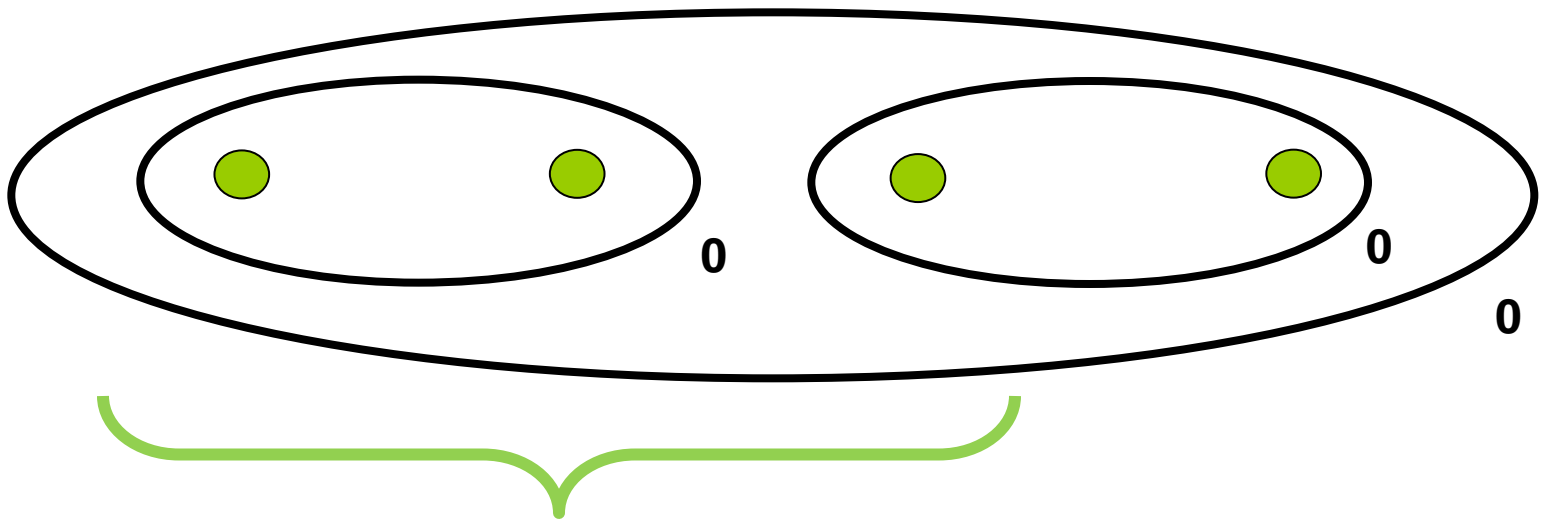


State of qubit is determined by
“charge” of two leftmost
particles

Transitions to this state are
leakage errors

Initializing a Qubit

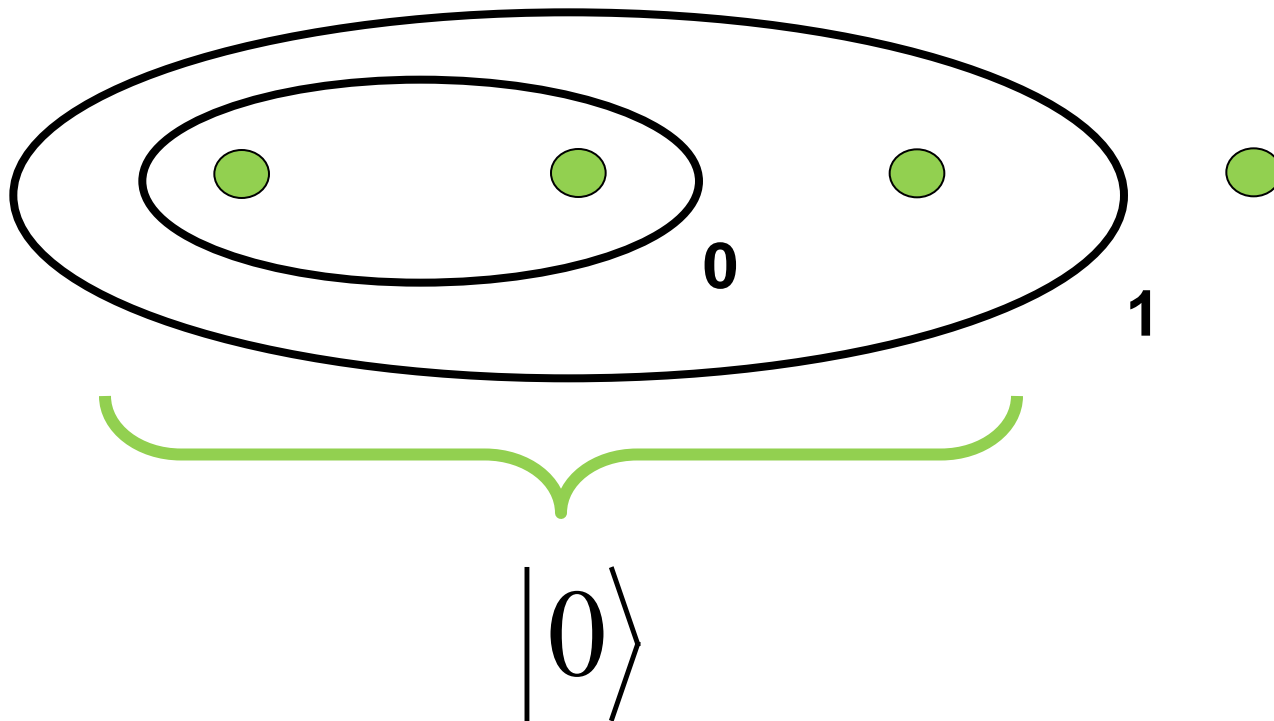
Pull two quasiparticle-quasihole pairs out of the “vacuum”.



These three particles have total q-spin 1

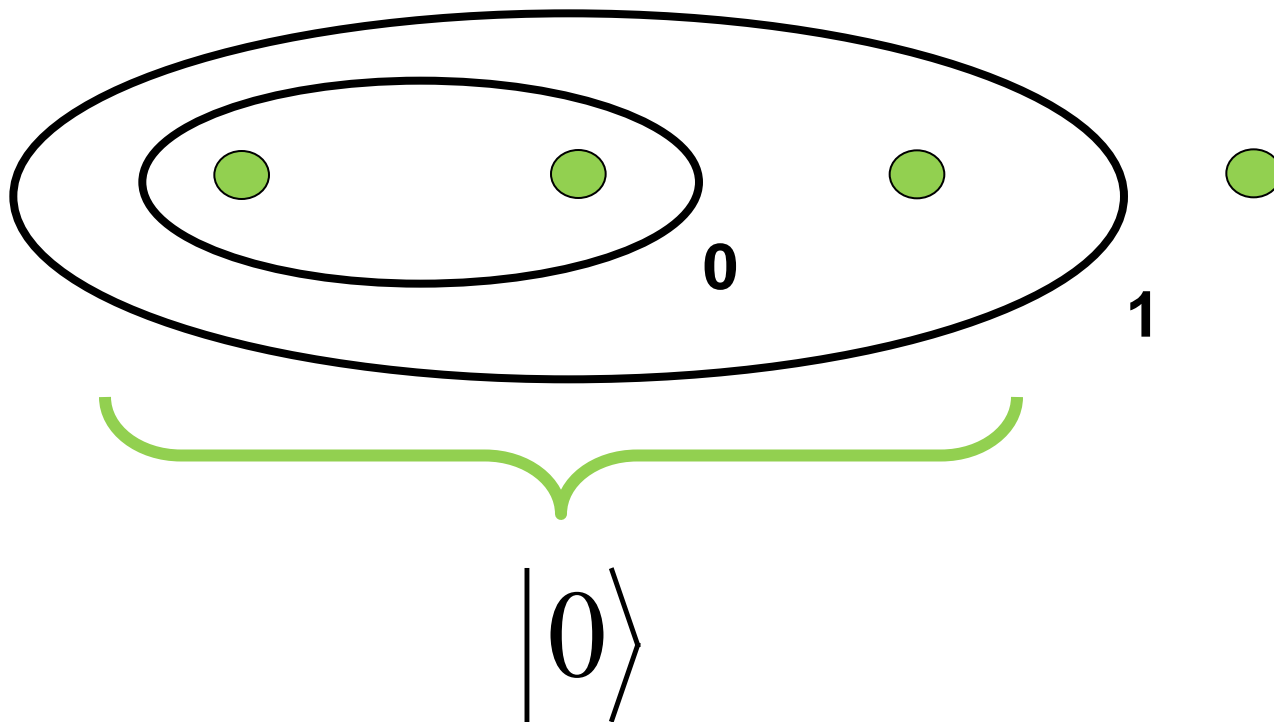
Initializing a Qubit

Pull two quasiparticle-quasihole pairs out of the “vacuum”.



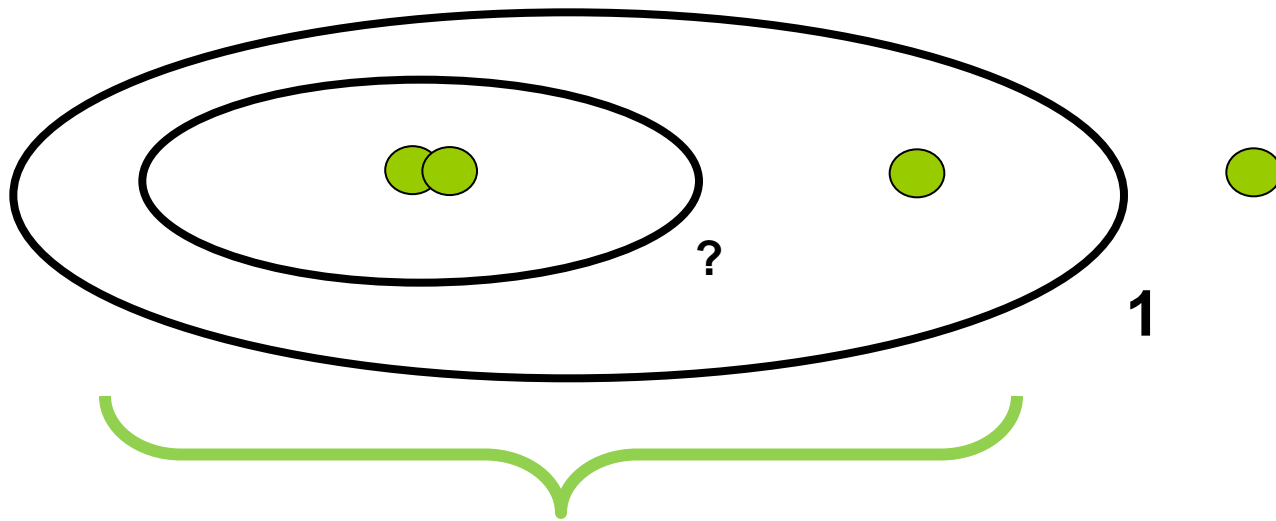
Initializing a Qubit

Pull two quasiparticle-quasihole pairs out of the “vacuum”.



Measuring a Qubit

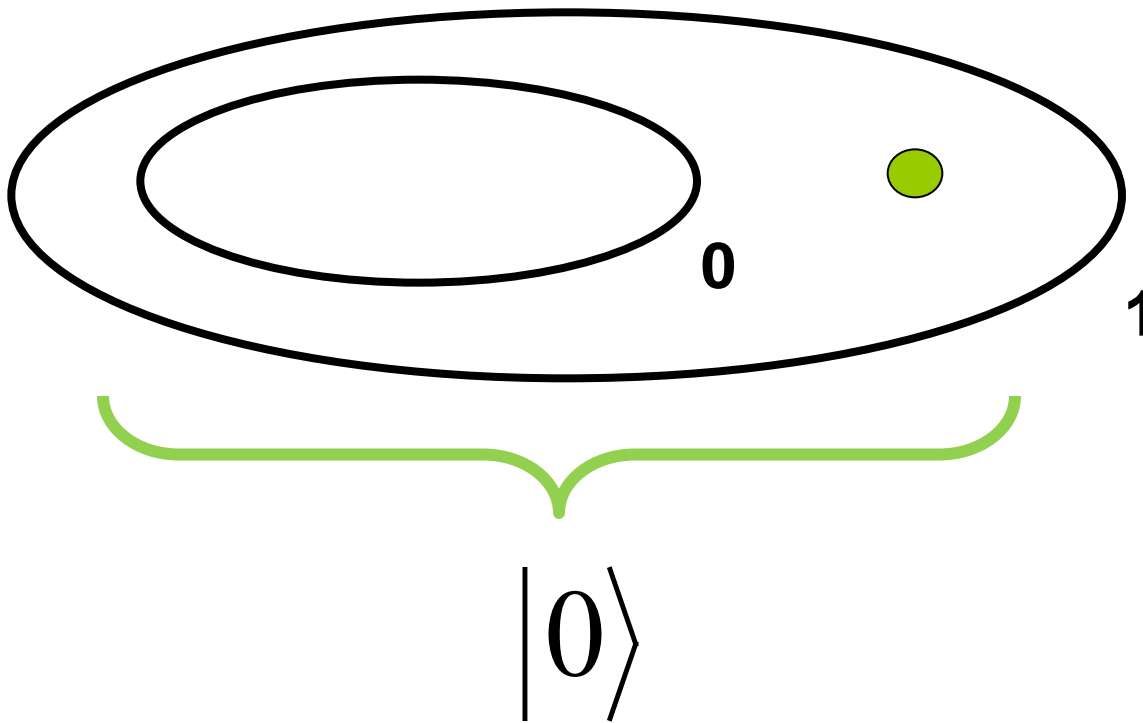
Try to fuse the leftmost quasiparticle-quasihole pair.



$$\alpha|0\rangle + \beta|1\rangle$$

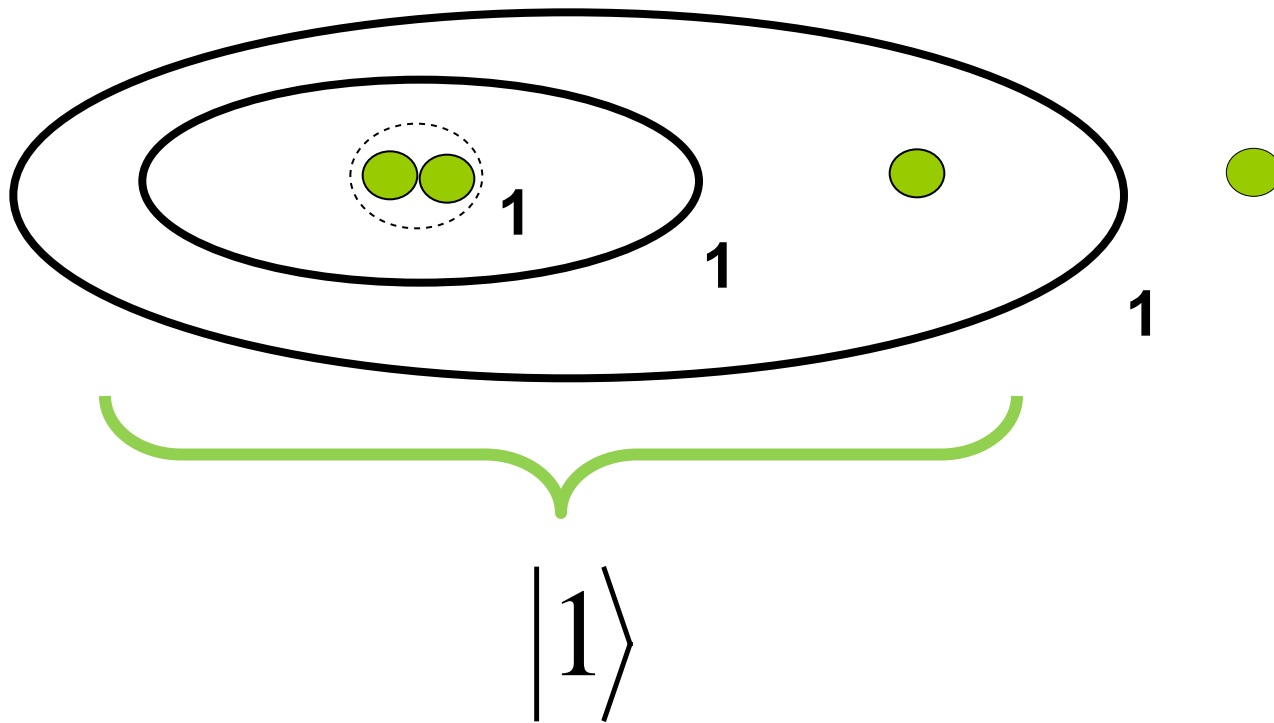
Measuring a Qubit

If they fuse back into the “vacuum” the result of the measurement is 0.

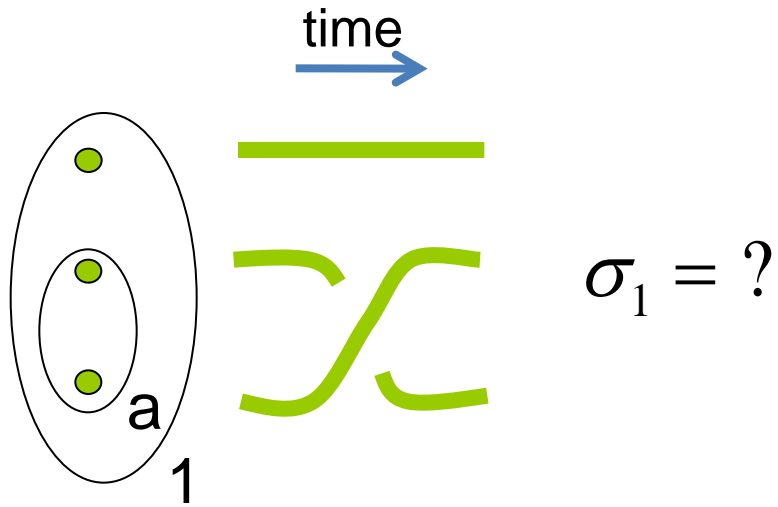


Measuring a Qubit

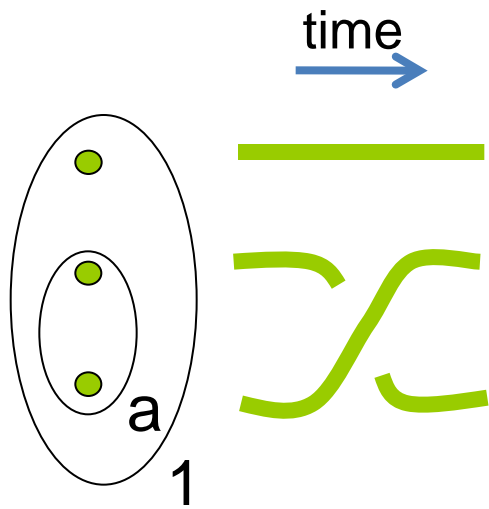
If they cannot fuse back into the “vacuum” the result of the measurement is 1



Elementary Braid Matrices



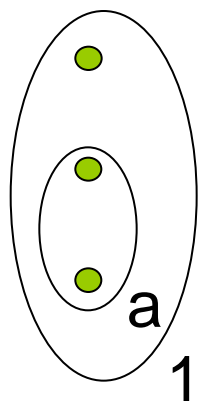
Elementary Braid Matrices



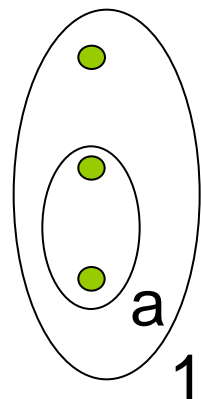
$$\sigma_1 = R = \begin{matrix} a = & 0 & 1 \\ \left(\begin{array}{cc} e^{-i4\pi/5} & 0 \\ 0 & e^{i3\pi/5} \end{array} \right) \end{matrix}$$

Elementary Braid Matrices

time
→

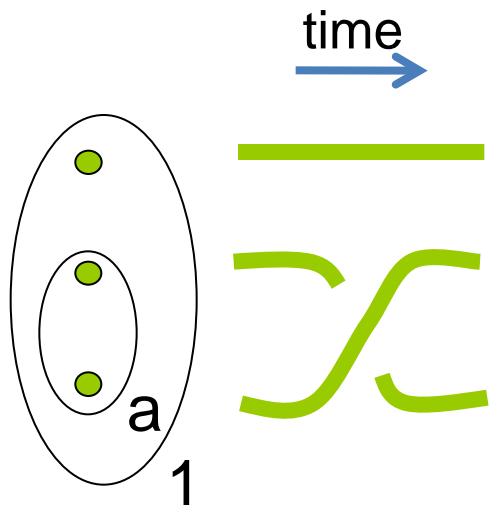


$$\sigma_1 = R = \begin{pmatrix} e^{-i4\pi/5} & 0 \\ 0 & e^{i3\pi/5} \end{pmatrix}$$

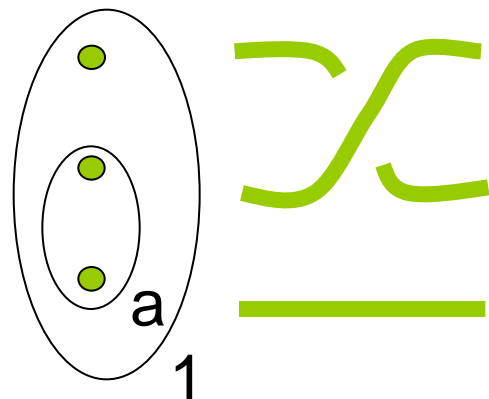


$$\sigma_2 = ?$$

Elementary Braid Matrices



$$\sigma_1 = R = \begin{pmatrix} e^{-i4\pi/5} & 0 \\ 0 & e^{i3\pi/5} \end{pmatrix}$$



$$\sigma_2 = F R F = \begin{pmatrix} -\tau e^{-i\pi/5} & \sqrt{\tau} e^{-i3\pi/5} \\ \sqrt{\tau} e^{-i3\pi/5} & -\tau \end{pmatrix}$$

$$\tau = (\sqrt{5} - 1) / 2 = \varphi^{-1}$$

Elementary Braid Matrices

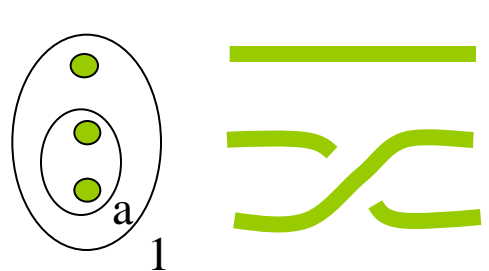


Diagram illustrating the elementary braid matrix σ_1 . On the left, a genus-1 torus with three marked points (green dots) is shown, with the bottom two points enclosed in a circle labeled 'a'. To the right, three horizontal strands are shown: the top strand is straight, the middle and bottom strands cross each other, and the bottom strand is straight.

$$\sigma_1 = \begin{pmatrix} e^{-i4\pi/5} & 0 \\ 0 & e^{i3\pi/5} \end{pmatrix}$$




Diagram illustrating the elementary braid matrix σ_2 . On the left, a genus-1 torus with three marked points (green dots) is shown, with the bottom two points enclosed in a circle labeled 'a'. To the right, three horizontal strands are shown: the top and middle strands cross each other, and the bottom strand is straight.

$$\sigma_2 = F \sigma_1 F = \begin{pmatrix} -\tau e^{-i\pi/5} & \sqrt{\tau} e^{-i3\pi/5} \\ \sqrt{\tau} e^{-i3\pi/5} & -\tau \end{pmatrix}$$

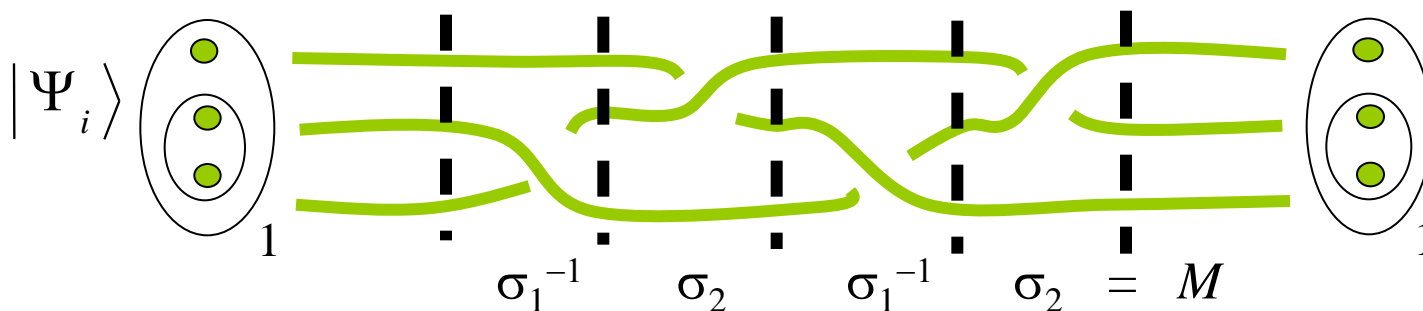
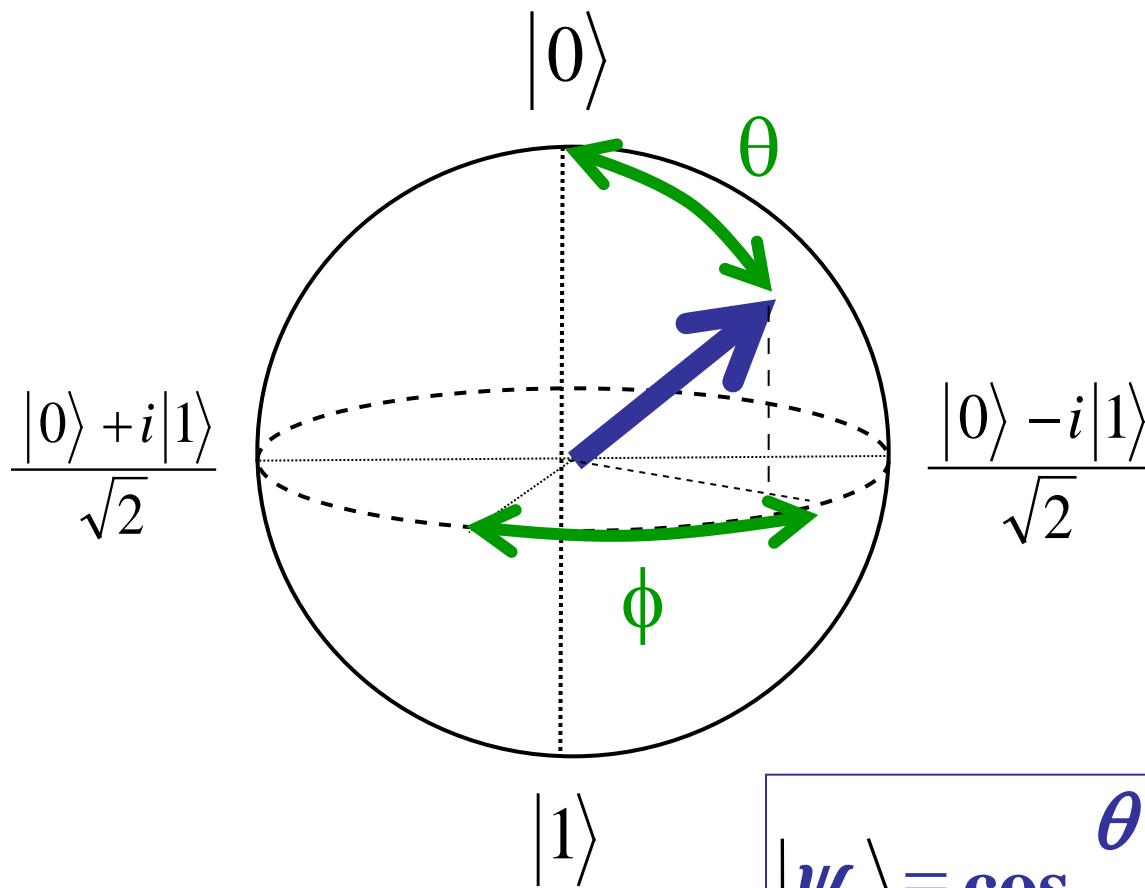


Diagram illustrating the monodromy matrix M . It shows a sequence of braiding operations on three strands. The initial state is a genus-1 torus with three marked points (green dots) and a circle labeled '1' around the bottom two points. The sequence of operations is σ_1^{-1} , σ_2 , σ_1^{-1} , and σ_2 . The final state is a genus-1 torus with three marked points (green dots) and a circle labeled '1' around the bottom two points. The entire sequence is labeled $= M$.

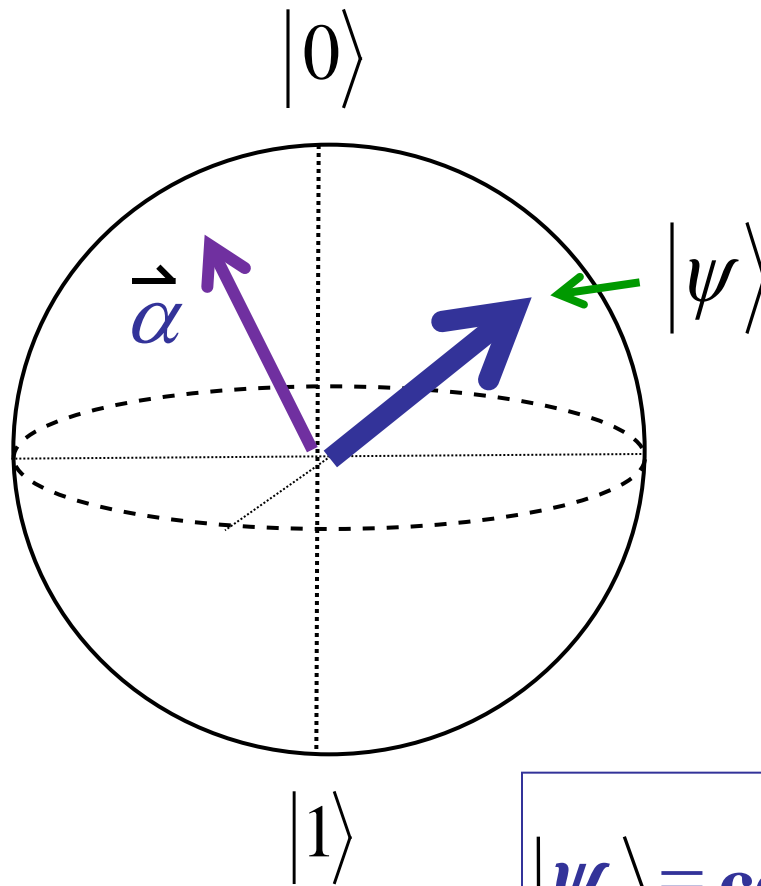
$$|\Psi_i\rangle \xrightarrow{\sigma_1^{-1} \sigma_2 \sigma_1^{-1} \sigma_2} |\Psi_f\rangle = M^T |\Psi_i\rangle$$

A Quantum Bit: A Continuum of States



$$|\psi\rangle = \cos \frac{\theta}{2} |0\rangle + \sin \frac{\theta}{2} e^{-i\phi} |1\rangle$$

Single Qubit Operations: Rotations



$\vec{\alpha}$ = rotation vector

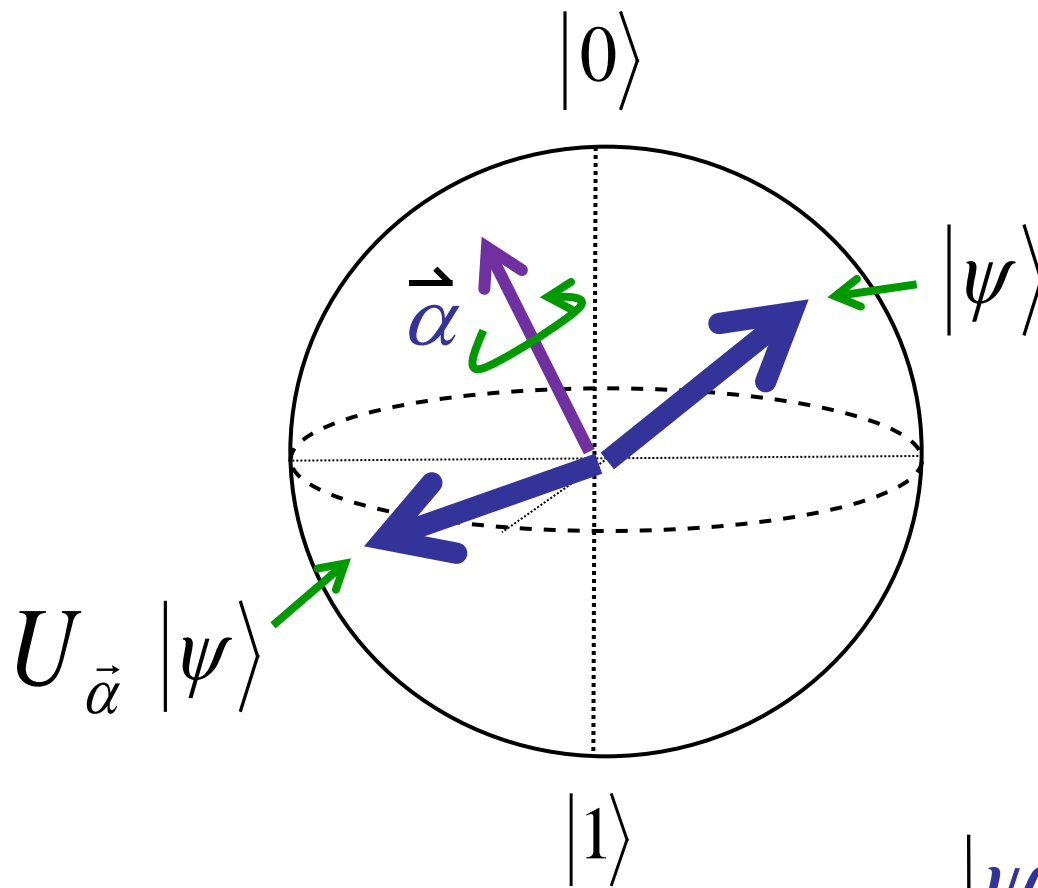
Direction of $\vec{\alpha}$ is
the rotation axis

Magnitude of $\vec{\alpha}$ is
the rotation angle

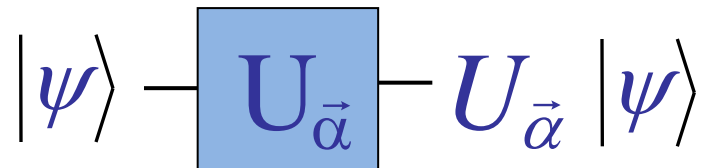
$$|\psi\rangle = \cos \frac{\theta}{2} |0\rangle + \sin \frac{\theta}{2} e^{-i\phi} |1\rangle$$

Single Qubit Operations: Rotations

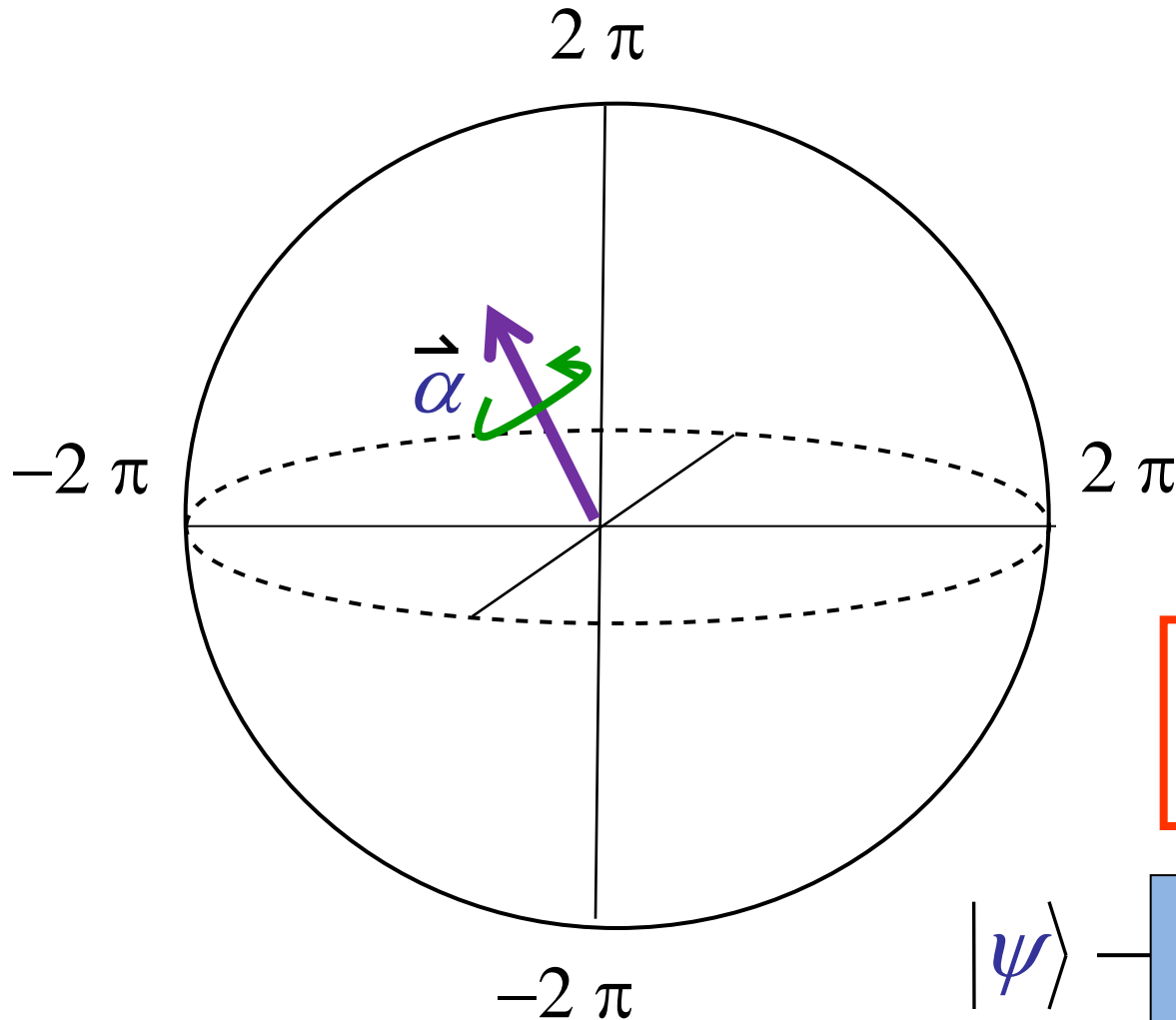
$\vec{\alpha}$ = rotation vector



$$U_{\vec{\alpha}} = \exp \frac{i \vec{\alpha} \cdot \vec{\sigma}}{2}$$



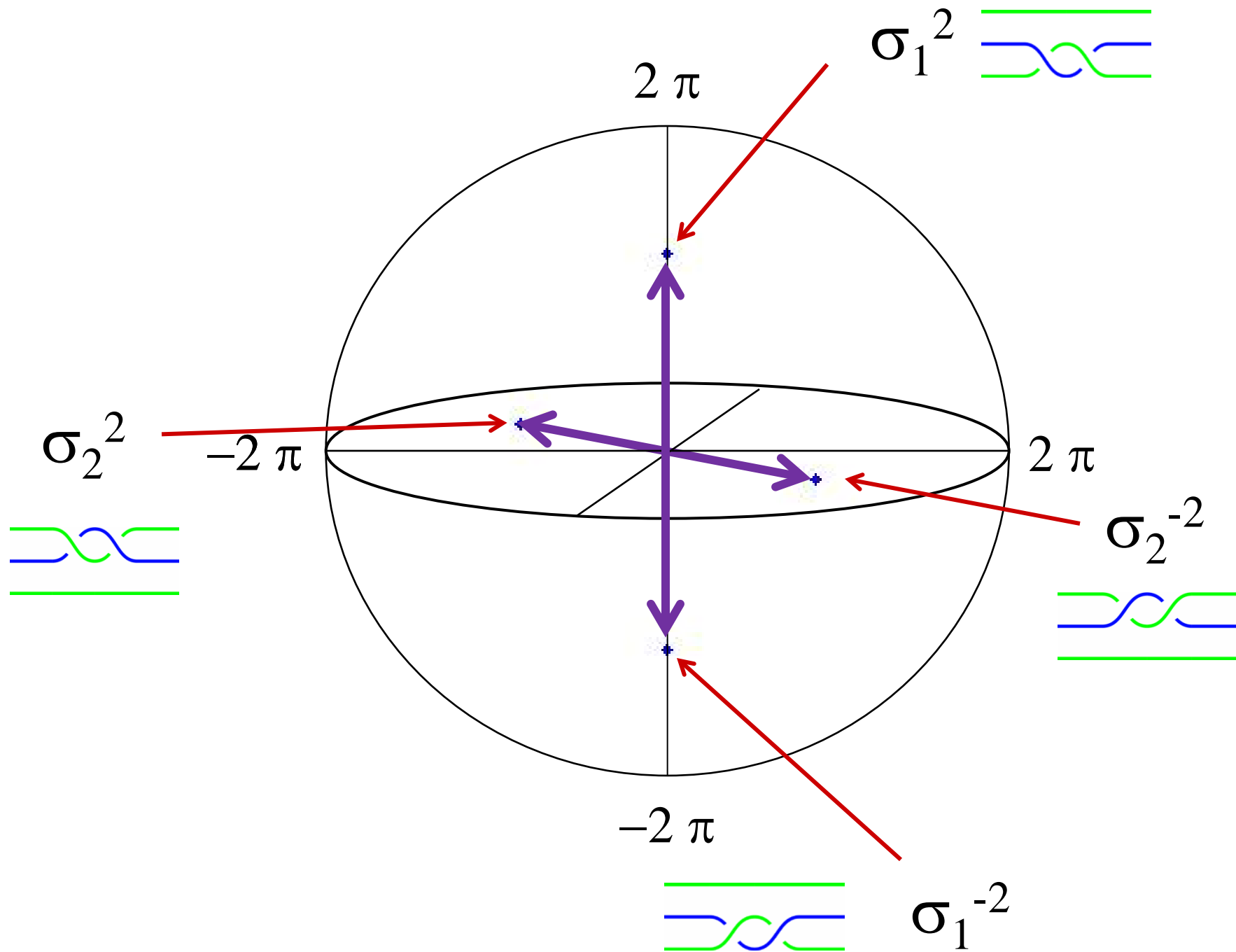
Single Qubit Operations: Rotations



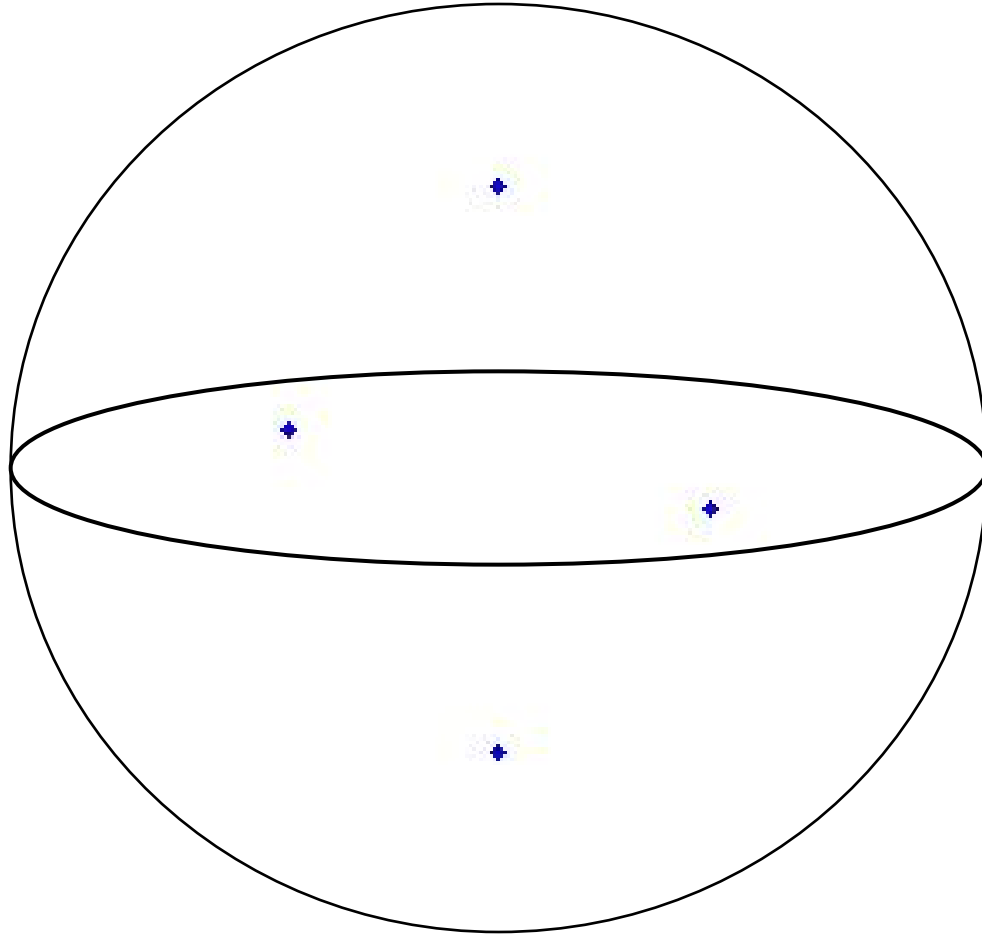
The set of all single qubit rotations lives in a solid sphere of radius 2π .

Qubits are sensitive:
Easy to make errors!

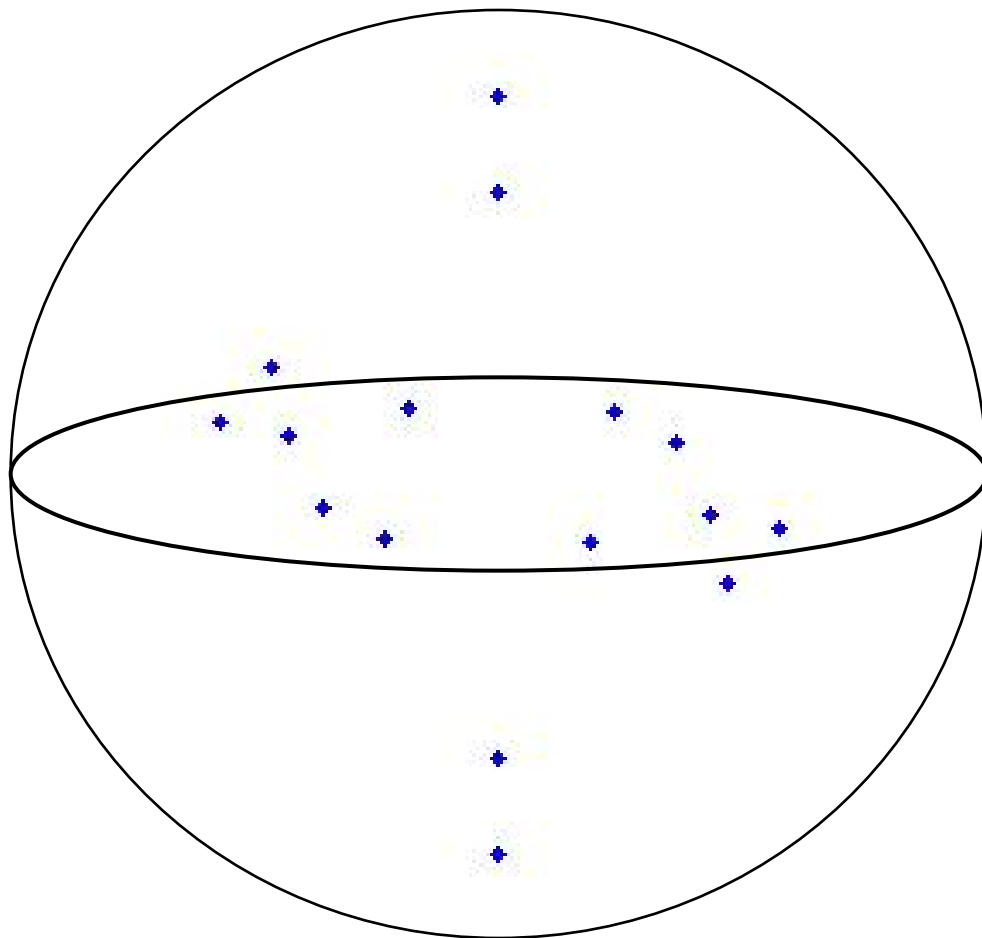
$$|\psi\rangle \rightarrow \boxed{U_{\vec{\alpha}}} U_{\vec{\alpha}} |\psi\rangle$$



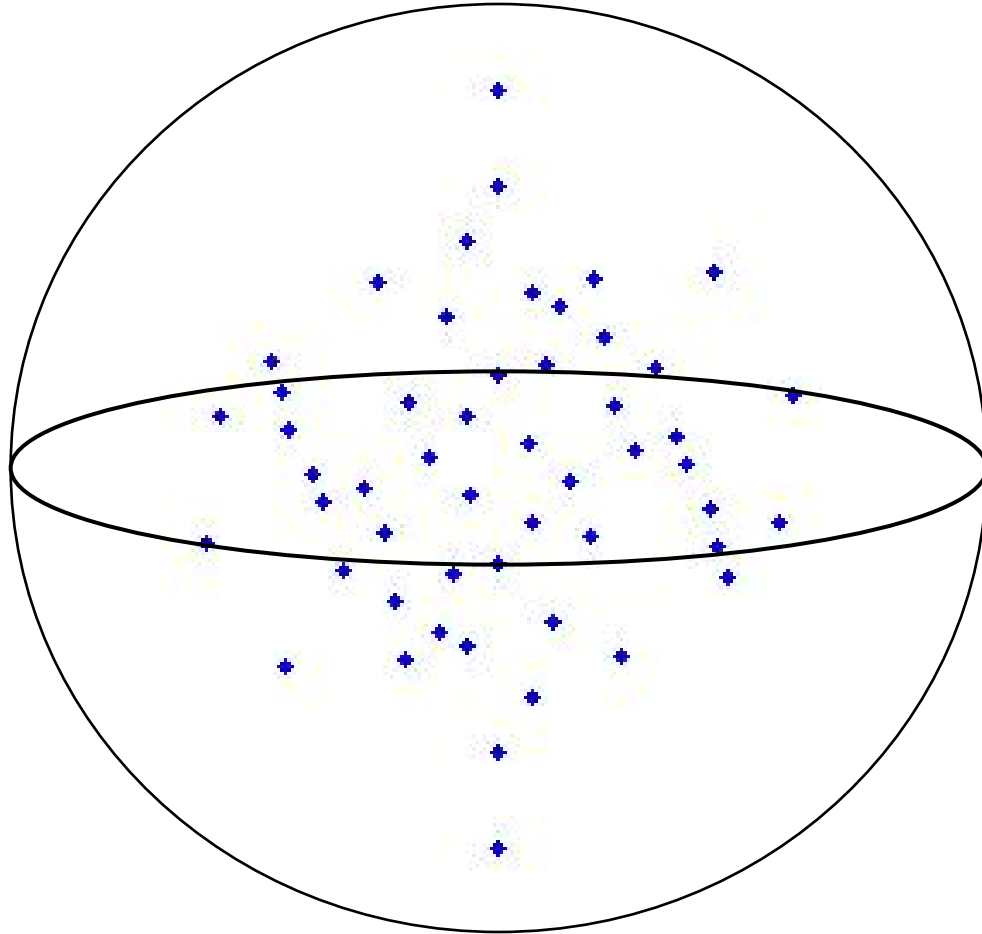
$N = 1$



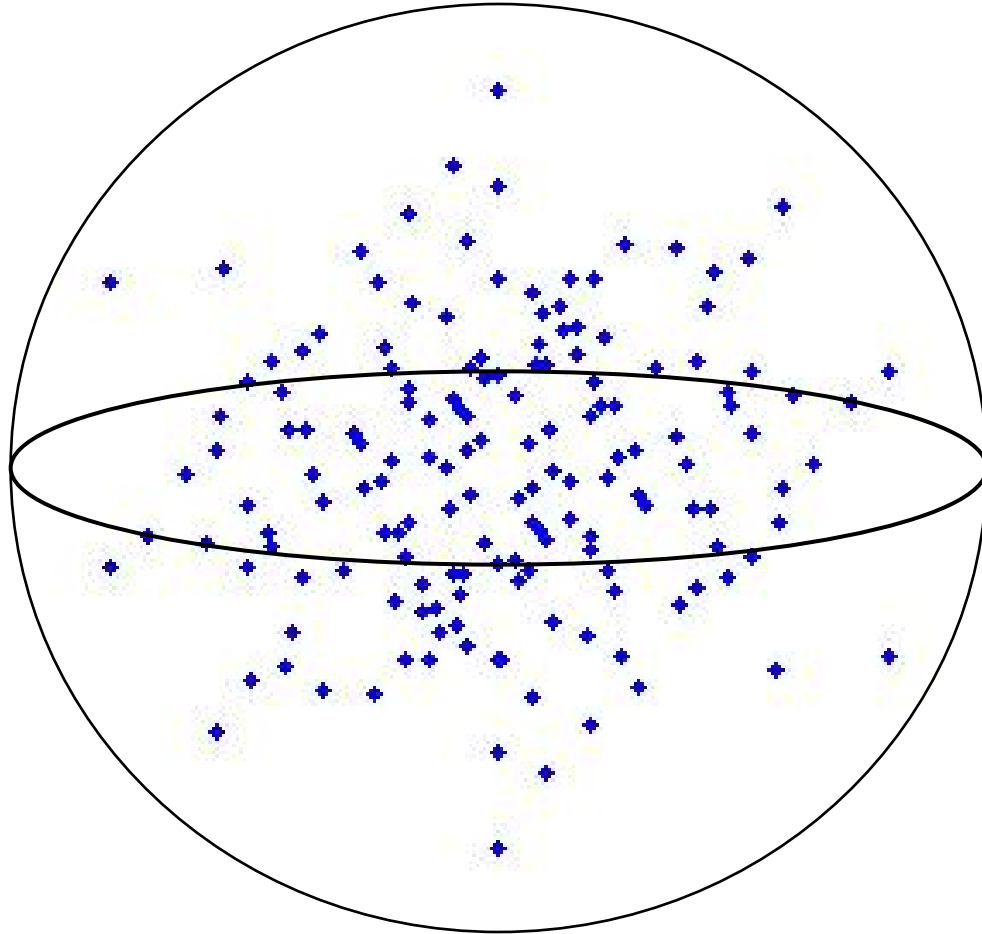
$N = 2$



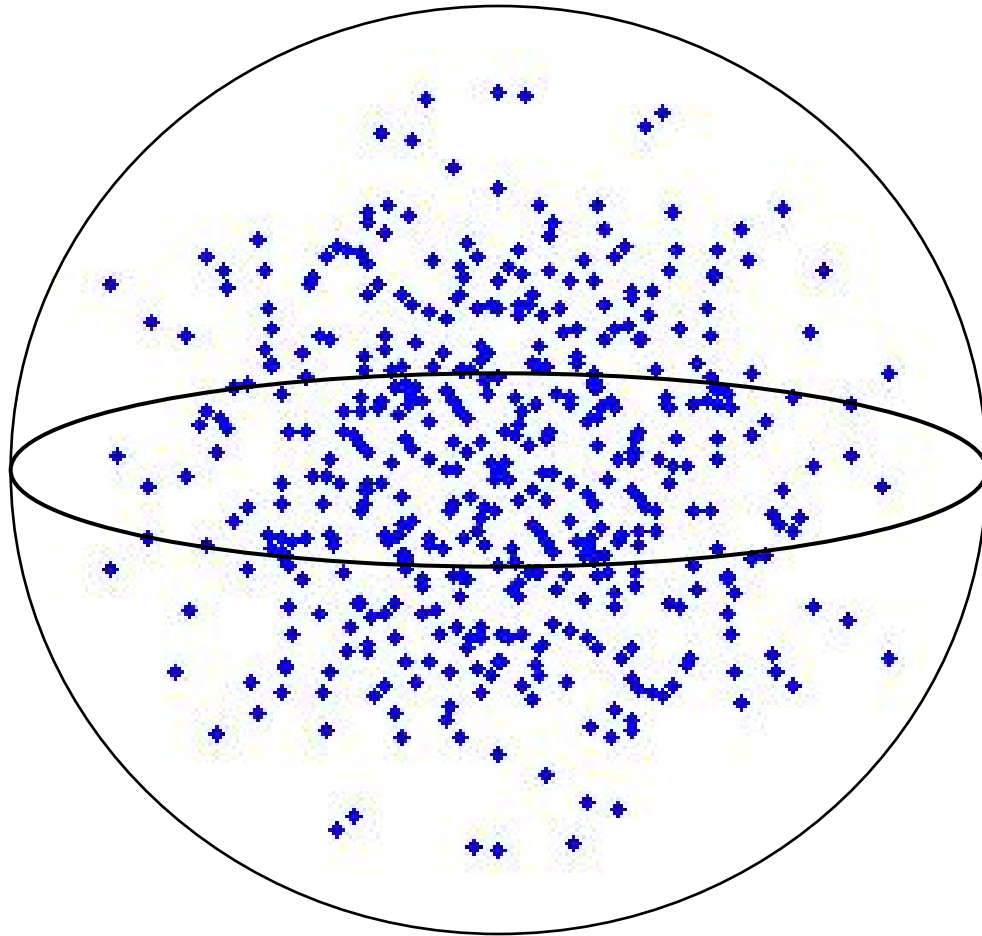
$N = 3$



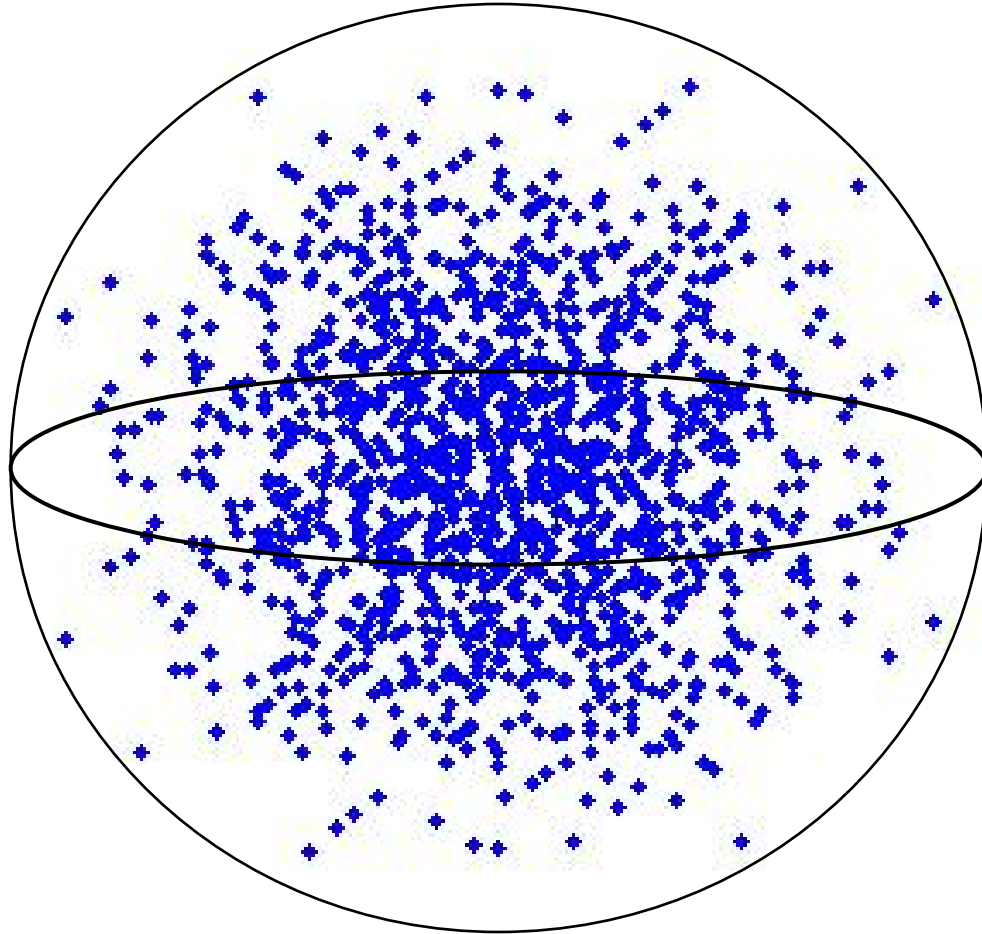
$N = 4$



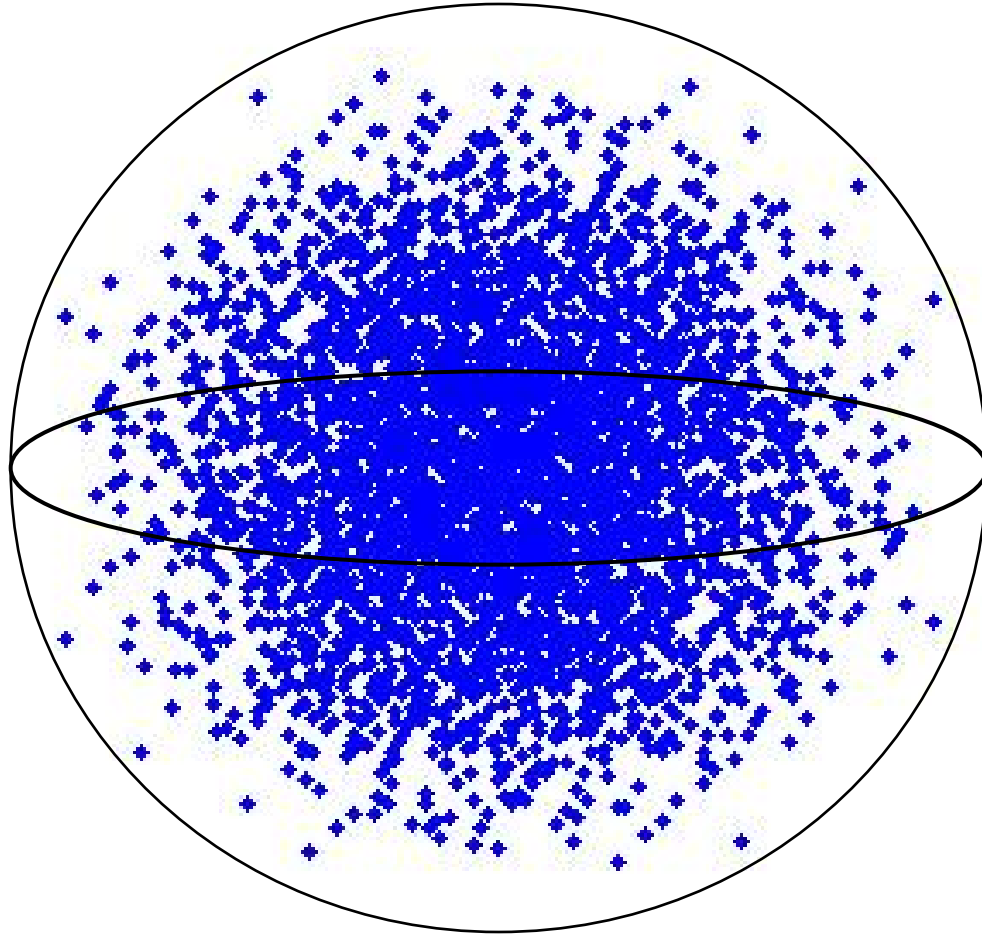
$N = 5$



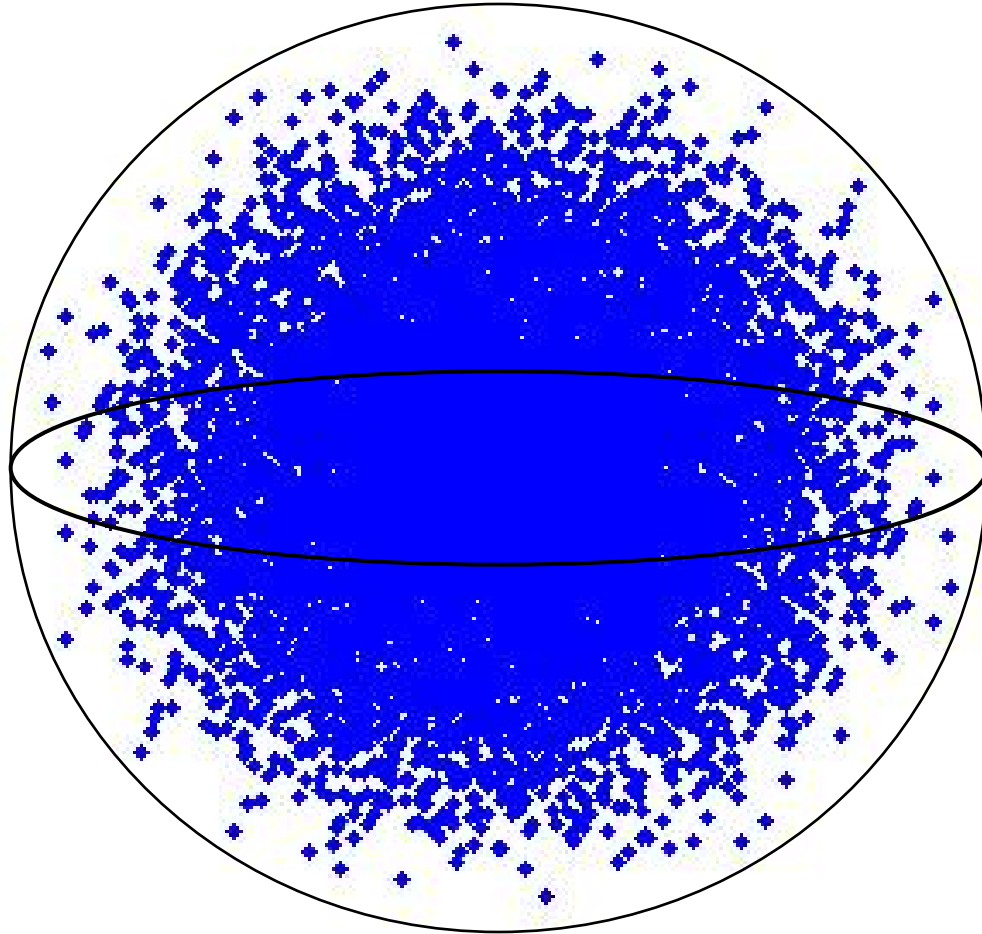
$N = 6$



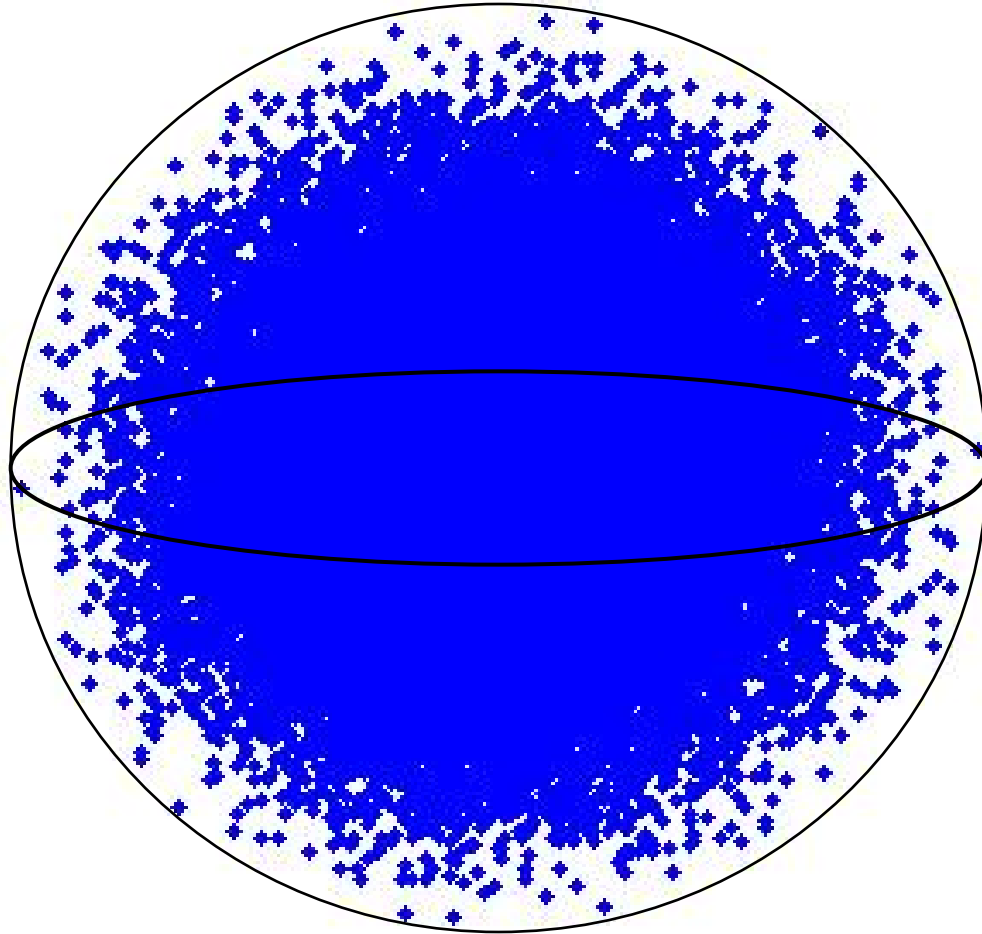
$N = 7$



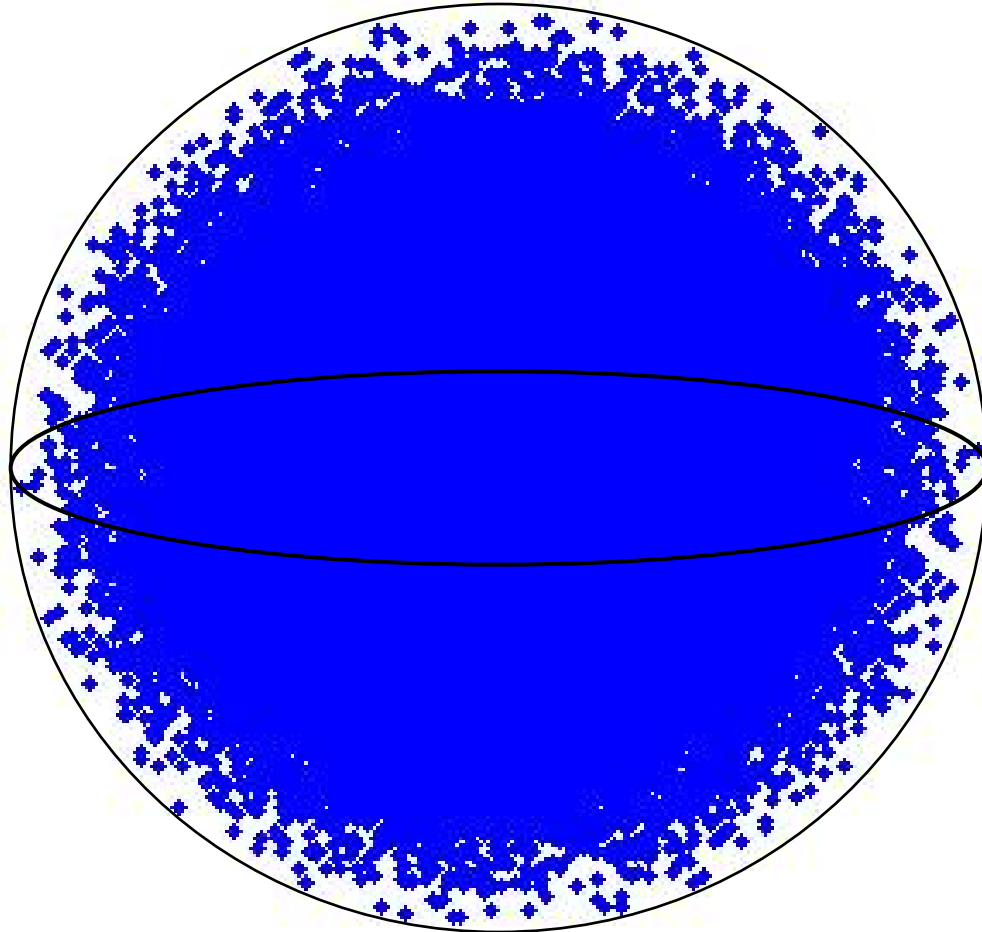
$N = 8$



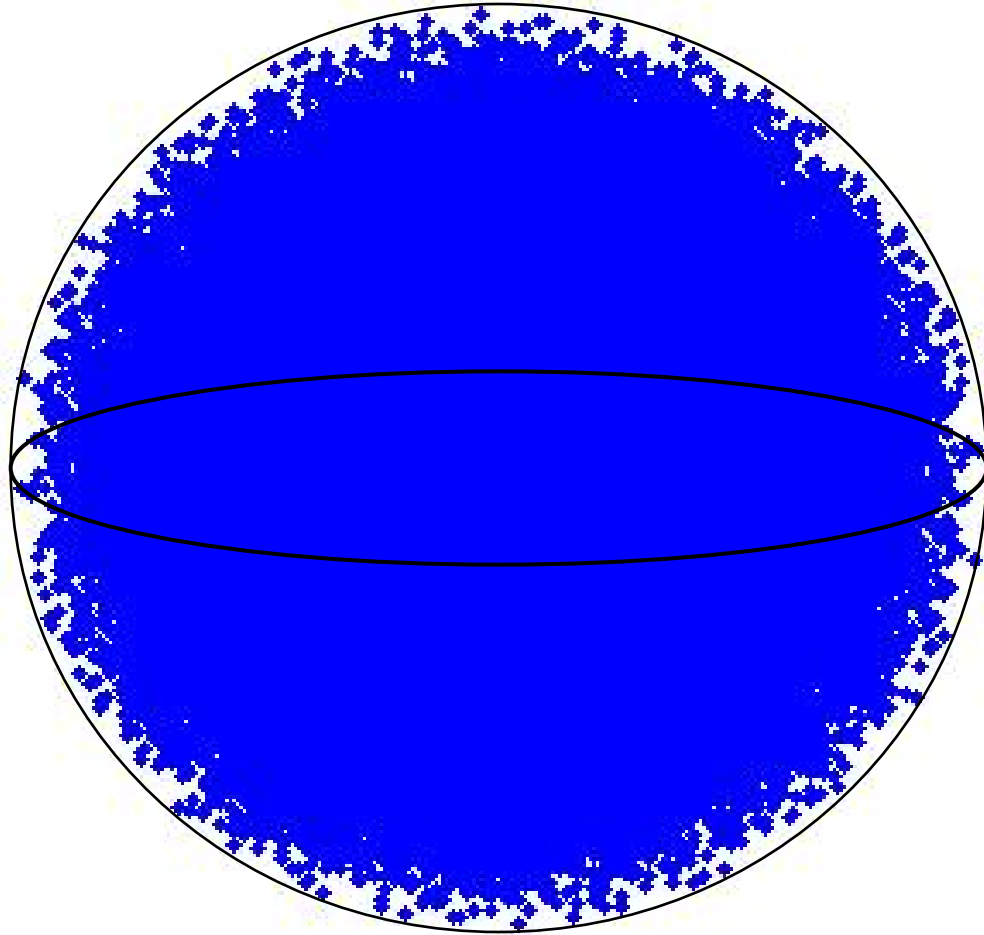
$N = 9$



$N = 10$



$N = 11$



Brute Force Search

$$\sigma_1^{-2}\sigma_2^{-4}\sigma_1^4\sigma_2^{-2}\sigma_1^2\sigma_2^2\sigma_1^{-2}\sigma_2^4\sigma_1^{-2}\sigma_2^4\sigma_1^2\sigma_2^{-4}\sigma_1^2\sigma_2^{-2}\sigma_1^2\sigma_2^{-2}\sigma_1^{-2} = \begin{pmatrix} 0 & i \\ i & 0 \end{pmatrix} + O(10^{-3})$$



Brute Force Search

$$\sigma_1^{-2} \sigma_2^{-4} \sigma_1^4 \sigma_2^{-2} \sigma_1^2 \sigma_2^2 \sigma_1^{-2} \sigma_2^4 \sigma_1^{-2} \sigma_2^4 \sigma_1^2 \sigma_2^{-4} \sigma_1^2 \sigma_2^{-2} \sigma_1^2 \sigma_2^{-2} \sigma_1^{-2} = \begin{pmatrix} 0 & i \\ i & 0 \end{pmatrix} + O(10^{-3})$$

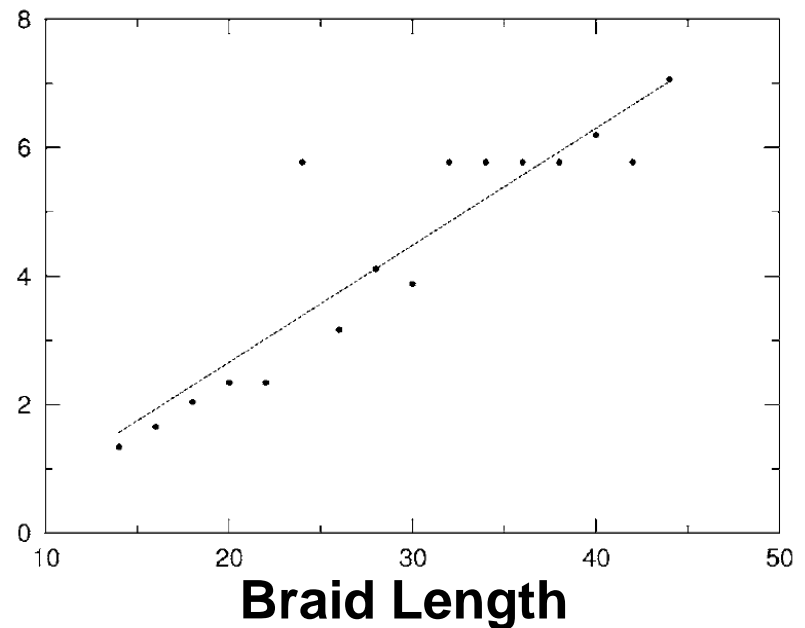
“error” ε



For brute force search:

Braid Length $\sim |\ln \varepsilon|$

$|\ln \varepsilon|$



Brute Force Search

$$\sigma_1^{-2}\sigma_2^{-4}\sigma_1^4\sigma_2^{-2}\sigma_1^2\sigma_2^2\sigma_1^{-2}\sigma_2^4\sigma_1^{-2}\sigma_2^4\sigma_1^2\sigma_2^{-4}\sigma_1^2\sigma_2^{-2}\sigma_1^2\sigma_2^{-2}\sigma_1^{-2} = \begin{pmatrix} 0 & i \\ i & 0 \end{pmatrix} + O(10^{-3})$$

“error” ε



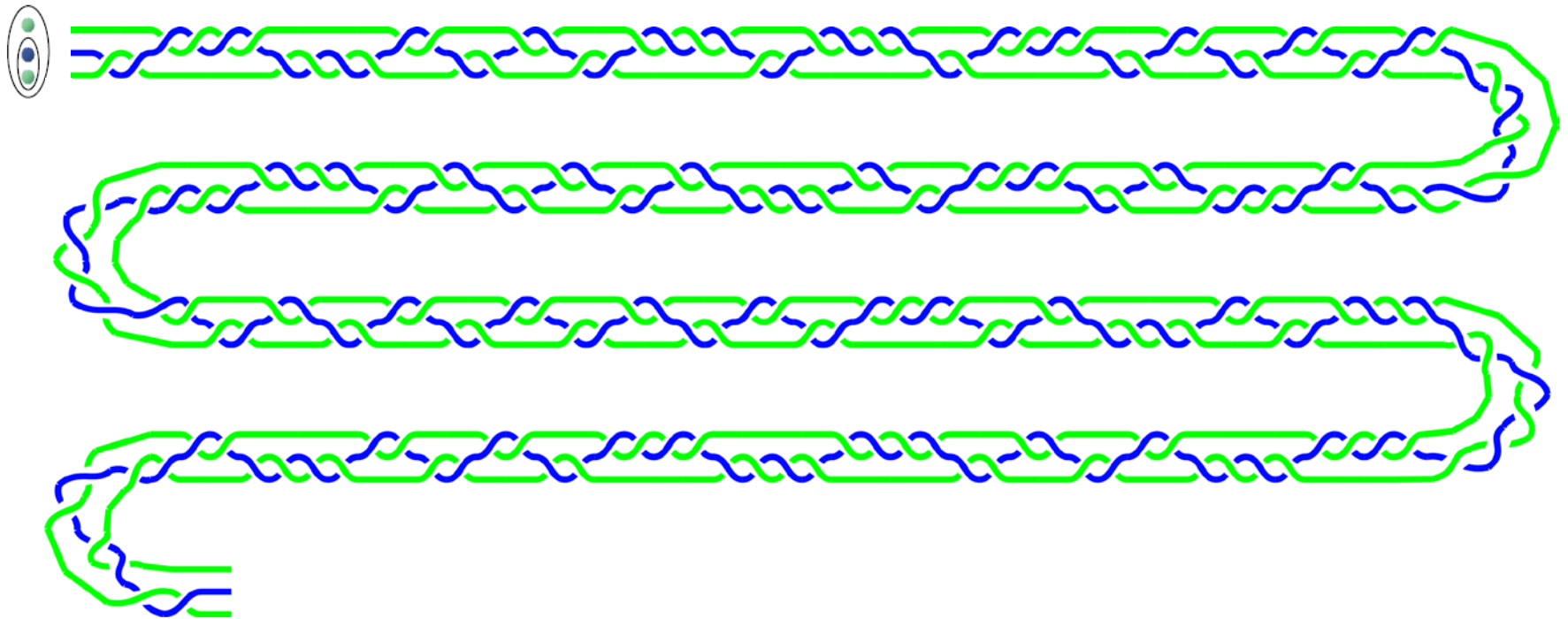
Brute force searching rapidly becomes infeasible as braids get longer.

Fortunately, a clever algorithm due to [Solovay and Kitaev](#) allows for systematic improvement of the braid given a sufficiently dense covering of $SU(2)$.

Solovay-Kitaev Construction

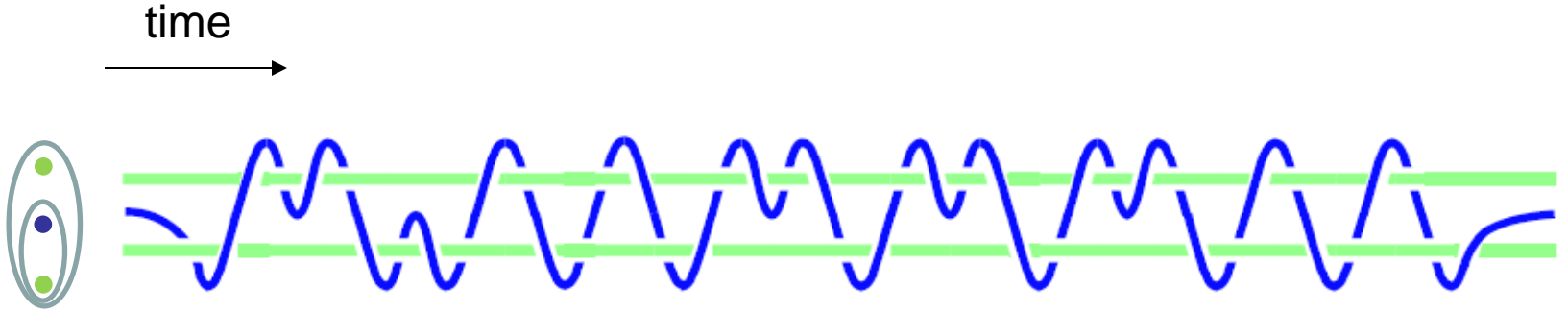
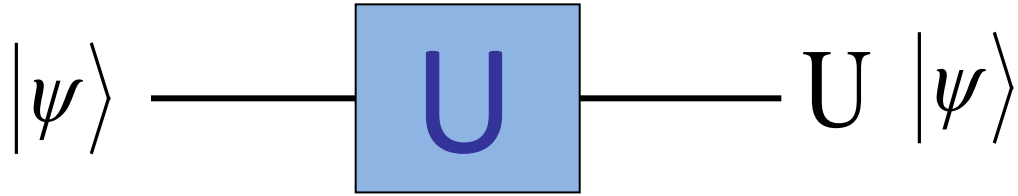
$$\begin{pmatrix} 0 & i \\ i & 0 \end{pmatrix} + O(10^{-5})$$

↑
“error” ε

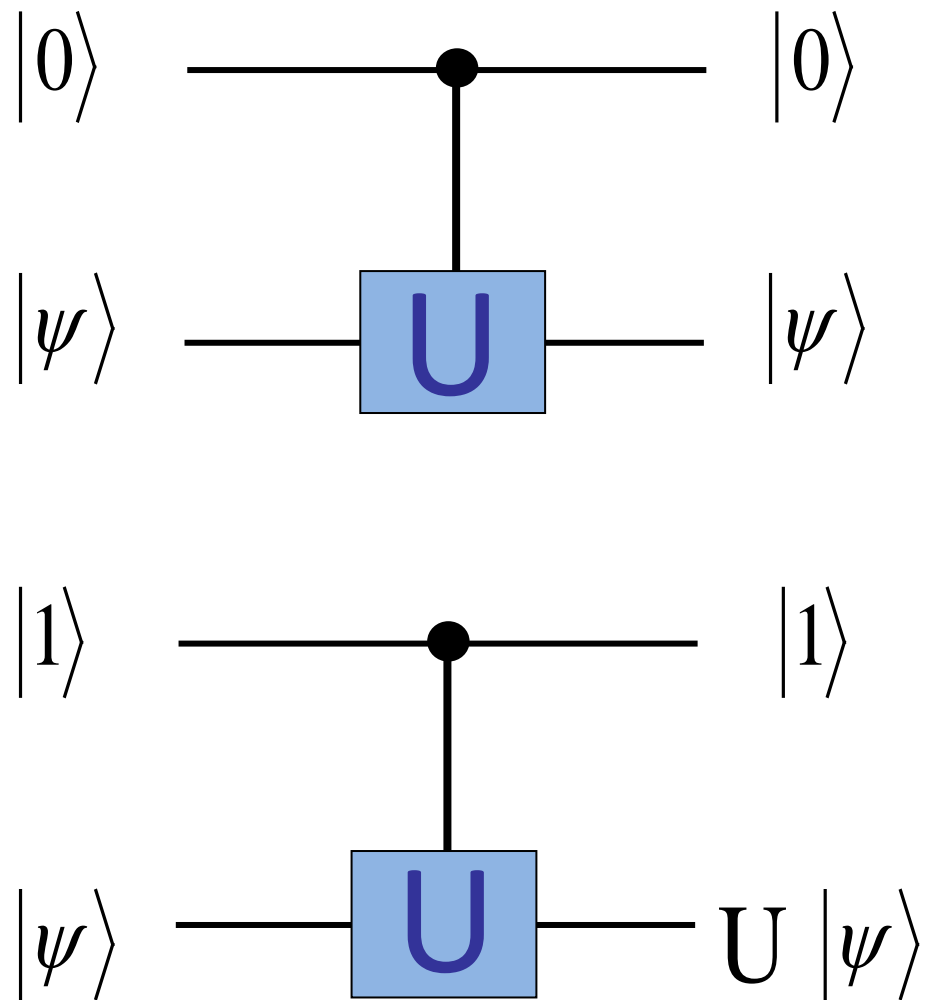


$$\text{Braid Length} \sim |\ln \varepsilon|^c, \quad c \approx 4$$

Single Qubit Operations

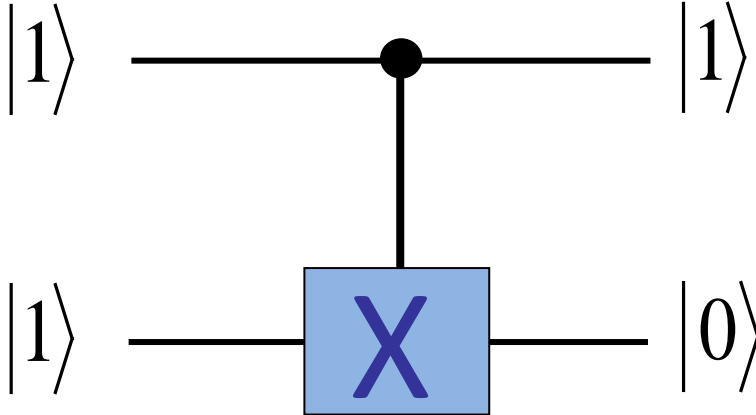
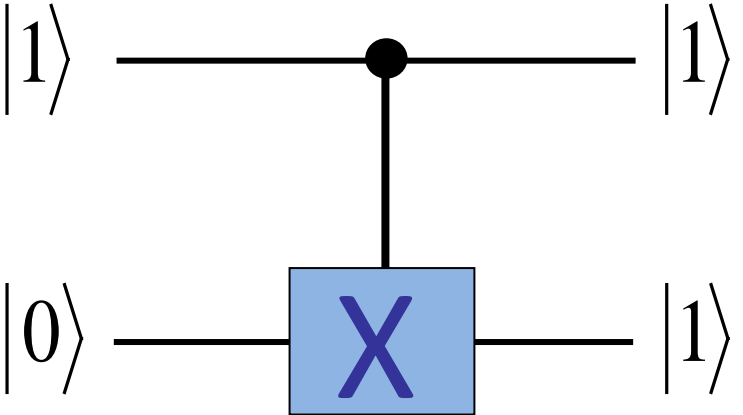
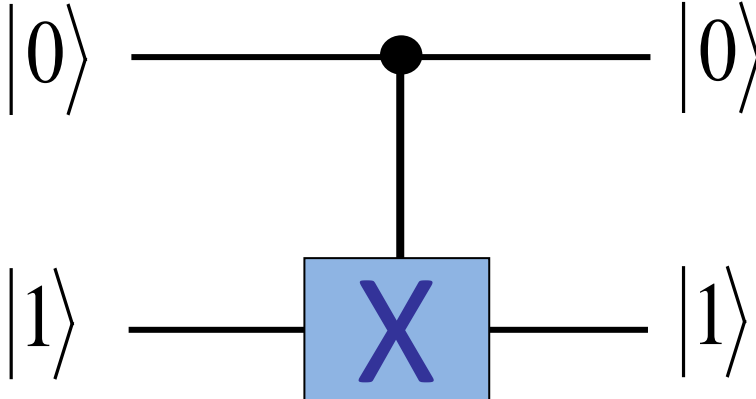
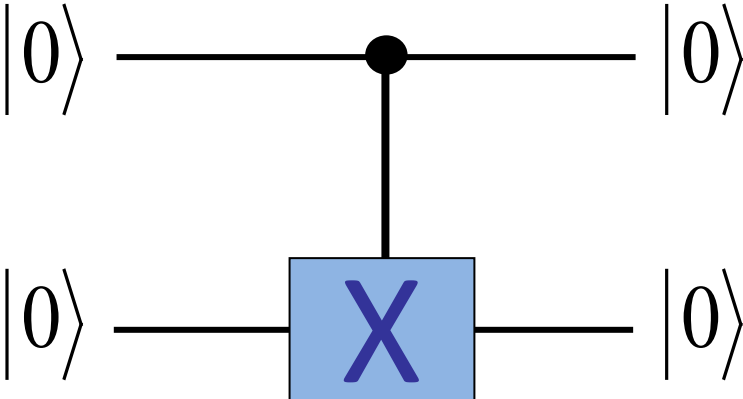


Two Qubit Gates

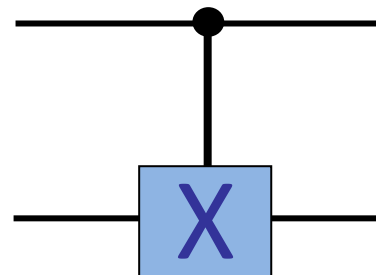


Controlled-NOT Gate

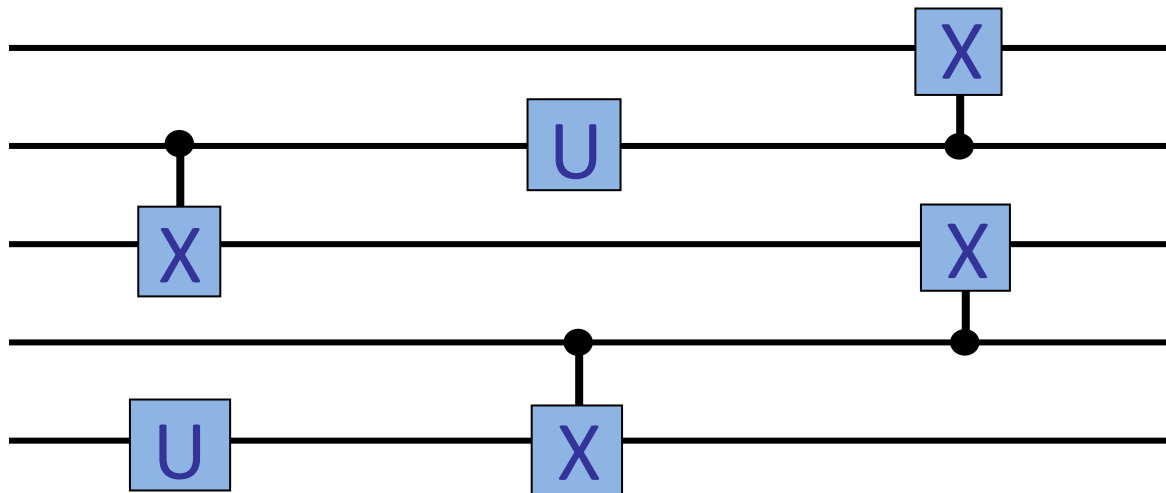
$$X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$



Universal Set of Gates



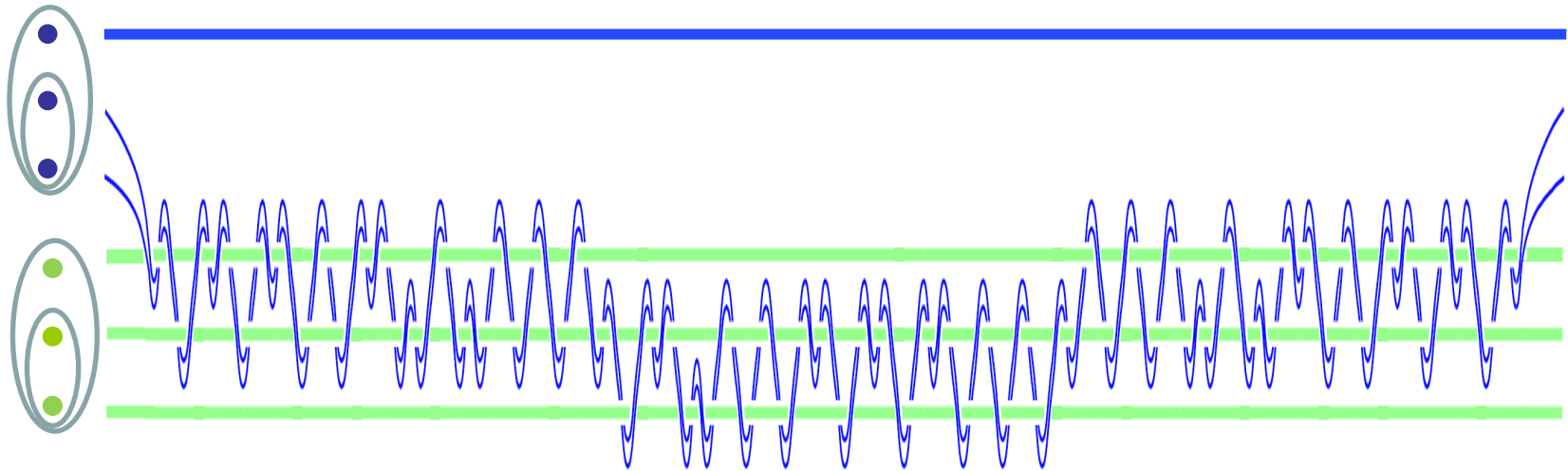
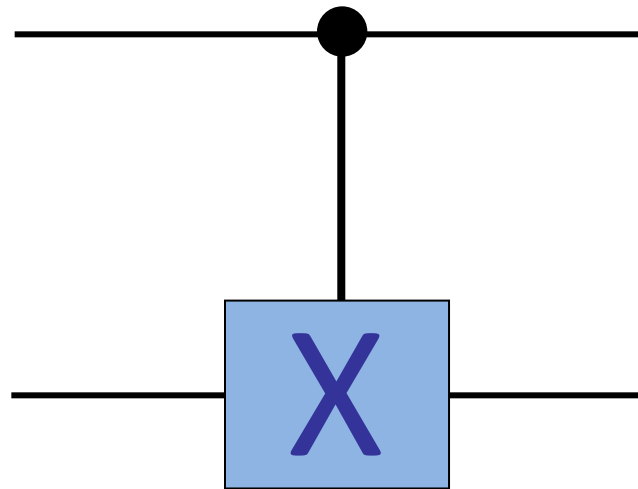
Quantum Circuit



Any unitary operation can be expressed as a quantum circuit

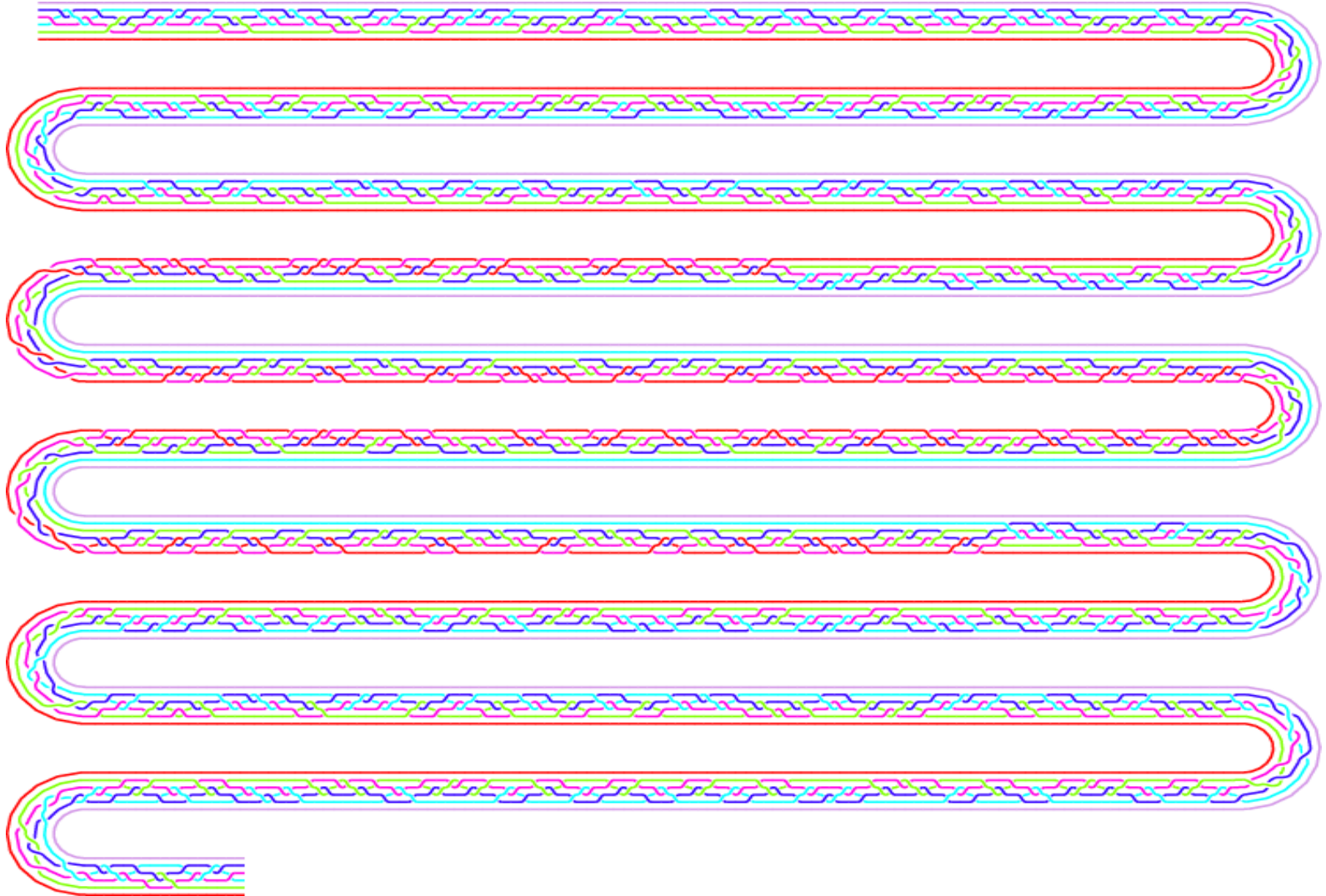
Two Qubit Gates (CNOT)

$$X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

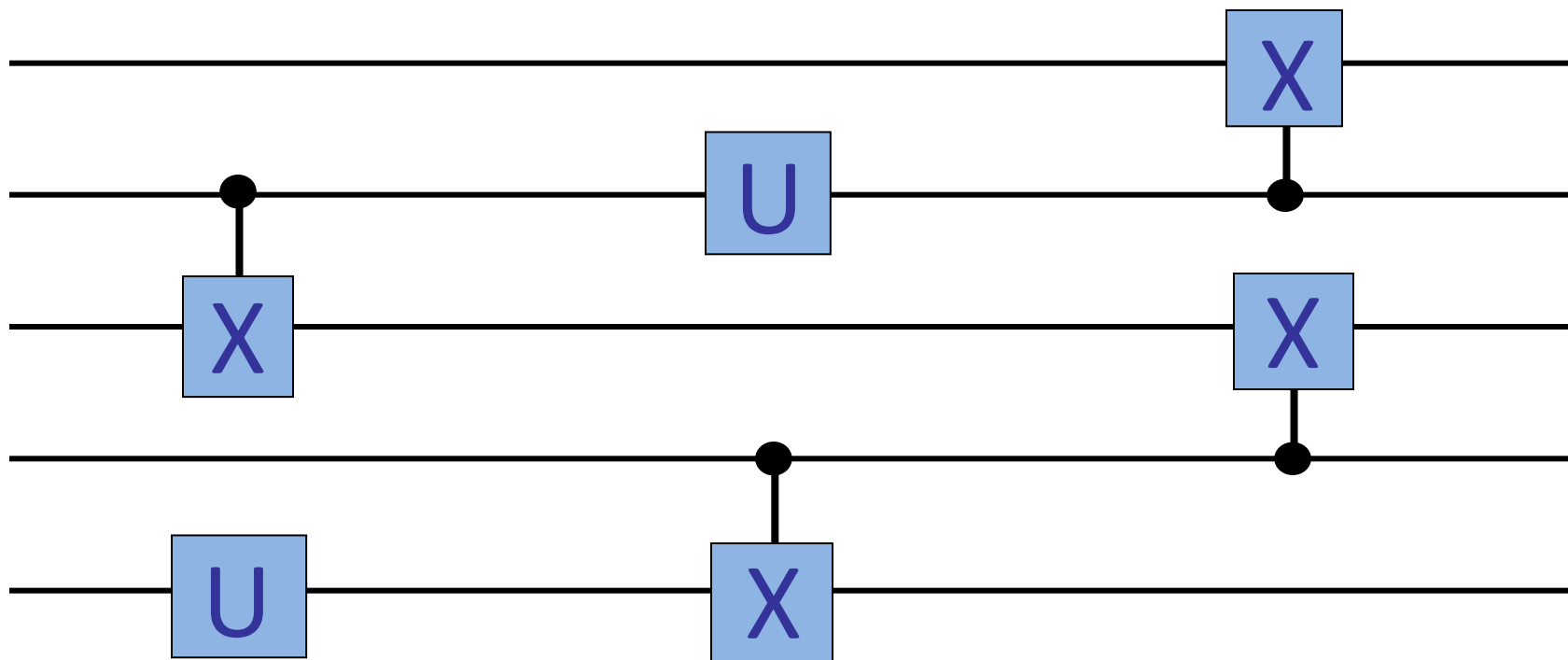


NEB, L. Hormozi, G. Zikos, & S. Simon, PRL 2005
L. Hormozi, G. Zikos, NEB, and S. H. Simon, PRB 2007

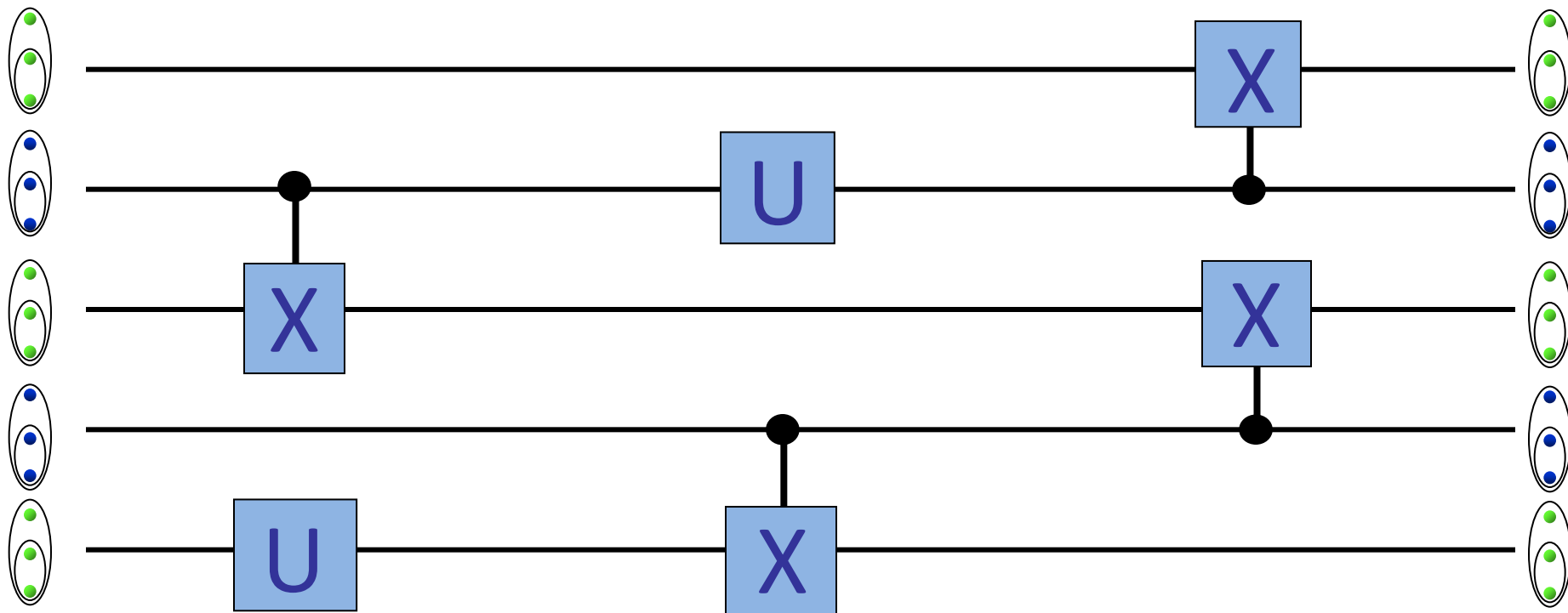
SK Improved Controlled-NOT Gate



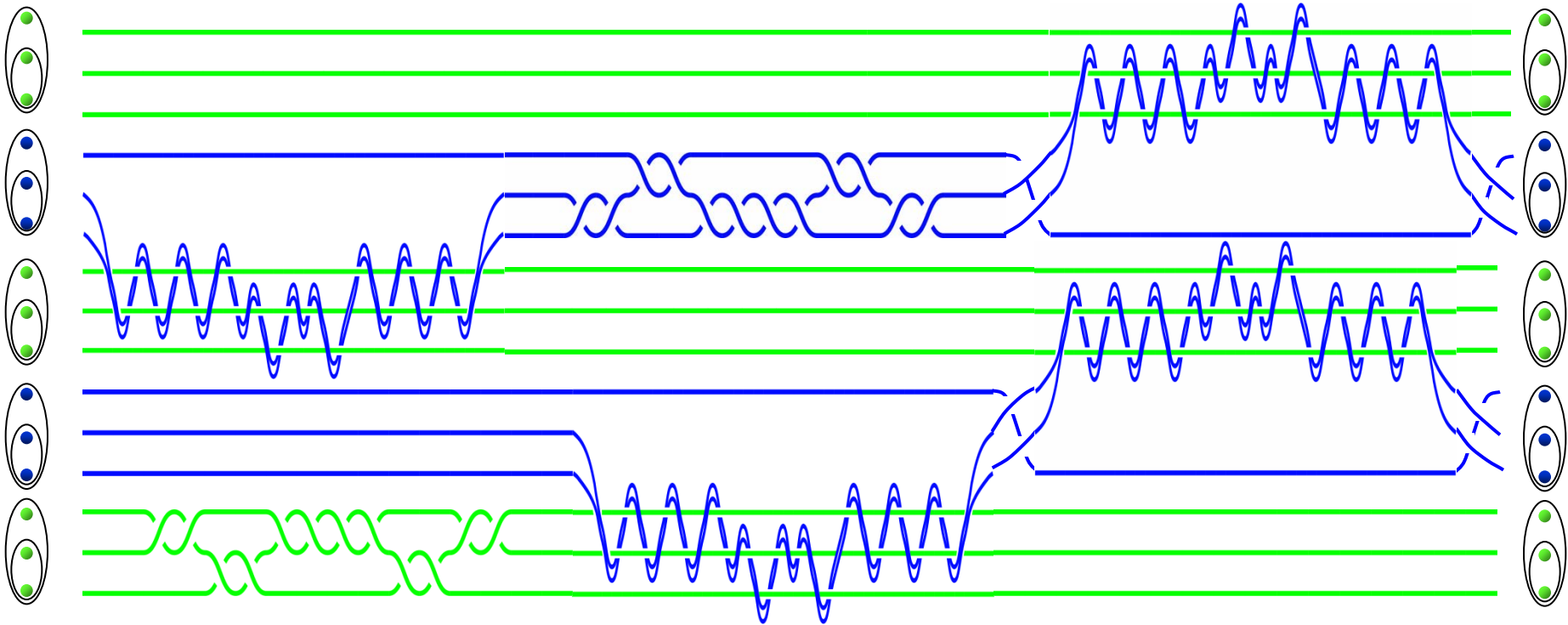
Quantum Circuit



Quantum Circuit



Braid



$SU(2)_k$ Nonabelian Particles

Quasiparticle excitations of the $\nu=12/5$ state (if it is a Read-Rezayi state) are described (up to abelian phases) by $SU(2)_3$ Chern-Simons theory.

In fact, Read and Rezayi proposed an infinite sequence of nonabelian states labeled by an index k with quasiparticle excitations described by $SU(2)_k$ Chern-Simons theory.

Read and Rezayi, 1999

Slingerland and Bais, 2001

Universal for quantum computation for $k=3$ and $k > 4$.

Freedman, Larsen, and Wang, 2001

But before $SU(2)_k$, there was just plain old $SU(2)$

Particles with Ordinary Spin: SU(2)

1. Particles have spin $s = 0, 1/2, 1, 3/2, \dots$

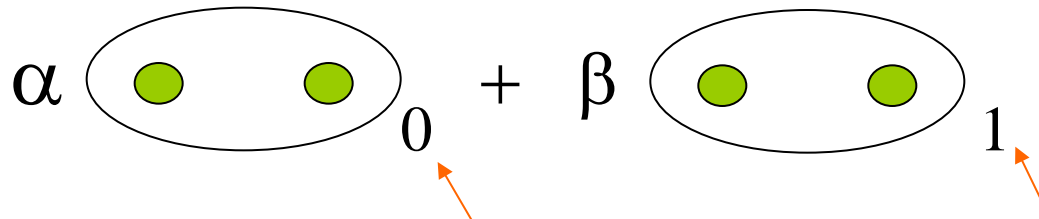


2. “Triangle Rule” for adding angular momentum:

$$s_1 \otimes s_2 = |s_1 - s_2| \oplus (|s_1 - s_2| + 1) \oplus \dots \oplus s_1 + s_2$$

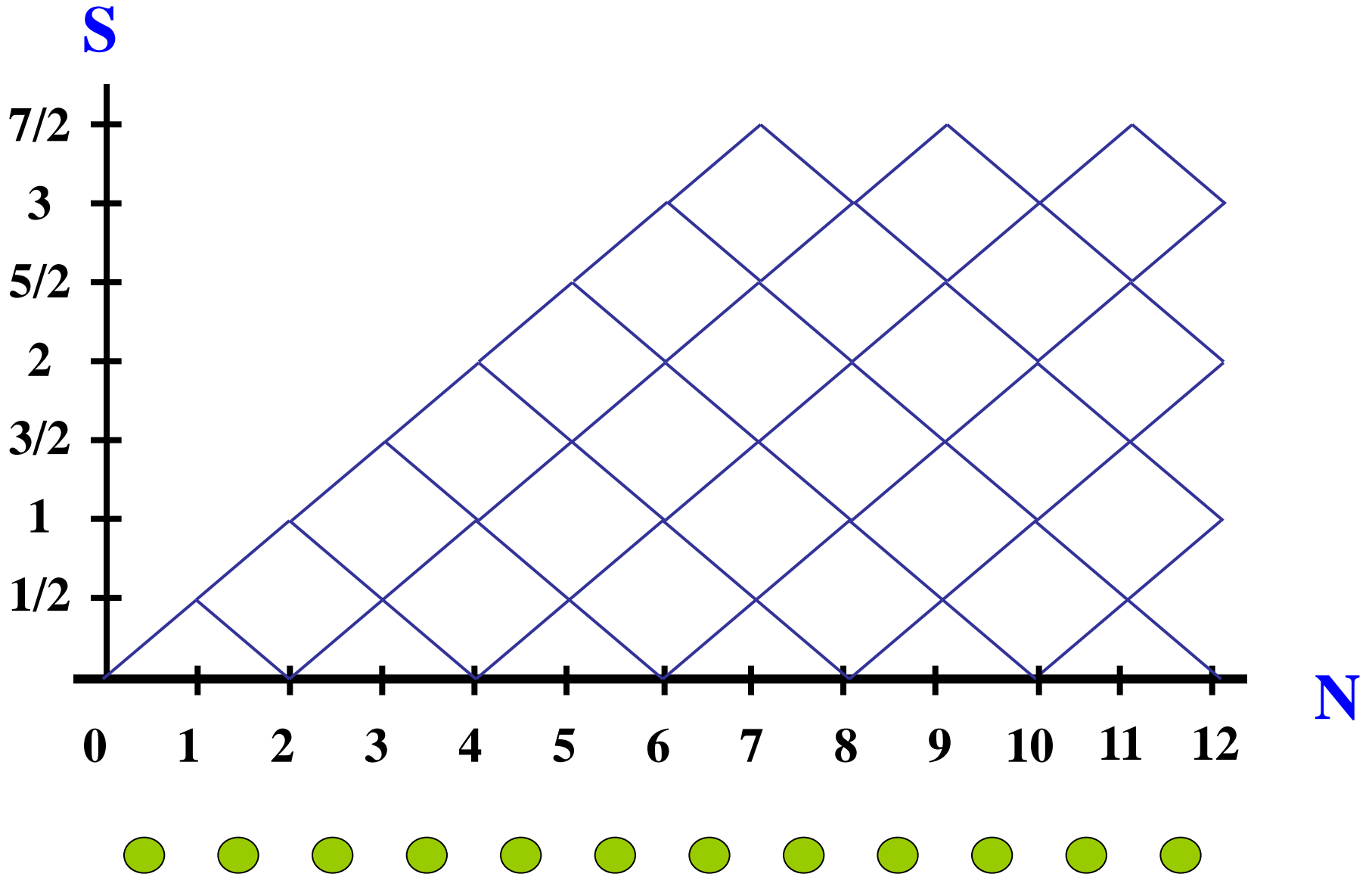
For example: $\frac{1}{2} \otimes \frac{1}{2} = 0 \oplus 1$

→ Two ● particles can have total spin **0** or **1**.

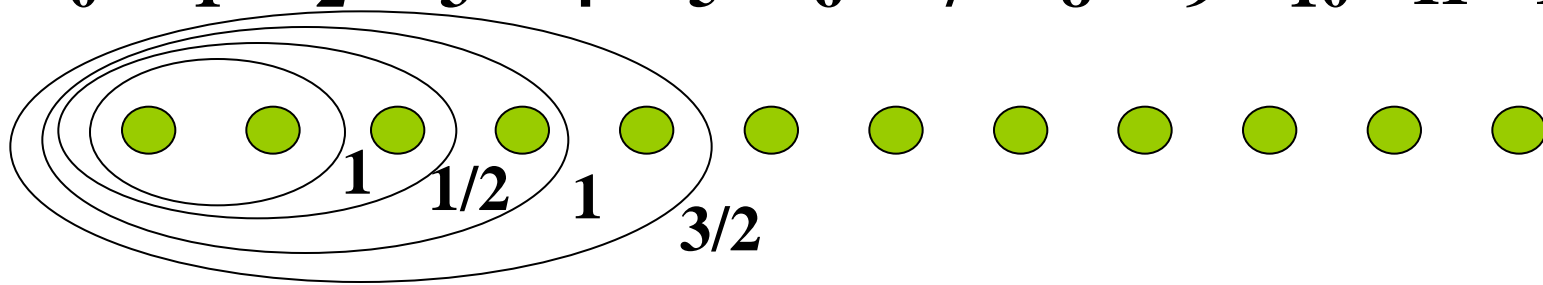
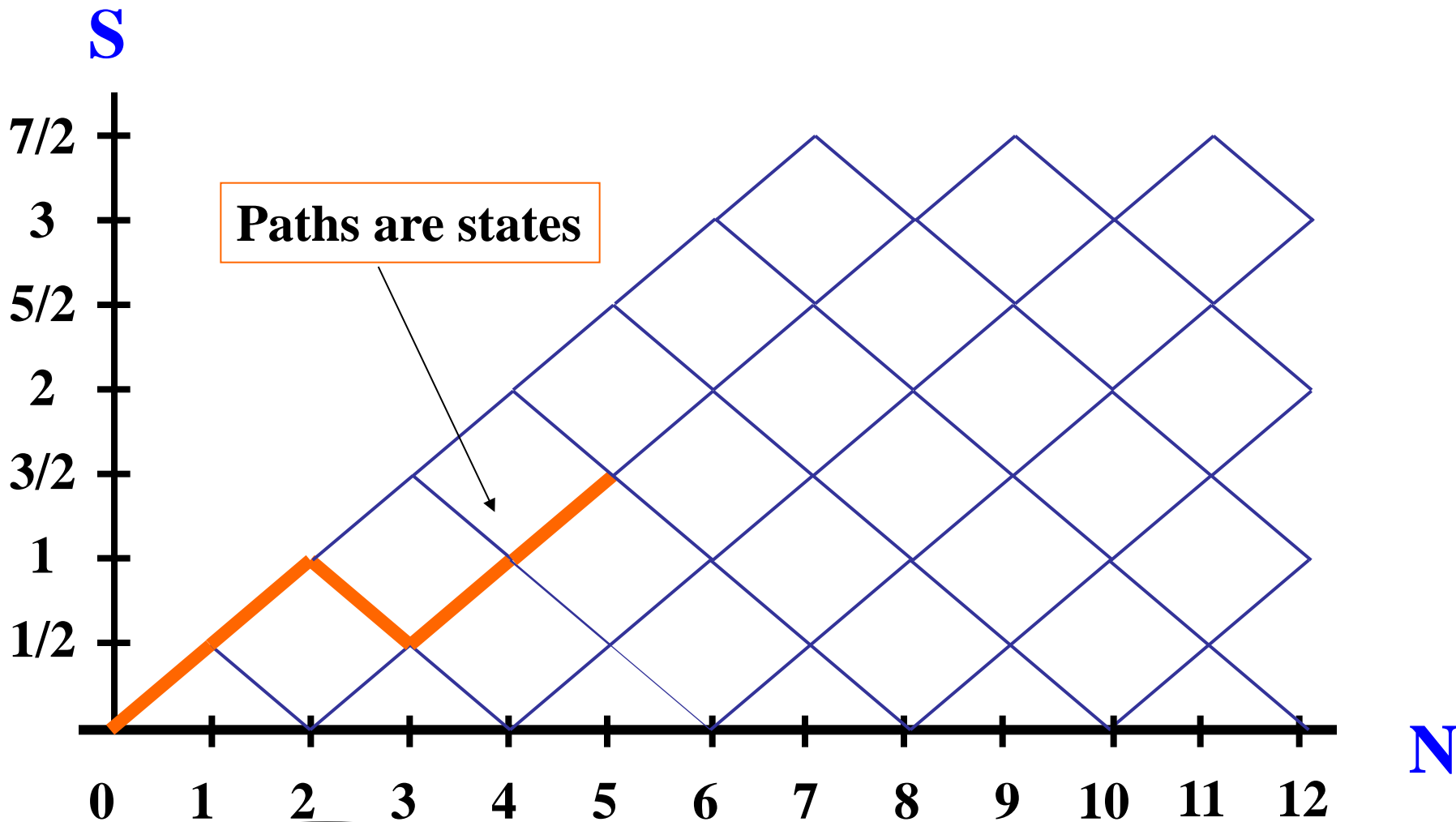


Numbers label total spin of particles inside ovals

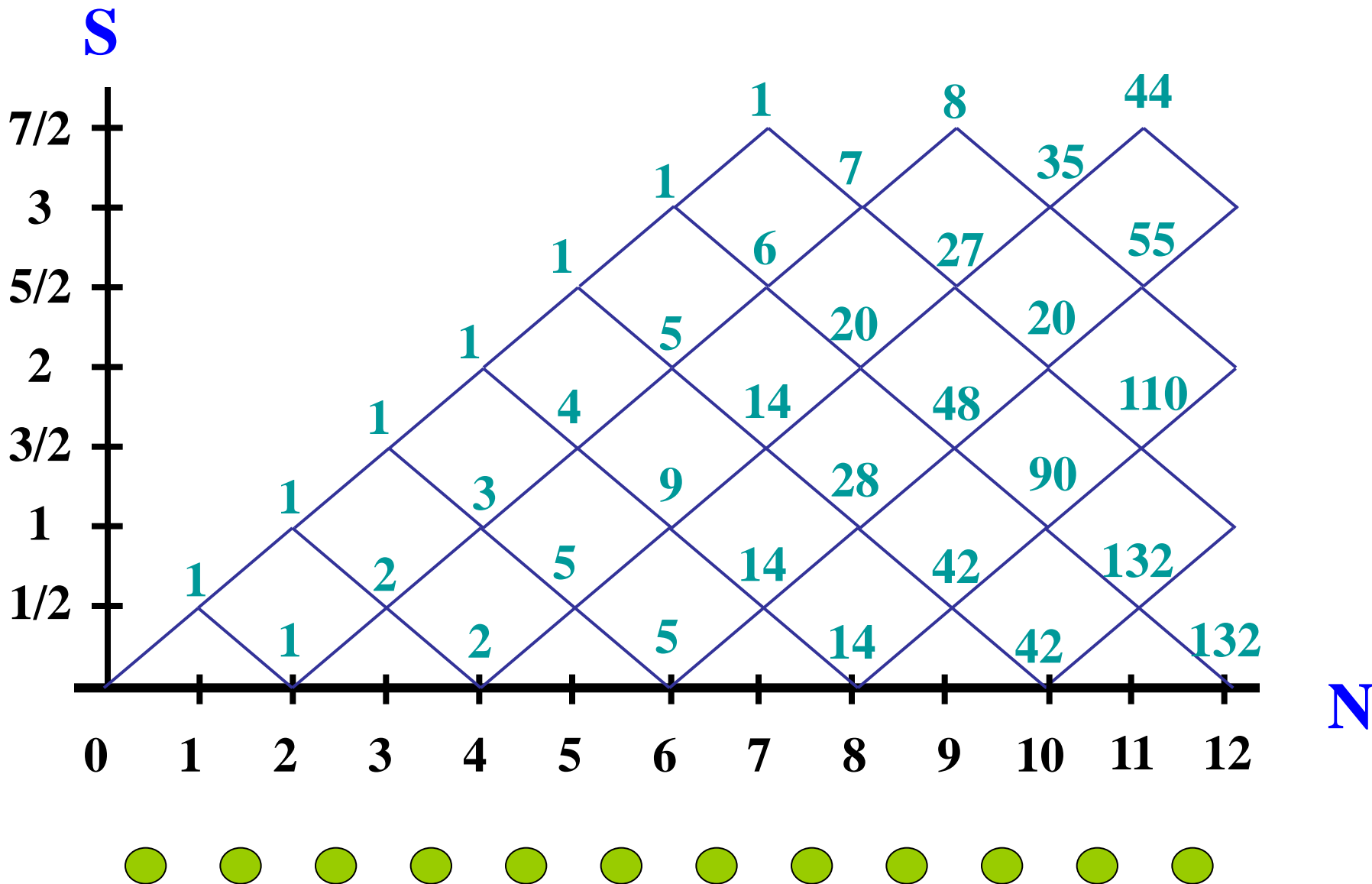
Hilbert Space



Hilbert Space



Hilbert Space



Particles with Ordinary Spin: SU(2)

1. Particles have spin $s = 0, 1/2, 1, 3/2, \dots$

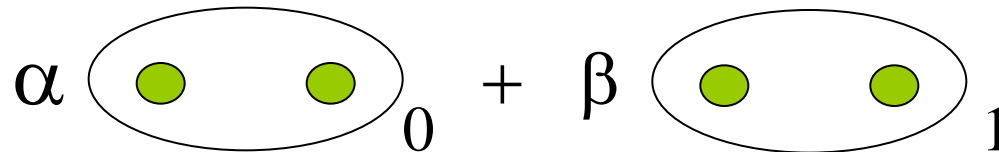


2. “Triangle Rule” for adding angular momentum:

$$s_1 \otimes s_2 = |s_1 - s_2| \oplus (|s_1 - s_2| + 1) \oplus \dots \oplus s_1 + s_2$$

For example: $\frac{1}{2} \otimes \frac{1}{2} = 0 \oplus 1$

→ Two  particles can have total spin **0** or **1**.



Nonabelian Particles: $SU(2)_k$

1. Particles have topological charge $s = 0, 1/2, 1, 3/2, \dots, k/2$

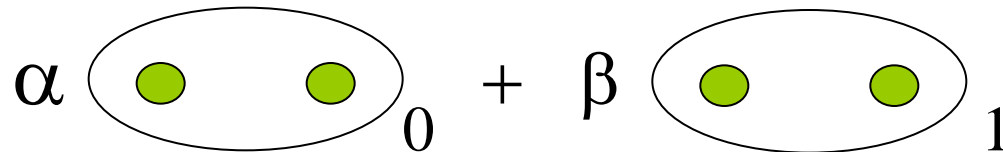


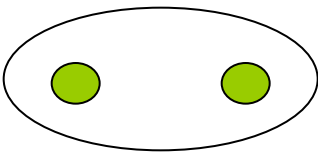
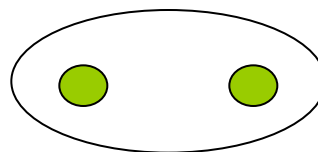
2. “Fusion Rule” for adding topological charge:

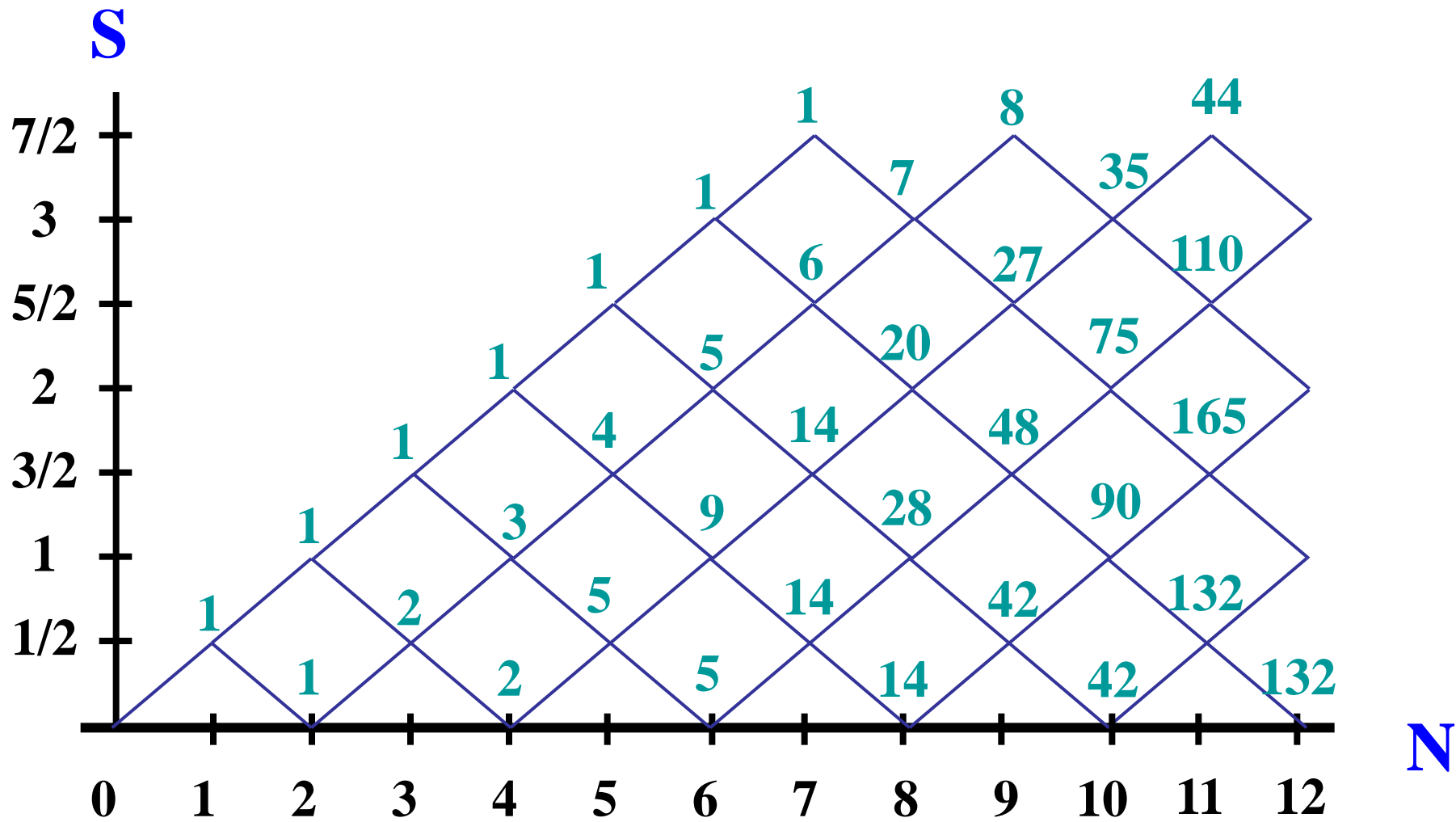
$$s_1 \otimes s_2 = |s_1 - s_2| \oplus (|s_1 - s_2| + 1) \oplus \dots \oplus \min[s_1 + s_2, s_1 + s_2 - k/2]$$

For example: $\frac{1}{2} \otimes \frac{1}{2} = 0 \oplus 1$

→ Two ● particles can have total topological charge **0** or **1**.



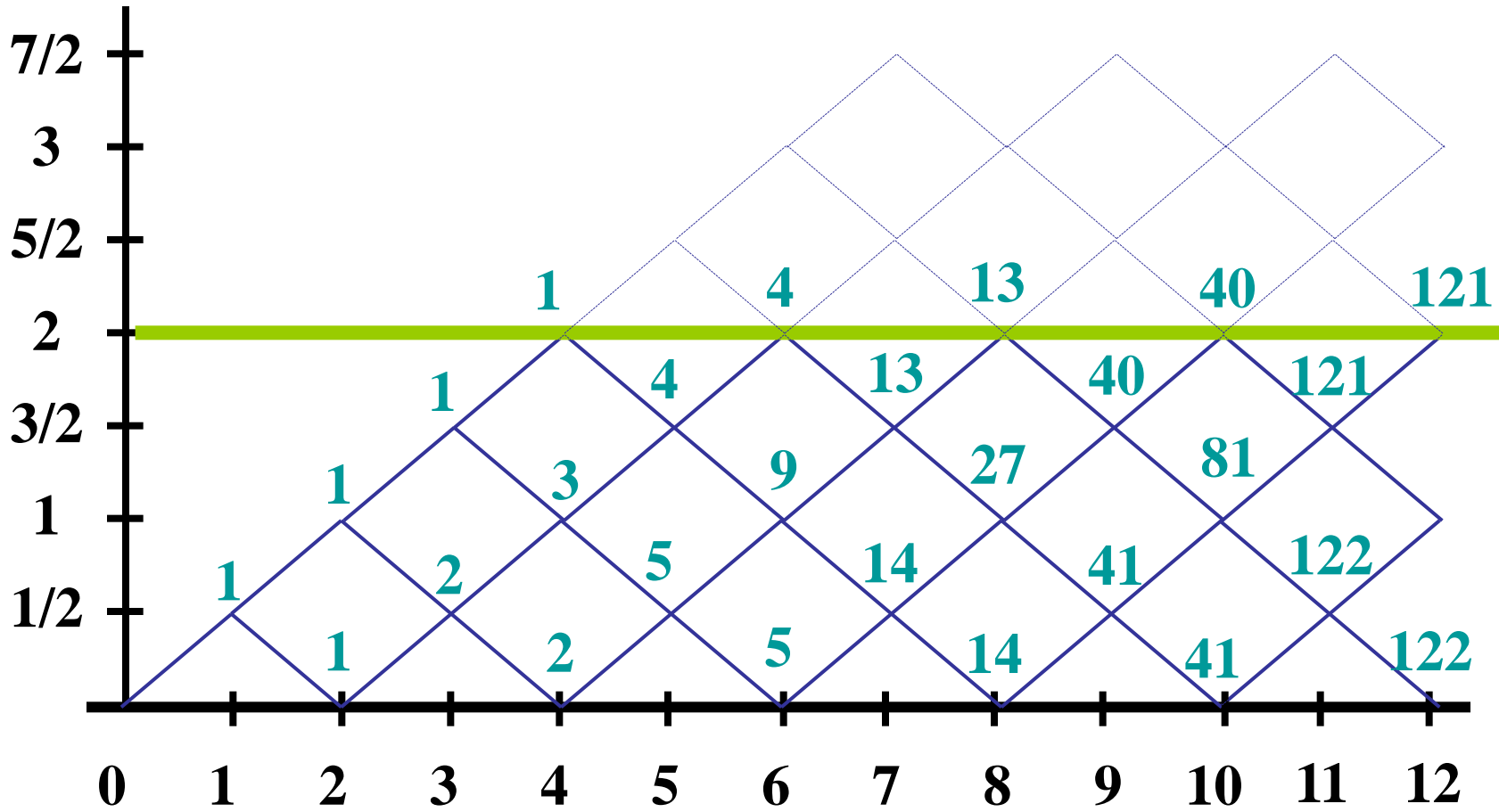
α  $_0$ + β  $_1$



$\text{Dim}(N) \sim 2^N$

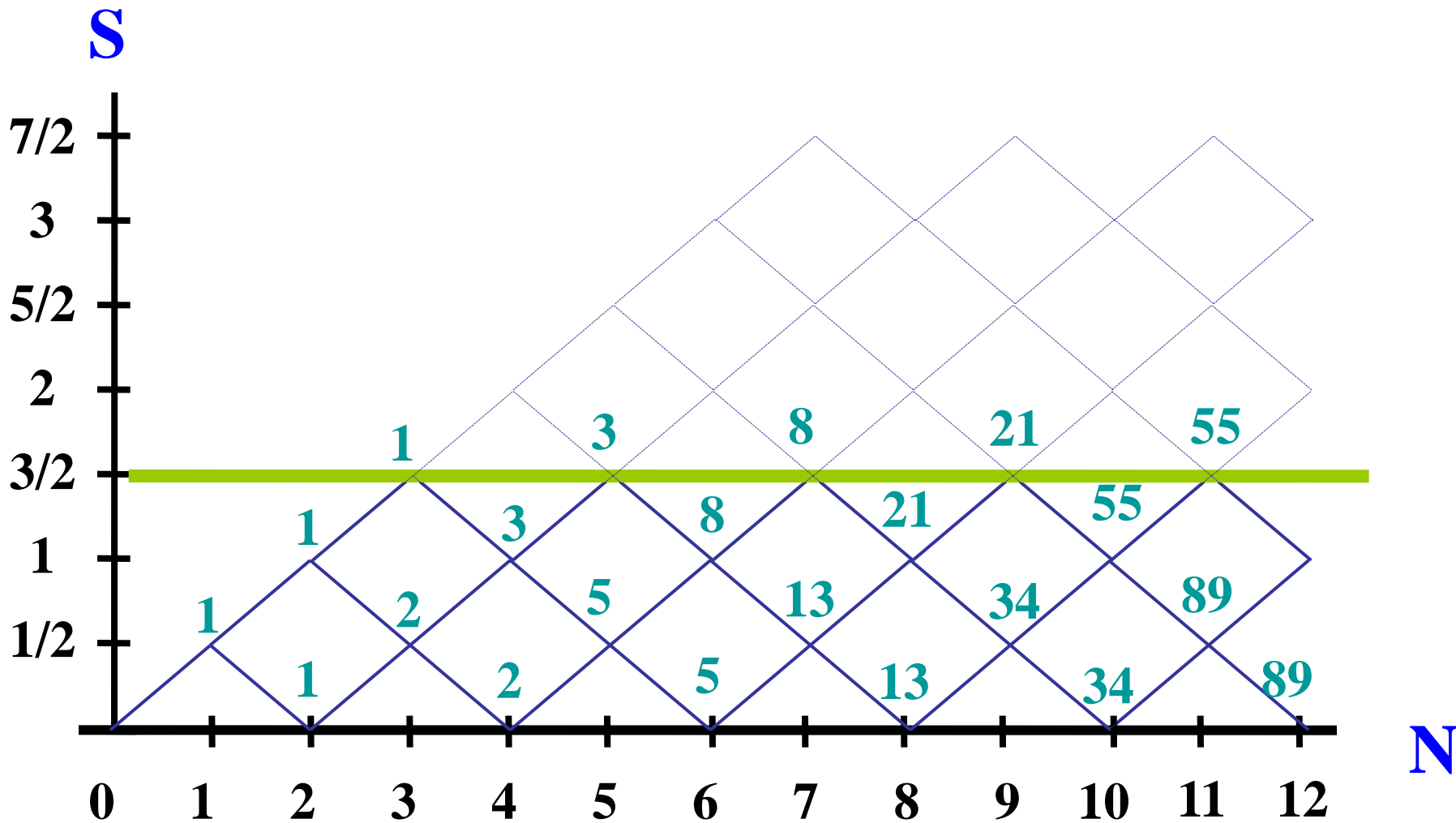
k = 4

S



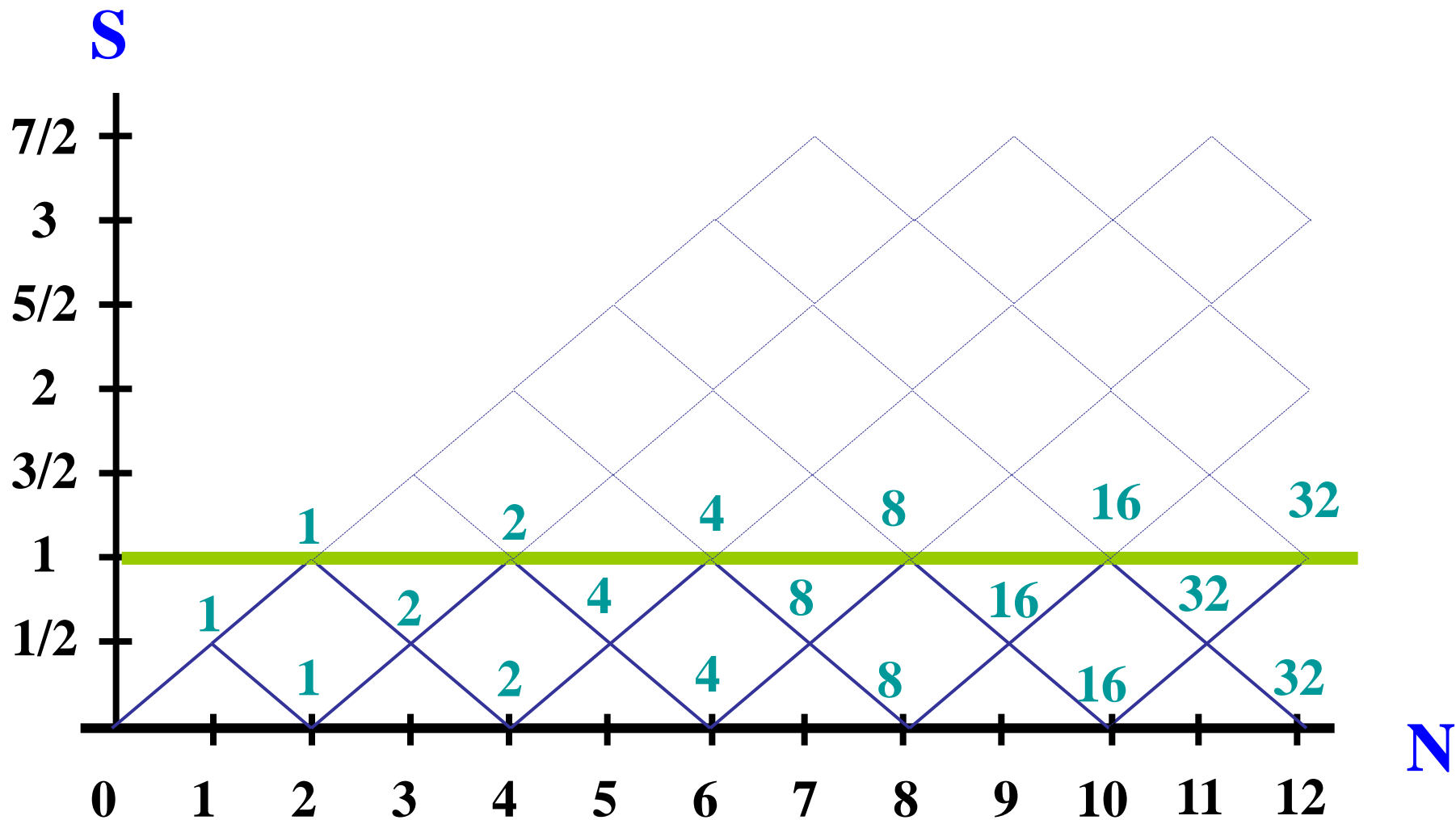
Dim(N) ~ 3^{N/2}

k = 3



Dim(N) = Fib(N+1) ~ ϕ^N

k = 2



Dim(N) = $2^{N/2}$

The F Matrix

$$\sum_a F_{ab}^c$$

$$\left(\begin{array}{cc|c} \frac{1}{[2]_q} & \frac{\sqrt{[3]_q}}{[2]_q} & 0 \\ \frac{\sqrt{[3]_q}}{[2]_q} & -\frac{1}{[2]_q} & 0 \\ \hline 0 & 0 & 1 \end{array} \right)$$

q-integers: $[m]_q = \frac{q^{m/2} - q^{-m/2}}{q^{1/2} - q^{-1/2}} ; q = e^{i2\pi/(k+2)}$

The R Matrix

Diagram 1 (State 0):

$$\text{Oval } \mathbf{0} \begin{array}{c} \text{---} \\ \text{---} \end{array} \begin{array}{c} \text{---} \\ \text{---} \end{array} = e^{i \frac{2k+1}{2k+4} \pi} \text{Oval } \mathbf{0} \begin{array}{c} \text{---} \\ \text{---} \end{array} \begin{array}{c} \text{---} \\ \text{---} \end{array}$$

Diagram 2 (State 1):

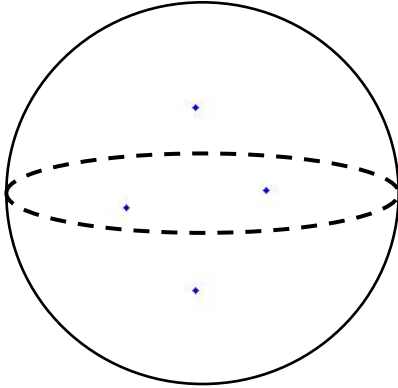
$$\text{Oval } \mathbf{1} \begin{array}{c} \text{---} \\ \text{---} \end{array} \begin{array}{c} \text{---} \\ \text{---} \end{array} = e^{i \frac{1}{2k+4} \pi} \text{Oval } \mathbf{1} \begin{array}{c} \text{---} \\ \text{---} \end{array} \begin{array}{c} \text{---} \\ \text{---} \end{array}$$

R matrix

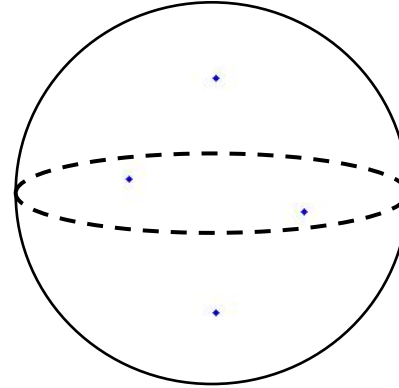
$$R = \begin{pmatrix} e^{i \frac{2k+1}{2k+4} \pi} & 0 \\ 0 & e^{i \frac{1}{2k+4} \pi} \end{pmatrix}$$

$N = 1$

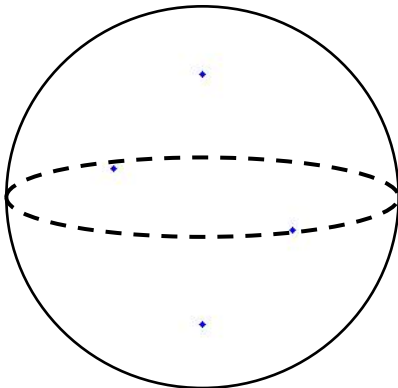
$k=2$



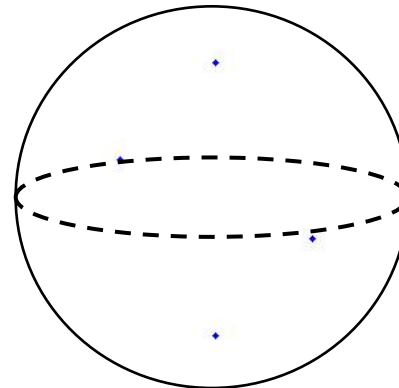
$k=3$



$k=4$

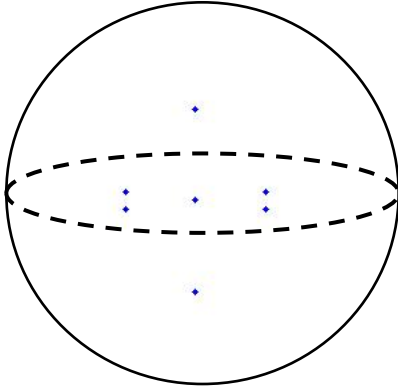


$k=5$

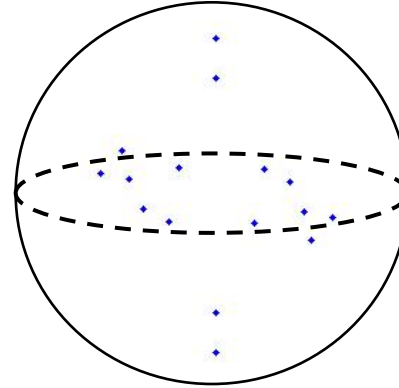


$N = 2$

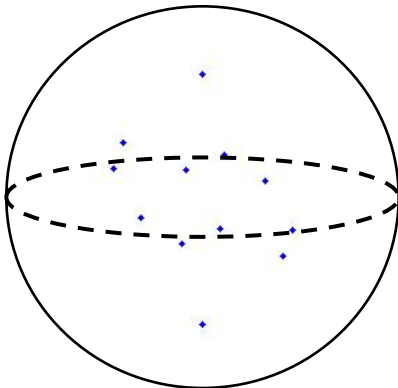
$k=2$



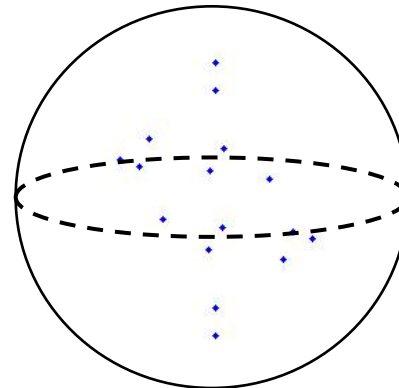
$k=3$



$k=4$

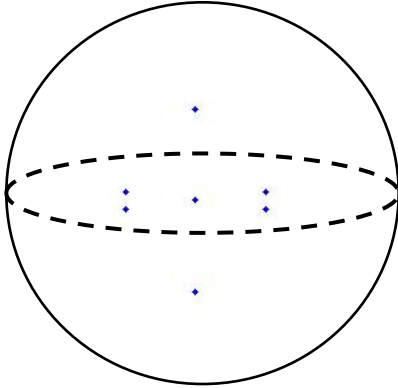


$k=5$

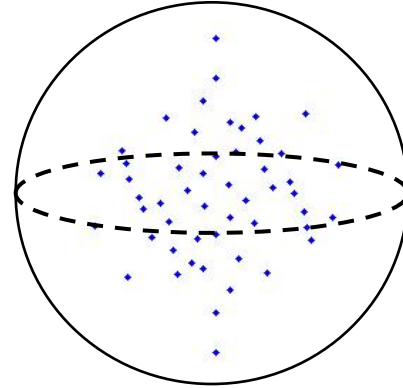


$N = 3$

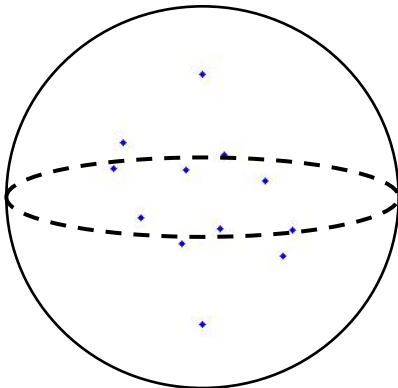
$k=2$



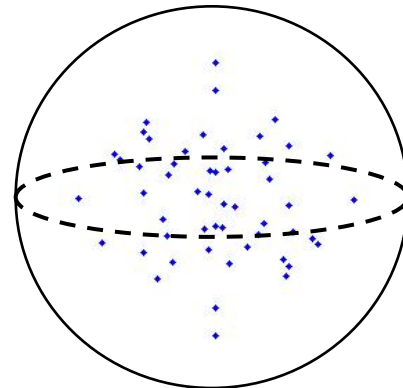
$k=3$



$k=4$

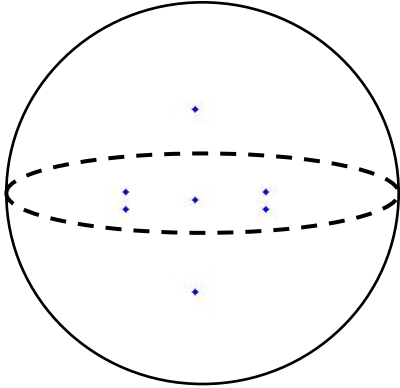


$k=5$

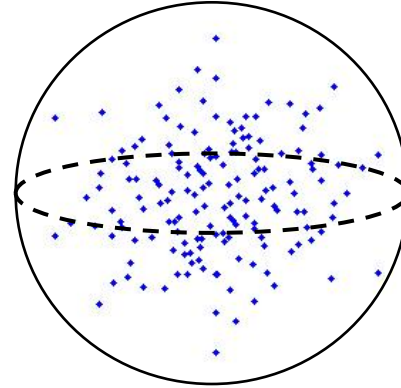


$N = 4$

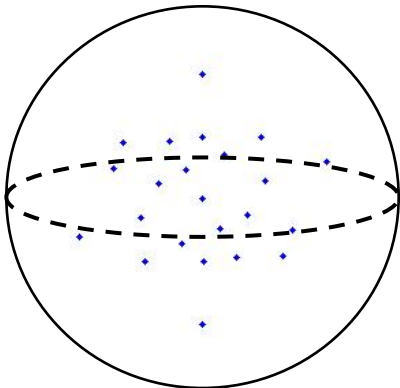
$k=2$



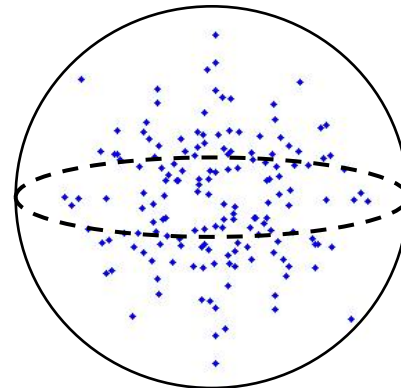
$k=3$



$k=4$

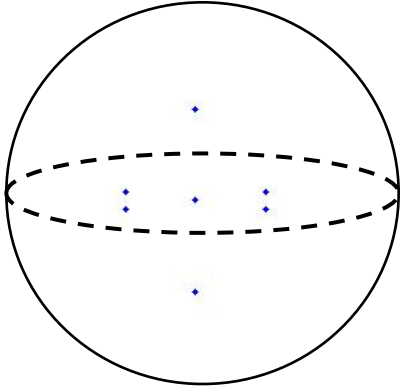


$k=5$

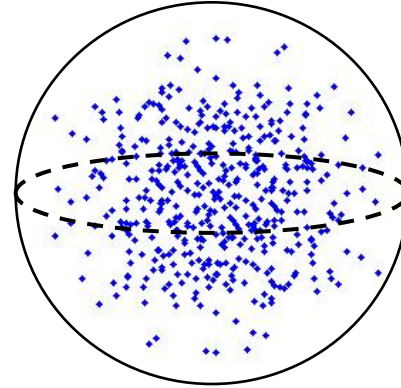


$N = 5$

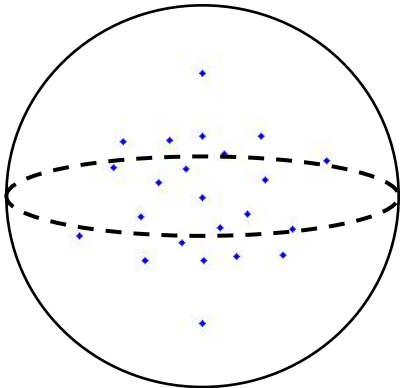
$k=2$



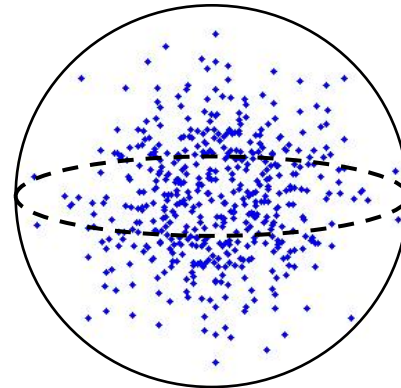
$k=3$



$k=4$

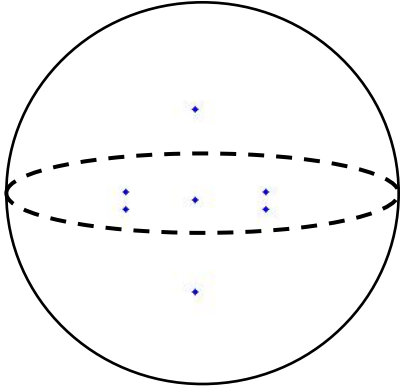


$k=5$

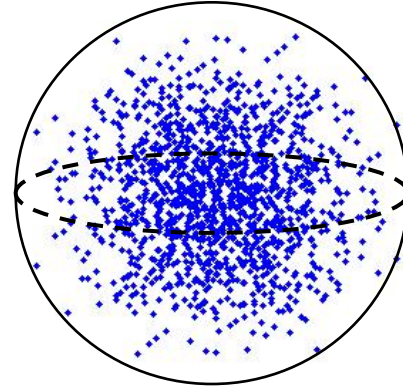


$N = 6$

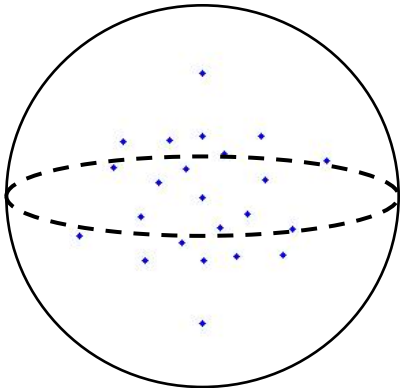
$k=2$



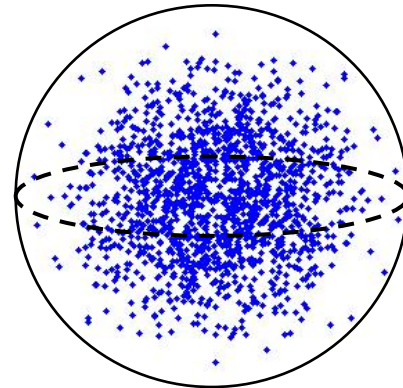
$k=3$



$k=4$

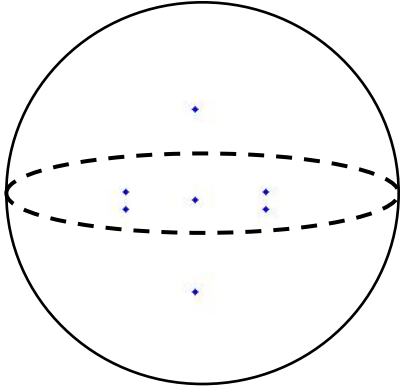


$k=5$

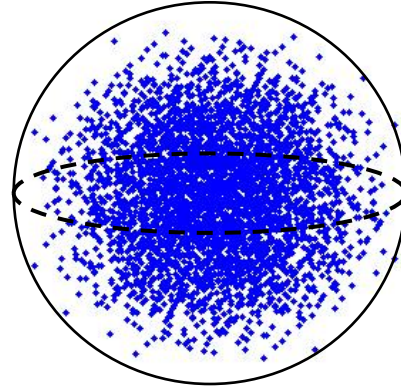


$N = 7$

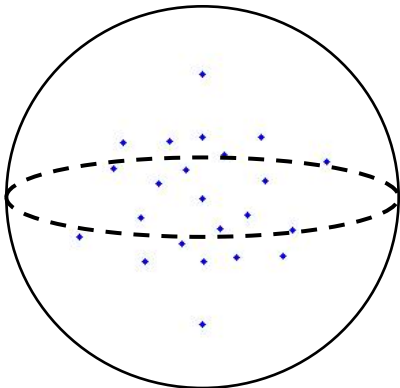
$k=2$



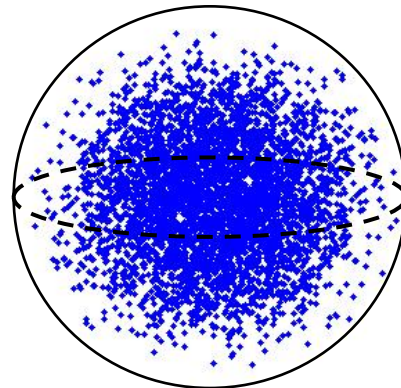
$k=3$



$k=4$

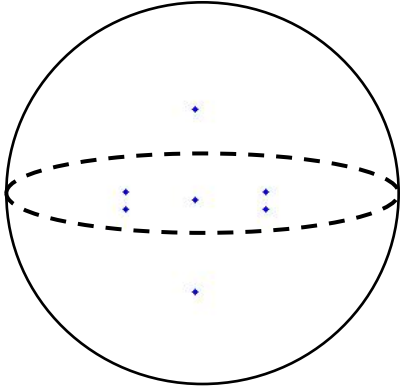


$k=5$

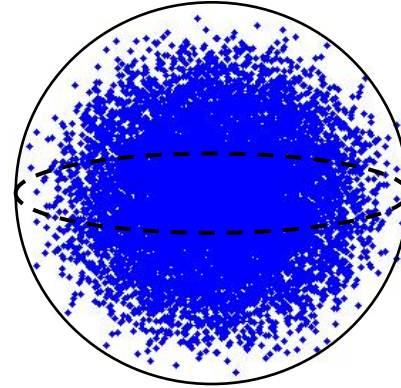


$N = 8$

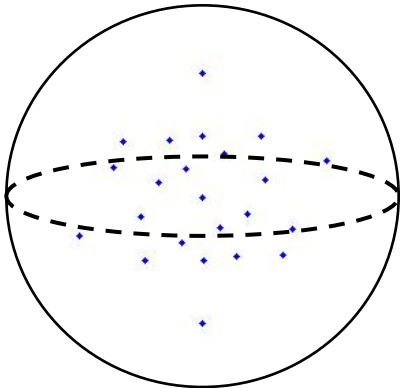
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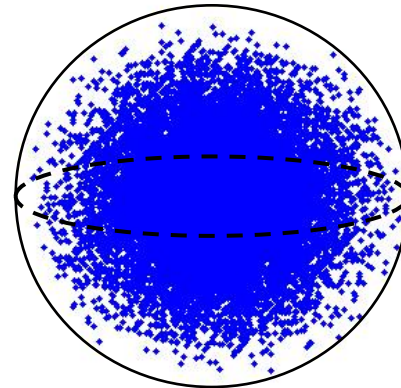
$k=3$



$k=4$

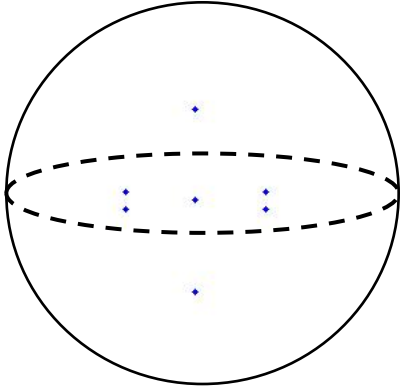


$k=5$

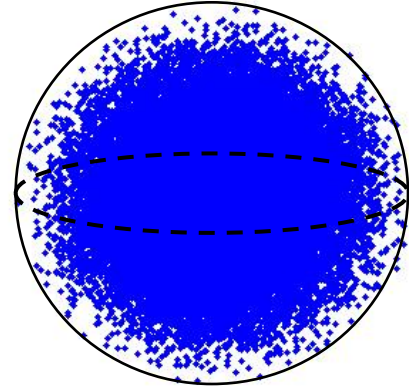


$N = 9$

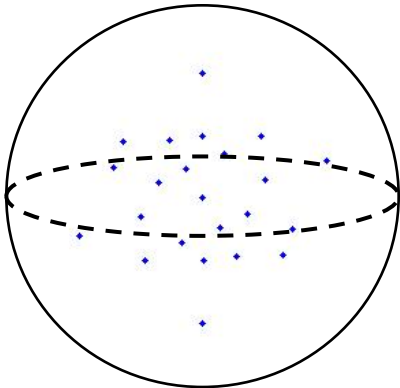
$k=2$



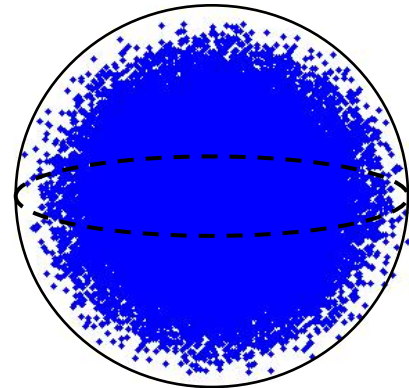
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$k=4$

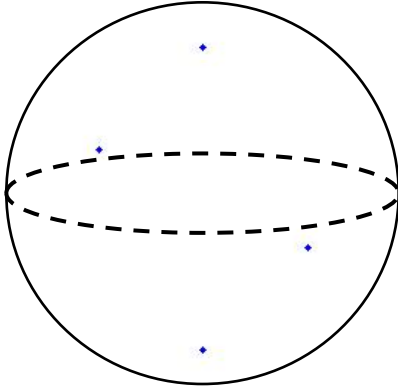


$k=5$

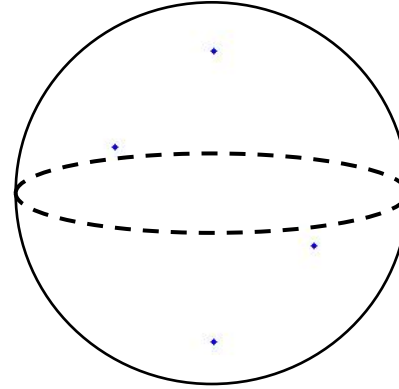


$N = 1$

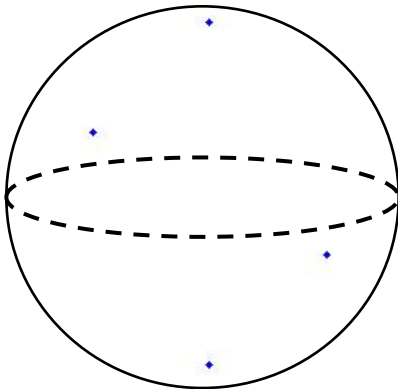
$k=6$



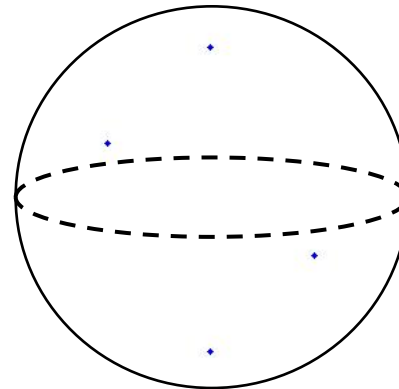
$k=7$



$k=8$

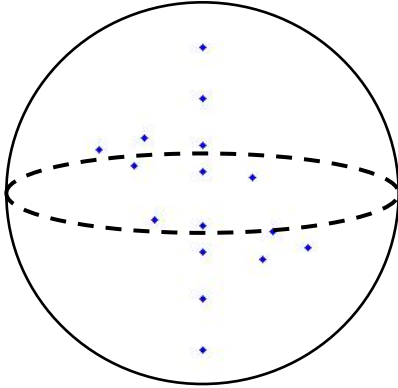


$k=9$

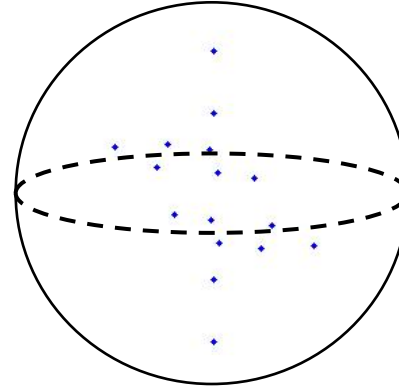


$N = 2$

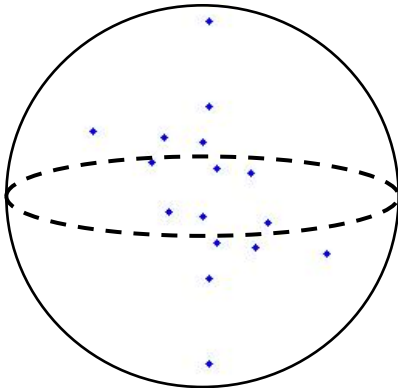
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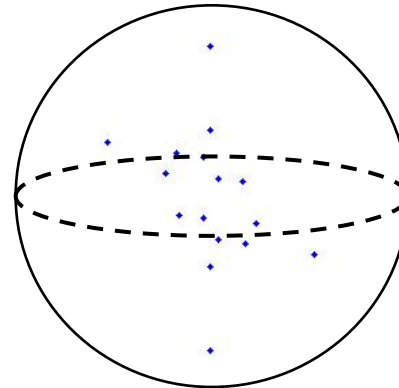
$k=7$



$k=8$

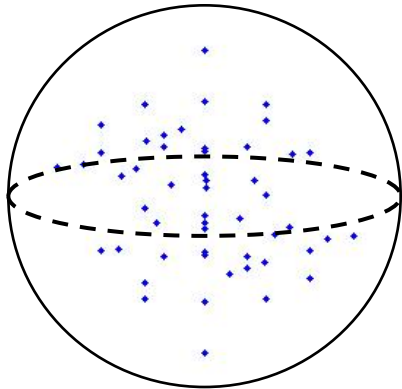


$k=9$

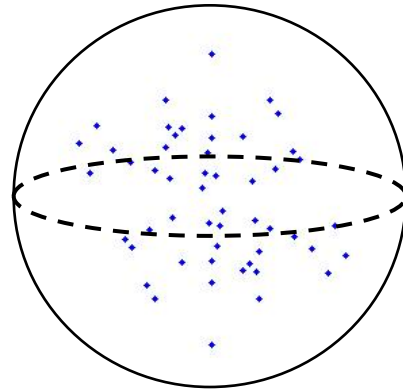


$N = 3$

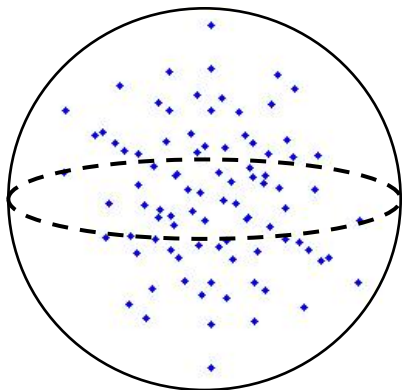
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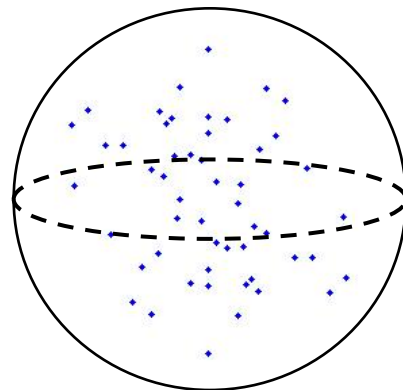
$k=7$



$k=8$

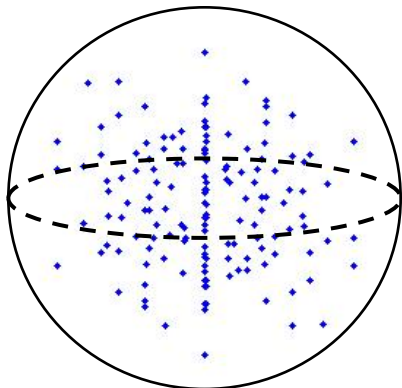


$k=9$

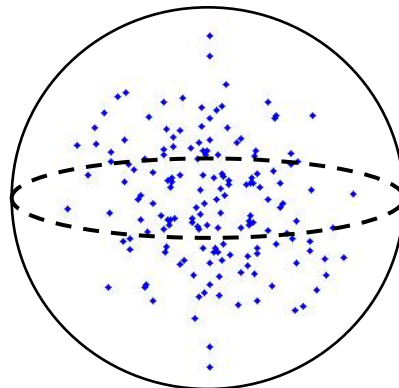


$N = 4$

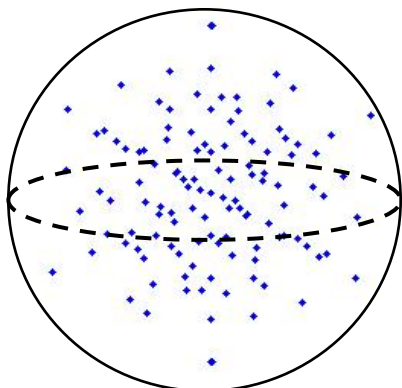
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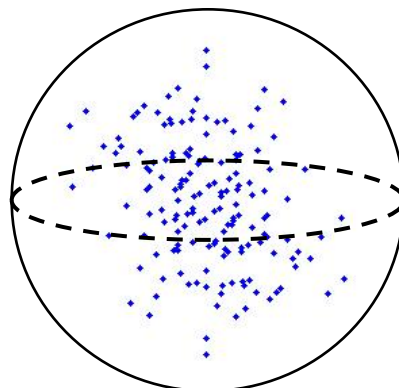
$k=7$



$k=8$

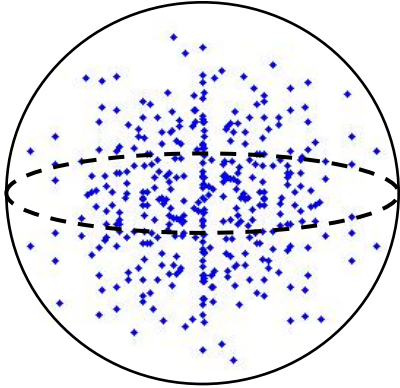


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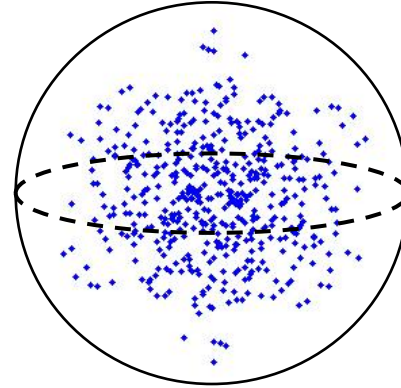


$N = 5$

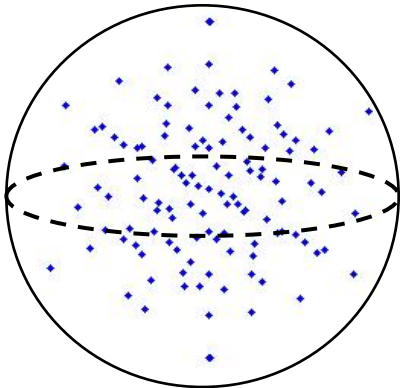
$k=6$



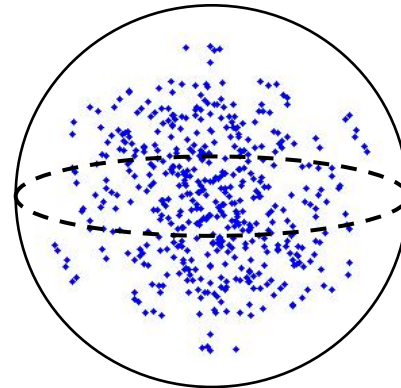
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$k=8$

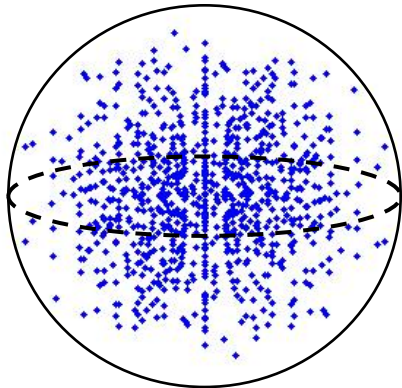


$k=9$

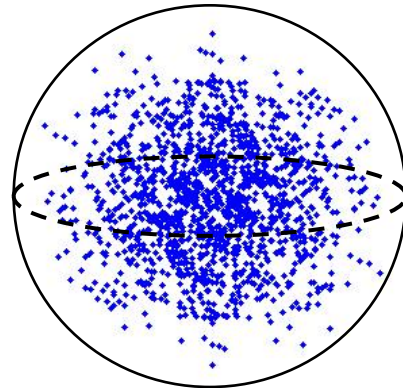


$N = 6$

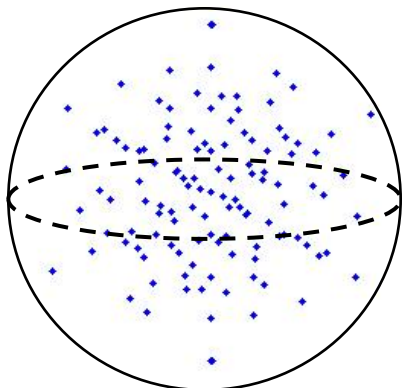
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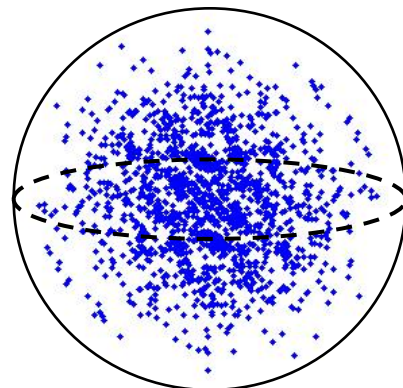
$k=7$



$k=8$

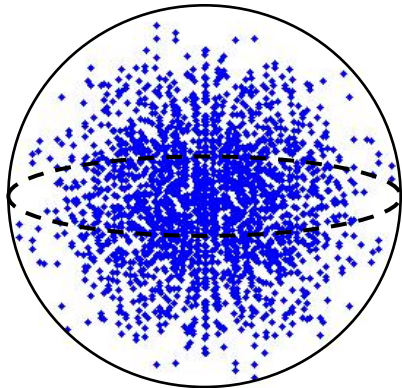


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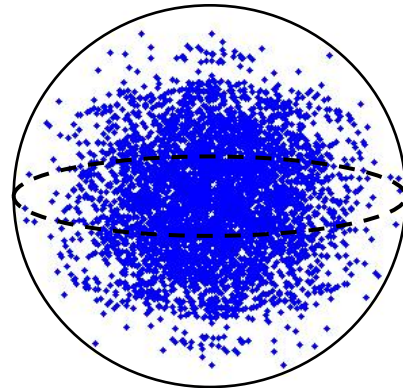


$N = 7$

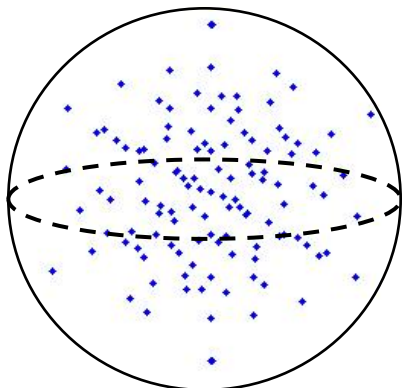
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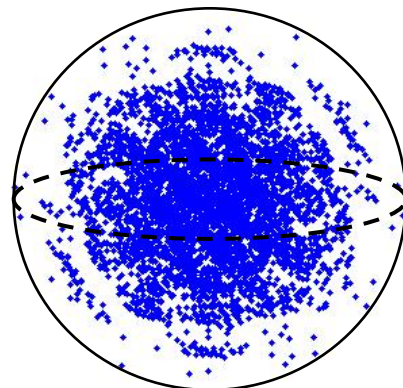
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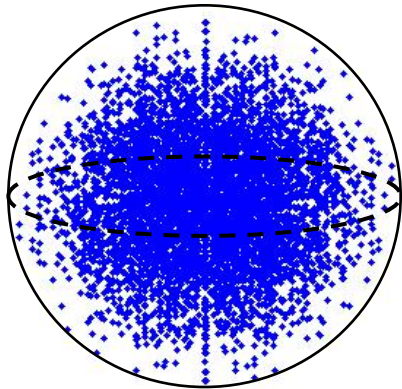


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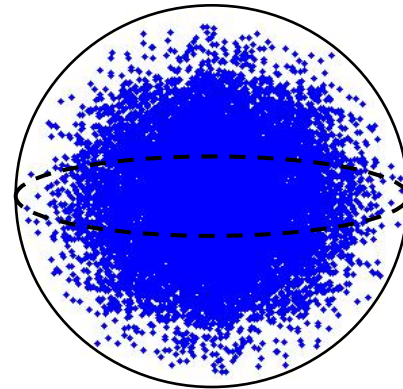


$N = 8$

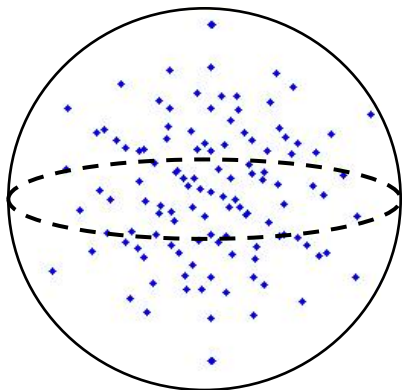
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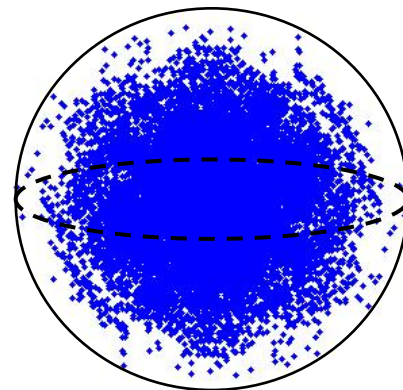
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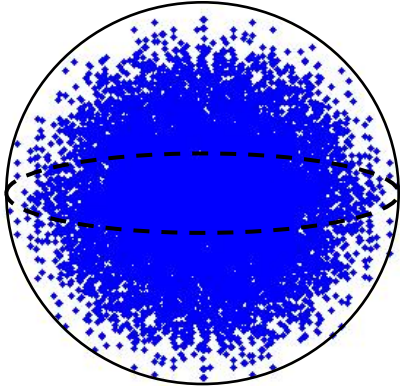


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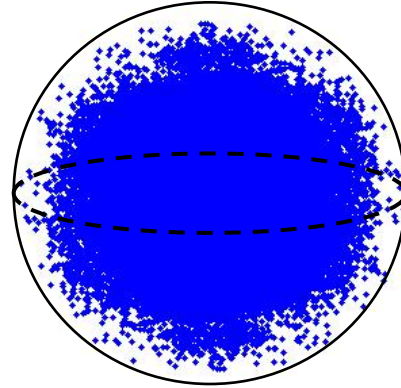


$N = 9$

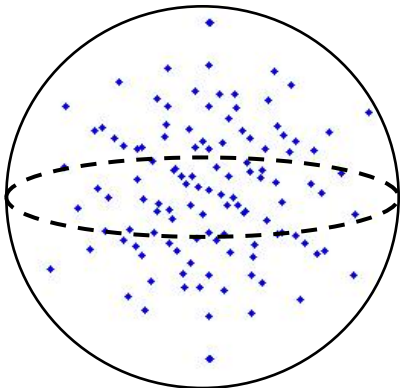
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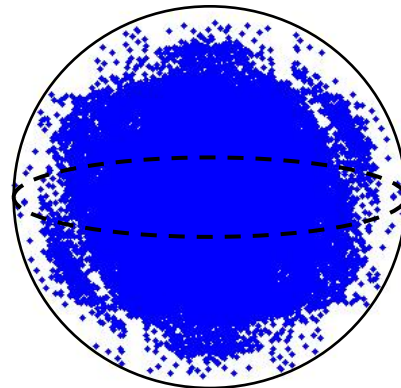
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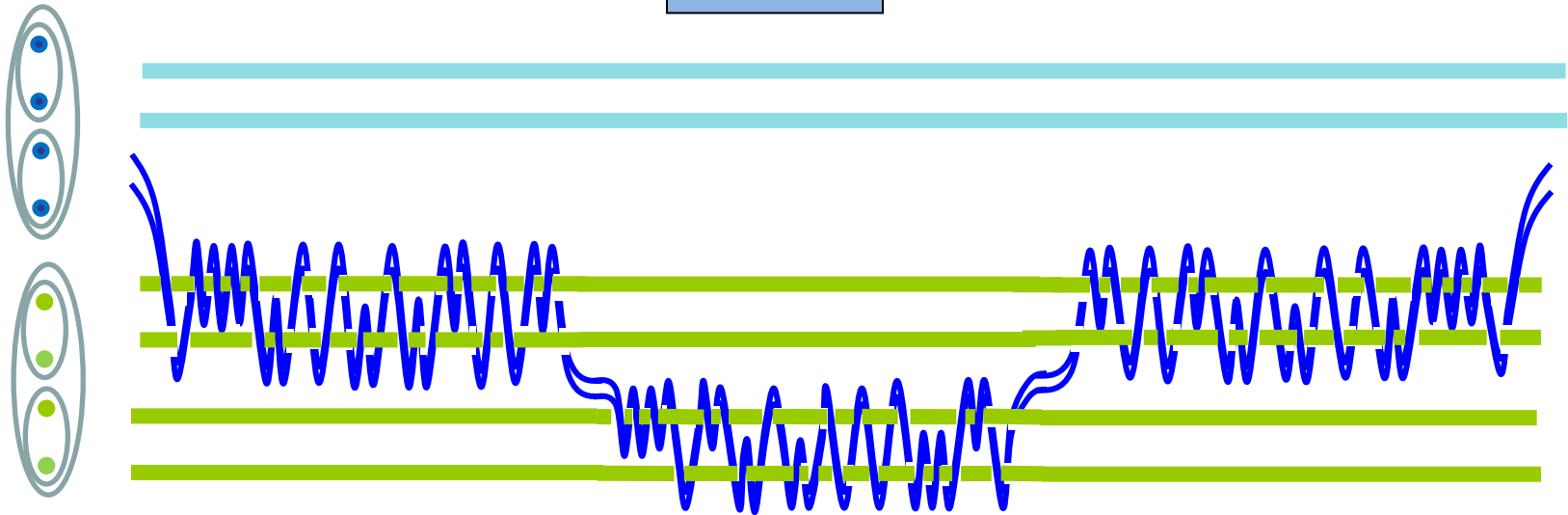
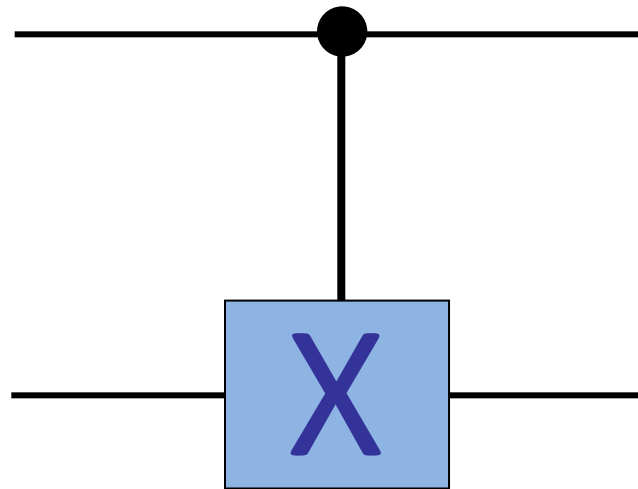
$k=8$



$k=9$



CNOT gate for $SU(2)_5$



“Surface Code” Approach

Key Idea: Use a quantum computer to *simulate* a theory of anyons.

The simulation “inherits” the fault-tolerance of the anyon theory.

Most promising approach is based on “Abelian anyons.” High error thresholds, but can’t compute purely by braiding.

Bravyi & Kitaev, 2003

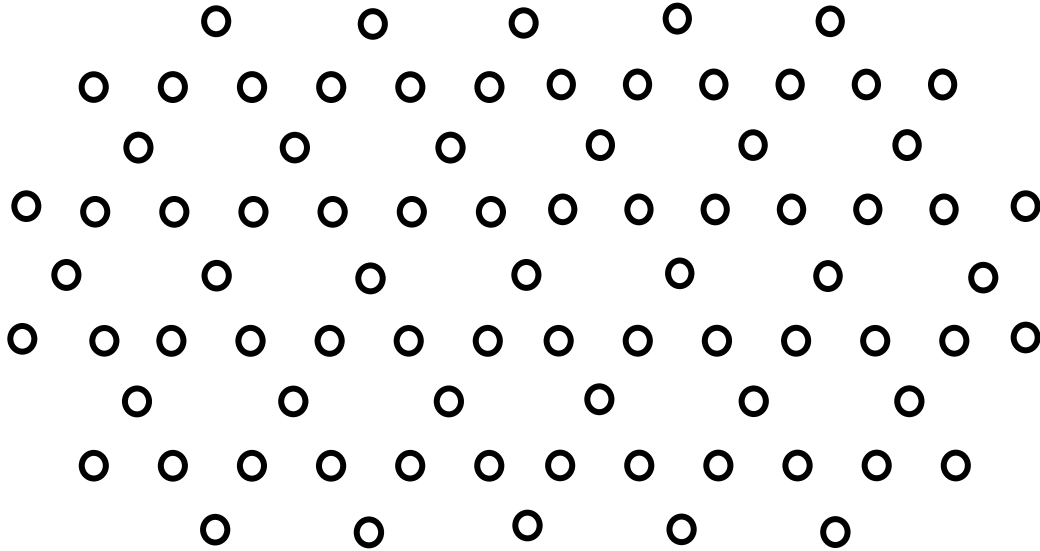
Raussendorf & Harrington, PRL 2007

Fowler, Stephens, and Groszkowski, PRA 2009

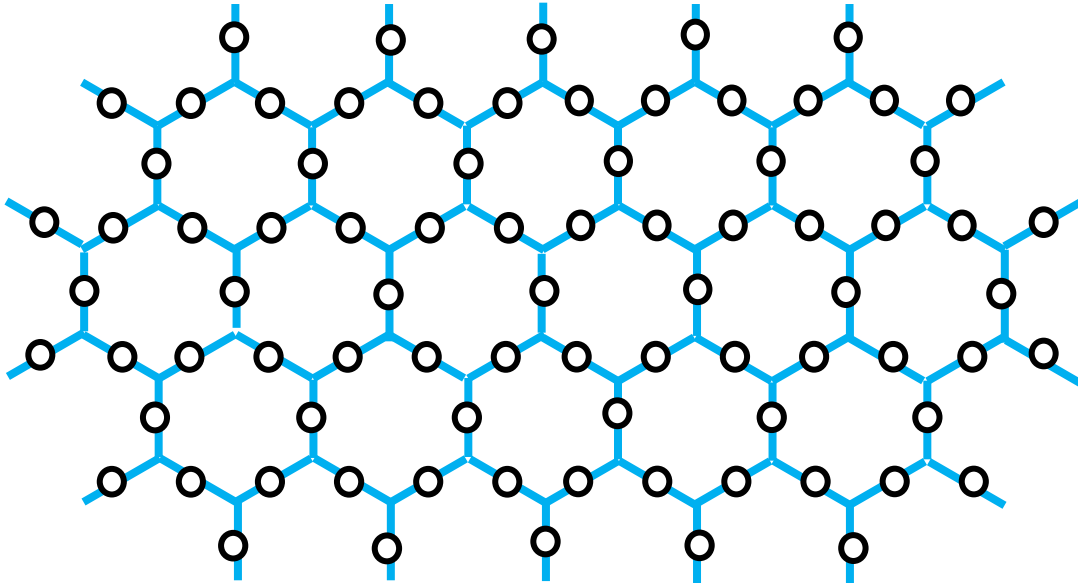
More speculative idea: simulate Fibonacci anyons.

Konig, Kuperberg, Reichardt, Ann. Phys. 2010

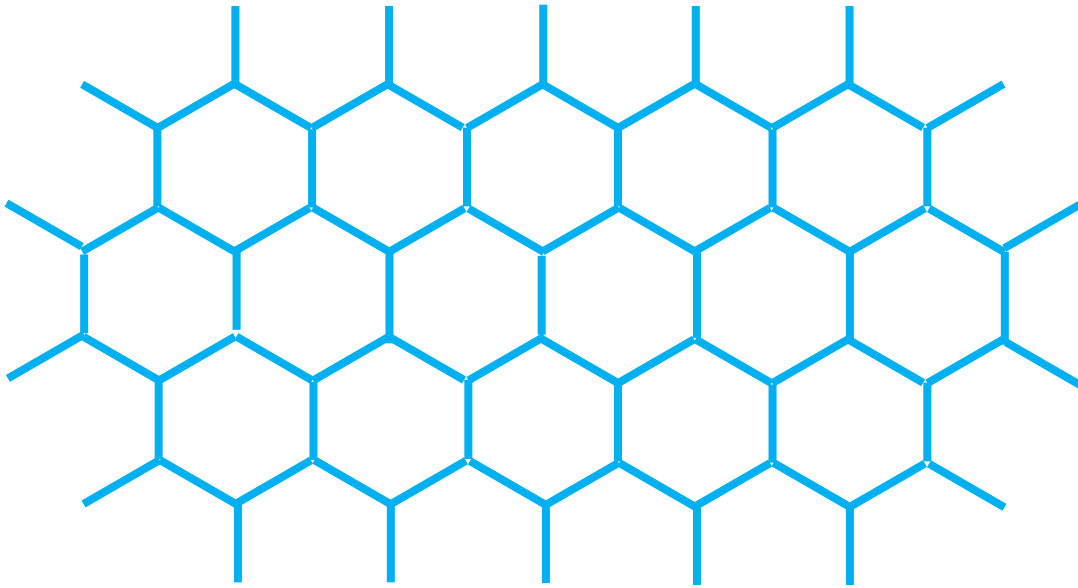
NEB, D.P. DiVincenzo, PRB 2012



2D Array of Qubits



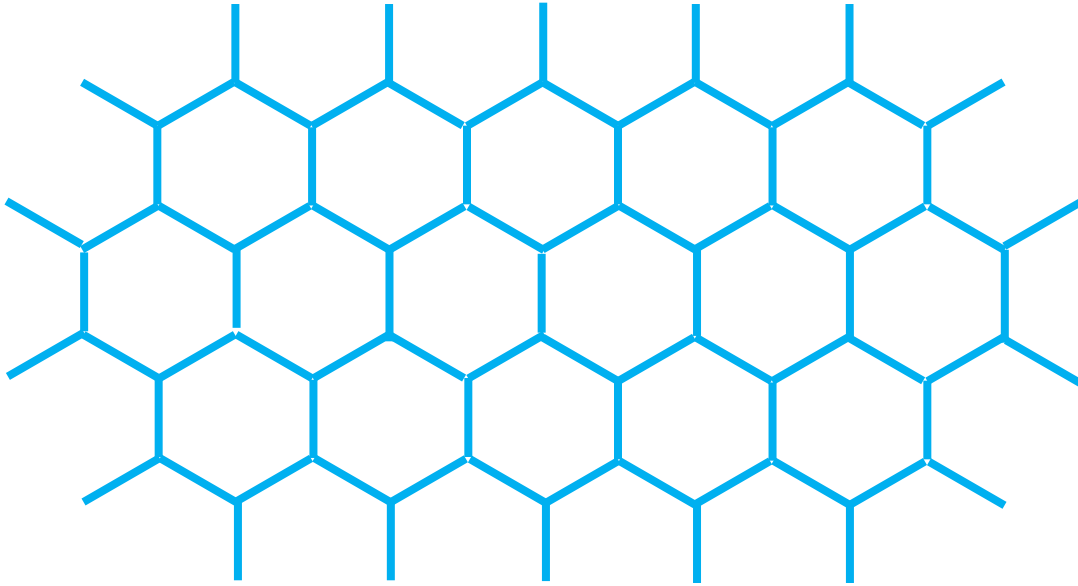
Trivalent Lattice



Trivalent Lattice

“Fibonacci” Levin-Wen Model

Levin & Wen, PRB 2005

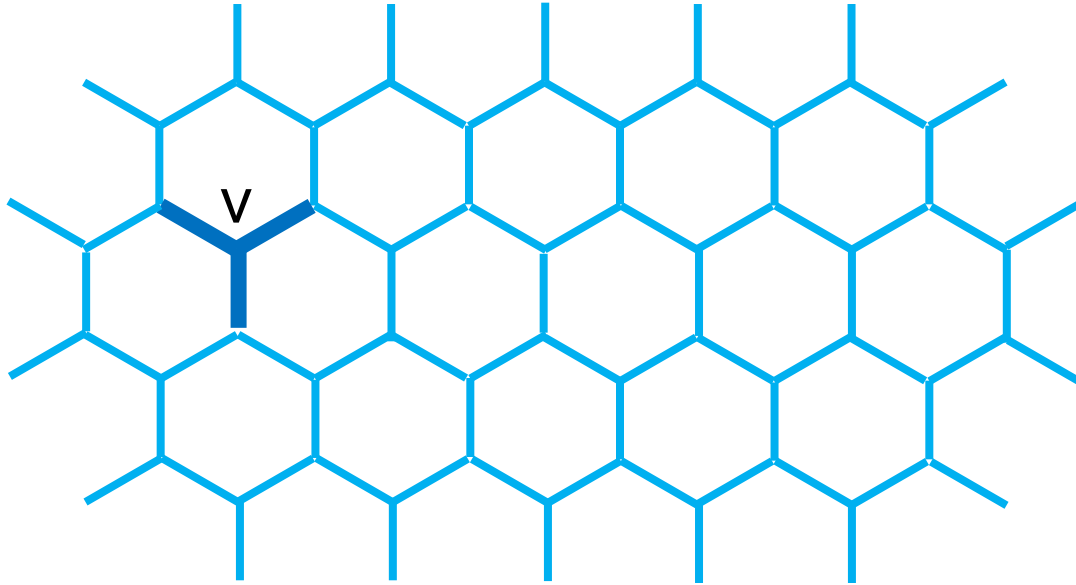


Trivalent Lattice

$$H = - \sum_v Q_v - \sum_p B_p$$

“Fibonacci” Levin-Wen Model

Levin & Wen, PRB 2005



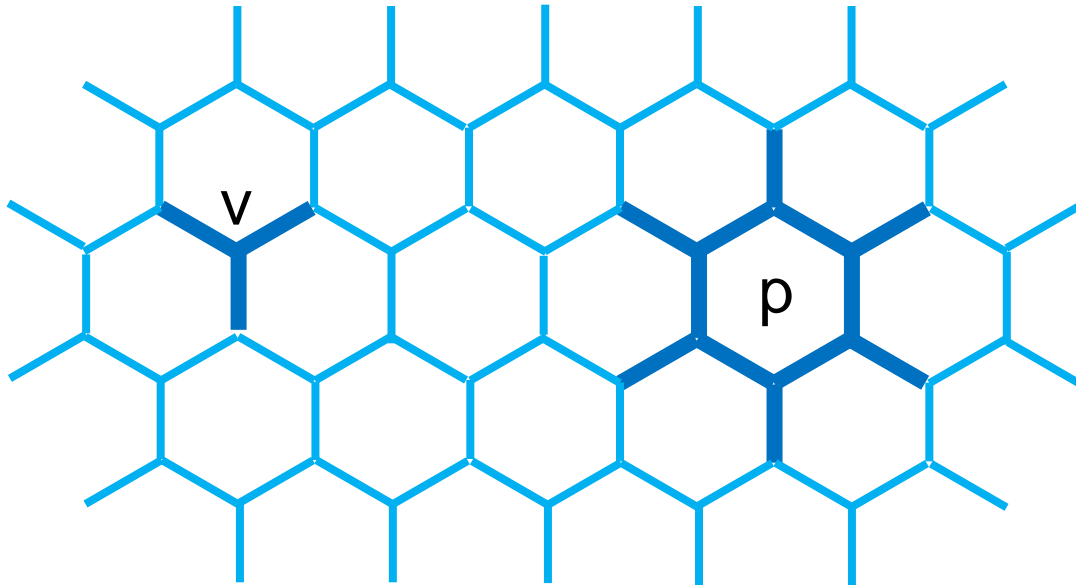
Trivalent Lattice

$$H = - \sum_v Q_v - \sum_p B_p$$

Vertex
Operator

$$Q_v = 0, 1$$

“Fibonacci” Levin-Wen Model Levin & Wen, PRB 2005



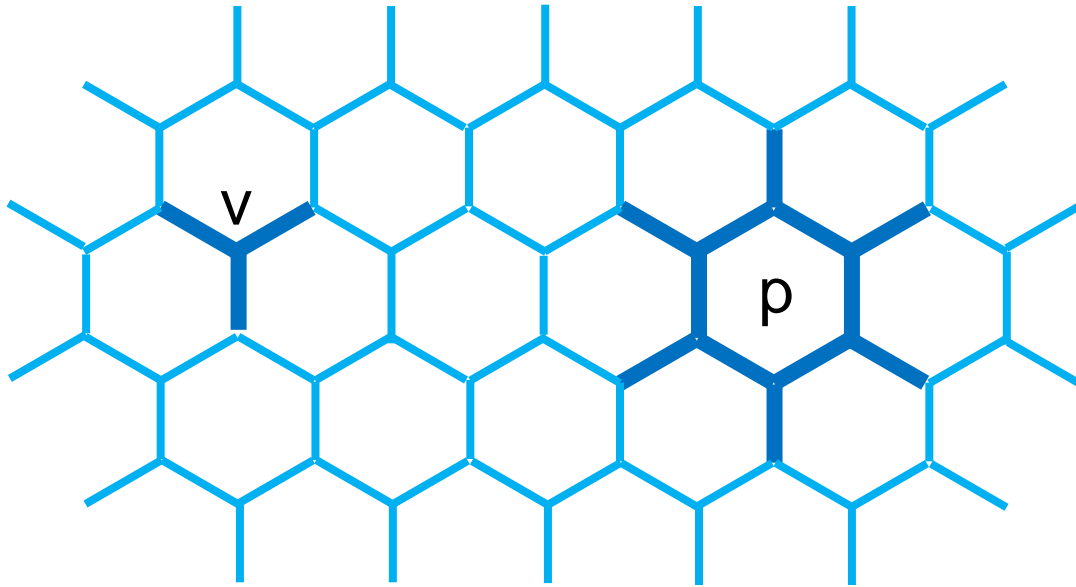
Trivalent Lattice

$$H = - \sum_v Q_v - \sum_p B_p$$

Vertex Operator Plaquette Operator

$Q_v = 0, 1$ $B_p = 0, 1$

“Fibonacci” Levin-Wen Model Levin & Wen, PRB 2005



Trivalent Lattice

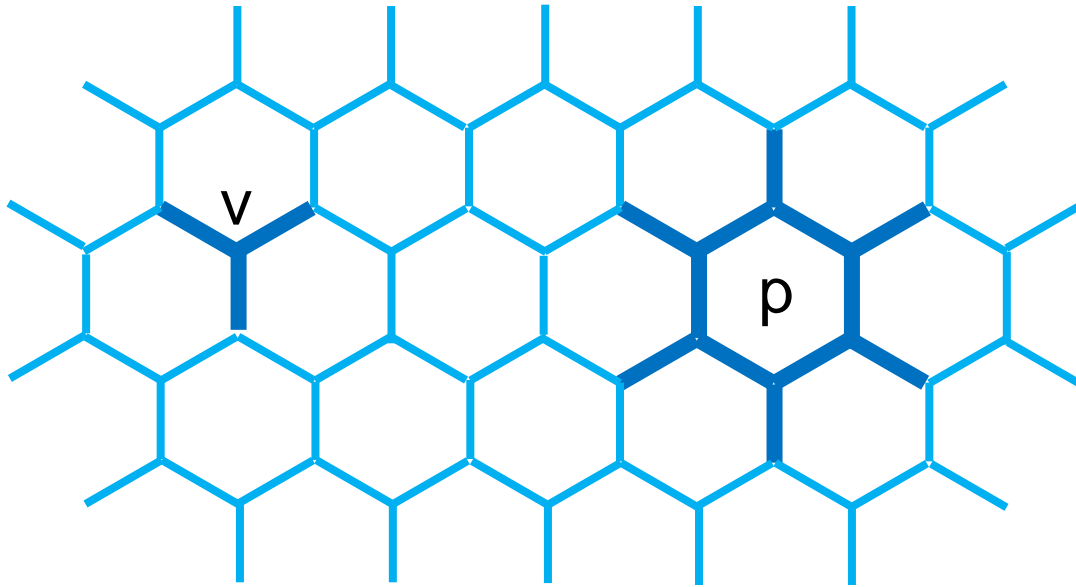
$$H = - \sum_v Q_v - \sum_p B_p$$

Vertex Operator $Q_v = 0, 1$

Plaquette Operator $B_p = 0, 1$

Ground State
 $Q_v = 1$ on each vertex
 $B_p = 1$ on each plaquette

“Fibonacci” Levin-Wen Model Levin & Wen, PRB 2005



Trivalent Lattice

$$H = - \sum_v Q_v - \sum_p B_p$$

Vertex Operator $Q_v = 0, 1$

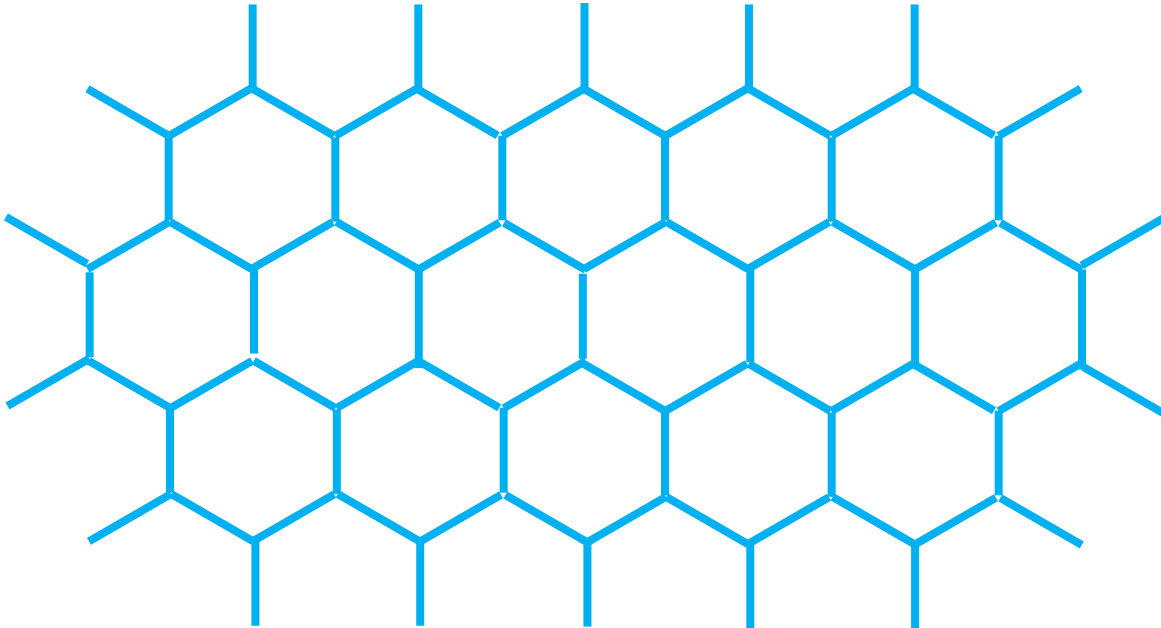
Plaquette Operator $B_p = 0, 1$

Ground State
 $Q_v = 1$ on each vertex
 $B_p = 1$ on each plaquette

Excited States are Fibonacci Anyons

Vertex Operator: Q_v

$$H = - \sum_v Q_v - \sum_p B_p$$



$$Q_v \left| \begin{array}{c} j \\ i \text{---} \mathbf{v} \\ k \end{array} \right\rangle = \delta_{ijk} \left| \begin{array}{c} j \\ i \text{---} \mathbf{v} \\ k \end{array} \right\rangle$$

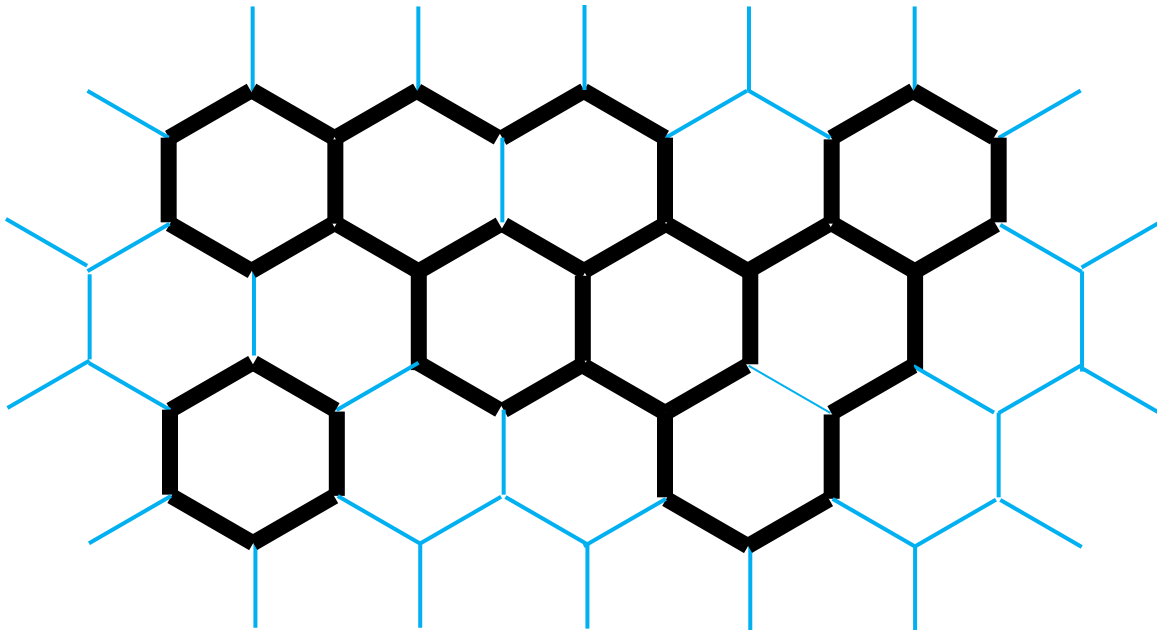
“Fibonacci” Levin-Wen Model

$$\delta_{100} = \delta_{010} = \delta_{001} = 0$$

All other $\delta_{ijk} = 1$

Vertex Operator: Q_v

$$H = - \sum_v Q_v - \sum_p B_p$$



$$\text{—} = |0\rangle$$

$$\text{—} = |1\rangle$$

$Q_v = 1$
on each vertex
→ “branching”
loop states.

$$Q_v \left| \begin{array}{c} j \\ i \text{---} \text{v} \\ k \end{array} \right\rangle = \delta_{ijk} \left| \begin{array}{c} j \\ i \text{---} \text{v} \\ k \end{array} \right\rangle$$

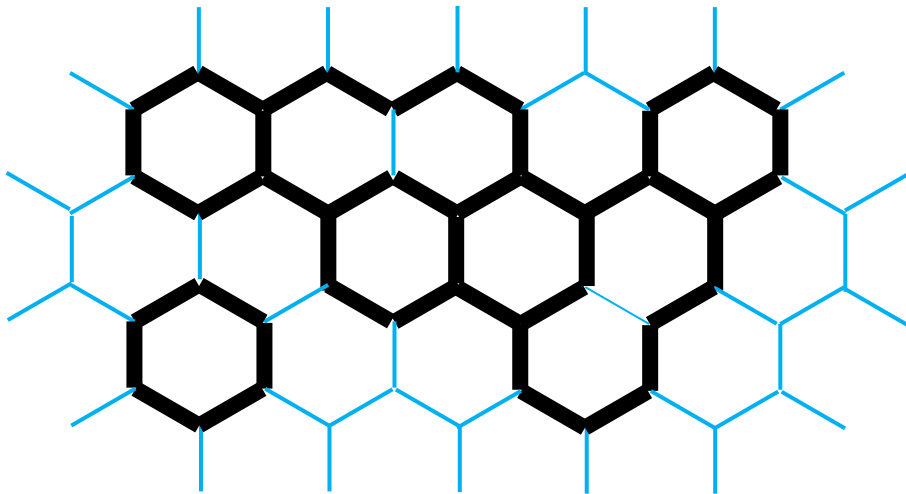
“Fibonacci” Levin-Wen Model

$$\delta_{100} = \delta_{010} = \delta_{001} = 0$$

All other $\delta_{ijk} = 1$

Plaquette Operator: B_p

$$H = - \sum_v Q_v - \sum_p B_p$$



$B_p = 1$
on each plaquette
→ superposition
of loop states

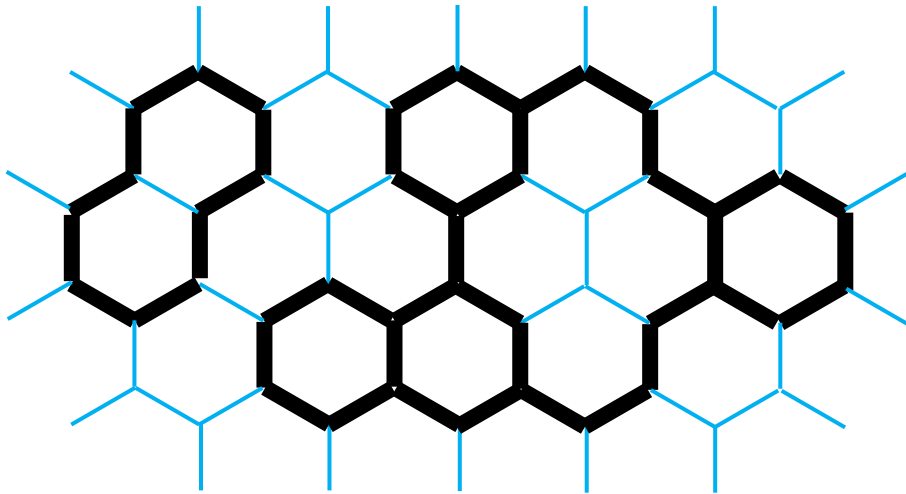
$$B_p^s \left| \begin{array}{c} b \\ a \quad i \quad j \quad c \\ \quad n \quad p \quad k \\ f \quad m \quad l \quad d \\ \quad e \end{array} \right\rangle = \sum_{i'j'k'l'm'n'} B_{p,ijklmn}^{s,i'j'k'l'm'n'} (abcdef) \left| \begin{array}{c} b \\ a \quad i' \quad j' \quad c \\ \quad n' \quad p \quad k' \\ f \quad m' \quad l' \quad d \\ \quad e \end{array} \right\rangle$$

$$B_p = \frac{B_p^0 + \varphi B_p^1}{1 + \varphi^2} \quad B_{p,ijklmn}^{s,i'j'k'l'm'n'} (abcdef) = F_{si'n'}^{ani} F_{sj'i'}^{bij} F_{sk'j'}^{cjk} F_{sl'k'}^{dkl} F_{sm'l'}^{elm} F_{sn'm'}^{fmn}$$

Very Complicated 12-qubit Interaction!

Plaquette Operator: B_p

$$H = - \sum_v Q_v - \sum_p B_p$$



$B_p = 1$
 on each plaquette
 → superposition
 of loop states

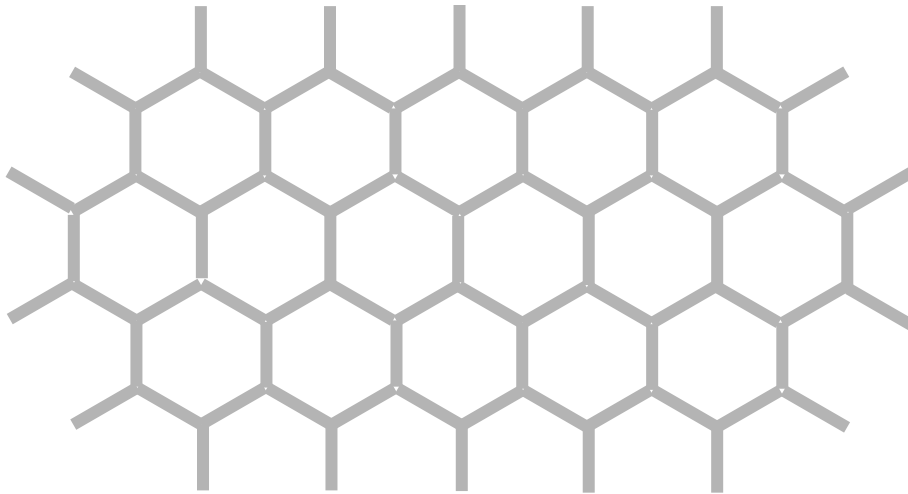
$$B_p^s \left| \begin{array}{c} b \\ a \quad i \quad j \quad c \\ \quad n \quad p \quad k \\ f \quad m \quad l \quad d \\ \quad \quad e \end{array} \right\rangle = \sum_{i'j'k'l'm'n'} B_{p,ijklmn}^{s,i'j'k'l'm'n'} (abcdef) \left| \begin{array}{c} b \\ a \quad i' \quad j' \quad c \\ \quad n' \quad p \quad k' \\ f \quad m' \quad l' \quad d \\ \quad \quad e \end{array} \right\rangle$$

$$B_p = \frac{B_p^0 + \varphi B_p^1}{1 + \varphi^2} \quad B_{p,ijklmn}^{s,i'j'k'l'm'n'} (abcdef) = F_{si'n'}^{ani} F_{sj'i'}^{bij} F_{sk'j'}^{cjk} F_{sl'k'}^{dkl} F_{sm'l'}^{elm} F_{sn'm'}^{fmn}$$

Very Complicated 12-qubit Interaction!

Plaquette Operator: B_p

$$H = - \sum_v Q_v - \sum_p B_p$$



$B_p = 1$
on each plaquette
➔ superposition
of loop states

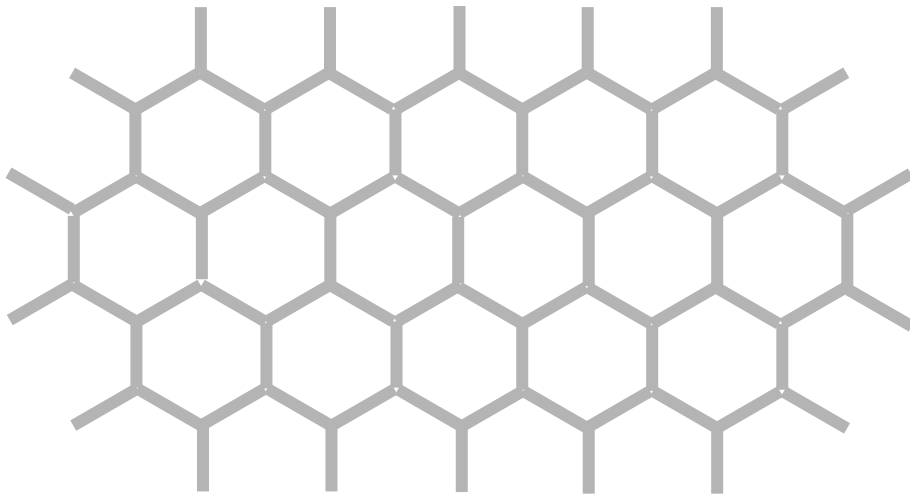
$$B_p^s \left| \begin{array}{c} b \\ a \quad i \quad j \quad c \\ \quad n \quad p \quad k \\ f \quad m \quad l \quad d \\ \quad e \end{array} \right\rangle = \sum_{i'j'k'l'm'n'} B_{p,ijklmn}^{s,i'j'k'l'm'n'} (abcdef) \left| \begin{array}{c} b \\ a \quad i' \quad j' \quad c \\ \quad n' \quad p \quad k' \\ f \quad m' \quad l' \quad d \\ \quad e \end{array} \right\rangle$$

$$B_p = \frac{B_p^0 + \varphi B_p^1}{1 + \varphi^2} \quad B_{p,ijklmn}^{s,i'j'k'l'm'n'} (abcdef) = F_{si'n'}^{ani} F_{sj'i'}^{bij} F_{sk'j'}^{cjk} F_{sl'k'}^{dkl} F_{sm'l'}^{elm} F_{sn'm'}^{fmn}$$

Very Complicated 12-qubit Interaction!

Plaquette Operator: B_p

$$H = - \sum_v Q_v - \sum_p B_p$$



$B_p = 1$
on each plaquette
➔ superposition
of loop states

$$B_p^s \left| \begin{array}{c} b \\ a \quad i \quad j \quad c \\ \quad n \quad p \quad k \\ f \quad m \quad l \quad d \\ \quad e \end{array} \right\rangle = \sum_{i'j'k'l'm'n'} B_{p,ijklmn}^{s,i'j'k'l'm'n'} (abcdef) \left| \begin{array}{c} b \\ a \quad i' \quad j' \quad c \\ \quad n' \quad p \quad k' \\ f \quad m' \quad l' \quad d \\ \quad e \end{array} \right\rangle$$

$$B_p = \frac{B_p^0 + \varphi B_p^1}{1 + \varphi^2} \quad B_{p,ijklmn}^{s,i'j'k'l'm'n'} (abcdef) = F_{si'n'}^{ani} F_{sj'i'}^{bij} F_{sk'j'}^{ckj} \boxed{F_{sl'k'}^{dkl}} F_{sm'l'}^{elm} F_{sn'm'}^{fmn}$$

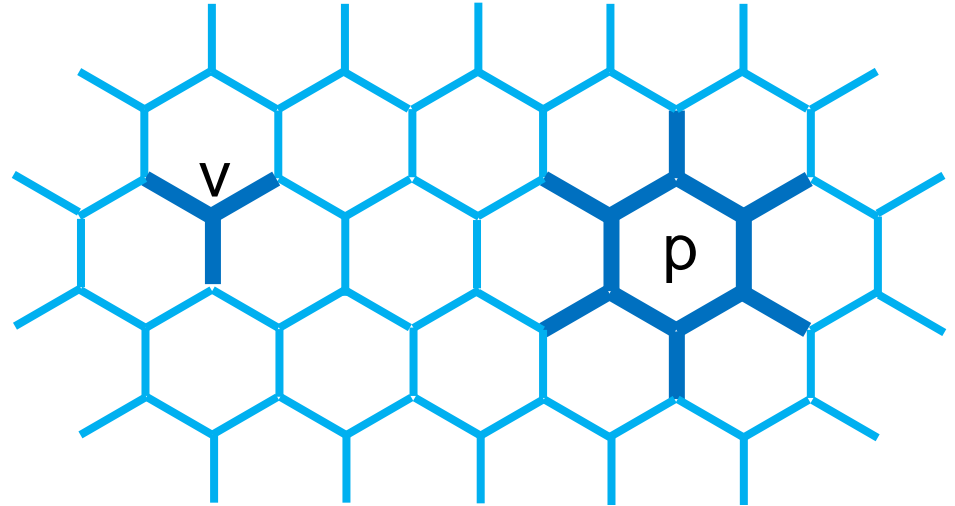
Very Complicated 12-qubit Interaction!

“Surface Code” Approach

Konig, Kuperberg, Reichardt, Ann. Phys. 2010

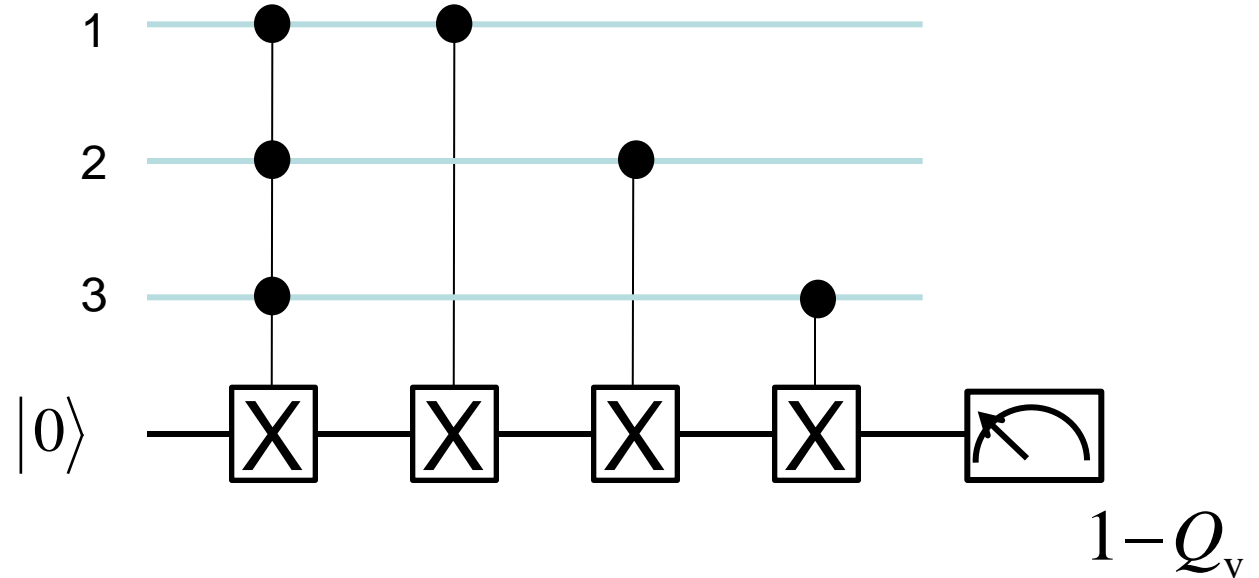
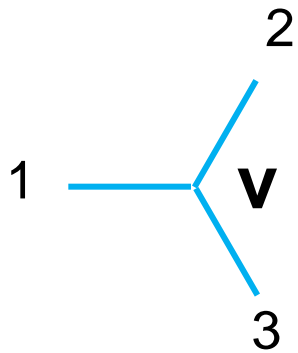
Use quantum computer to repeatedly measure Q_v and B_p .

If $Q_v=1$ on each vertex and $B_p=1$ on each plaquette, good. If not, treat as an “error.”

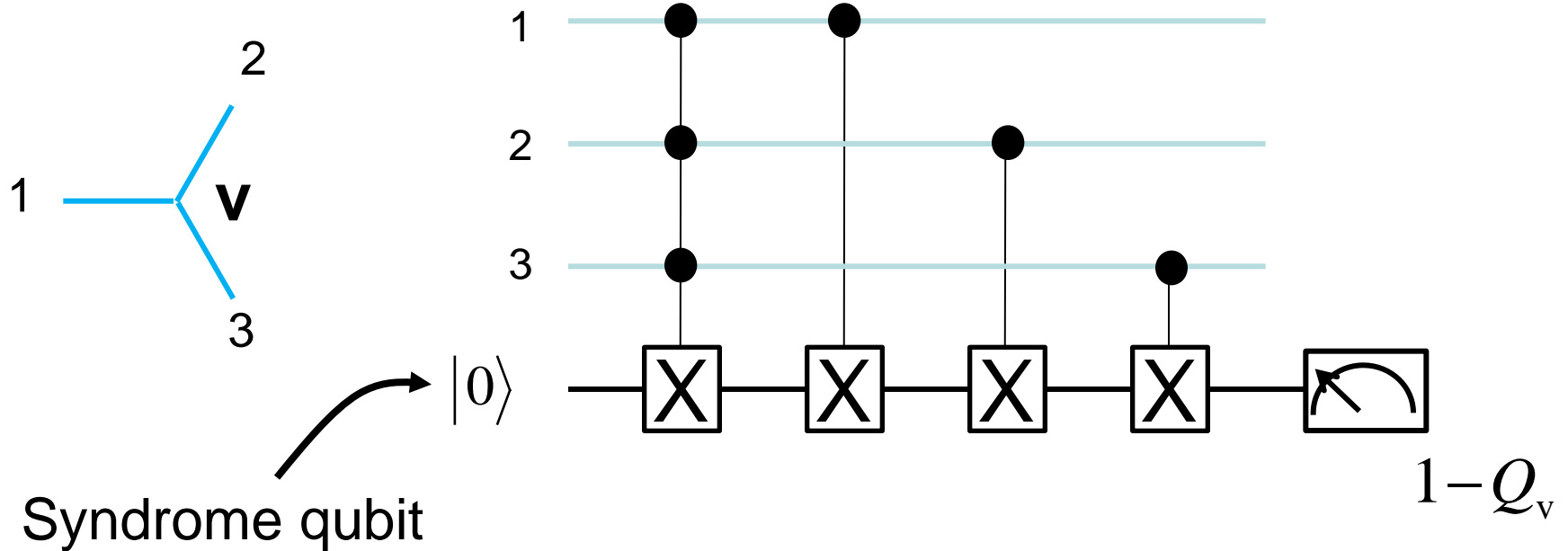


How hard is it to measure Q_v and B_p with a quantum computer?

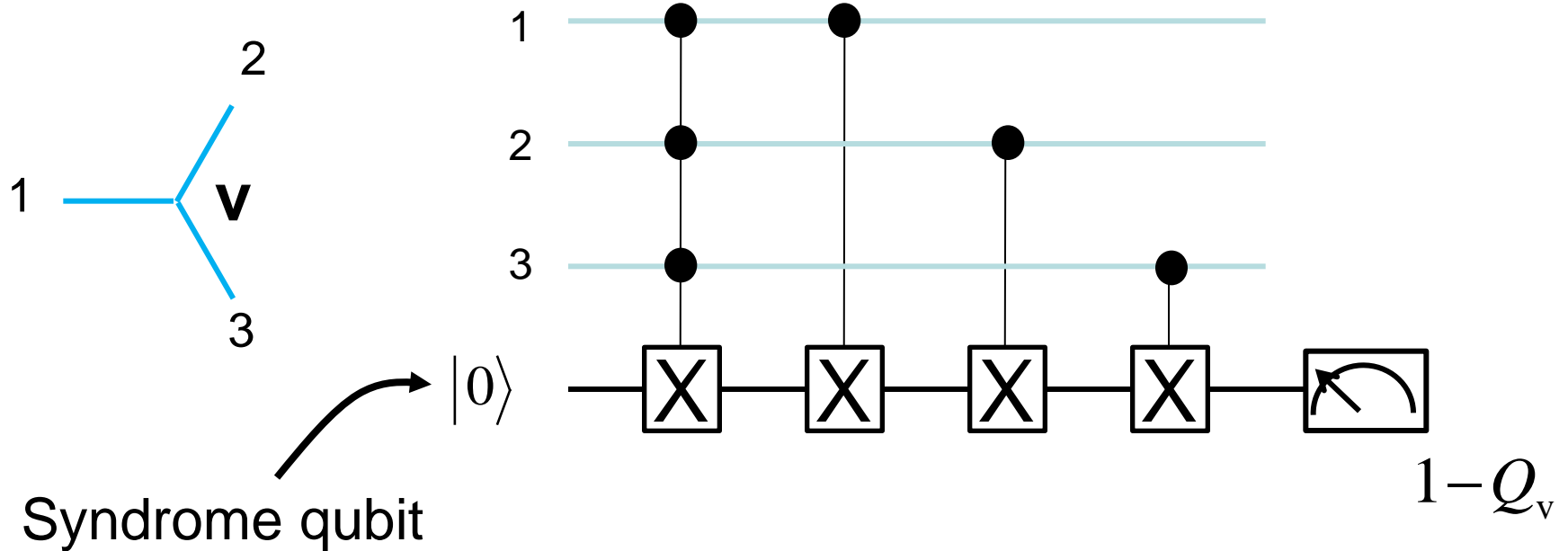
Quantum Circuit for Measuring Q_v



Quantum Circuit for Measuring Q_v

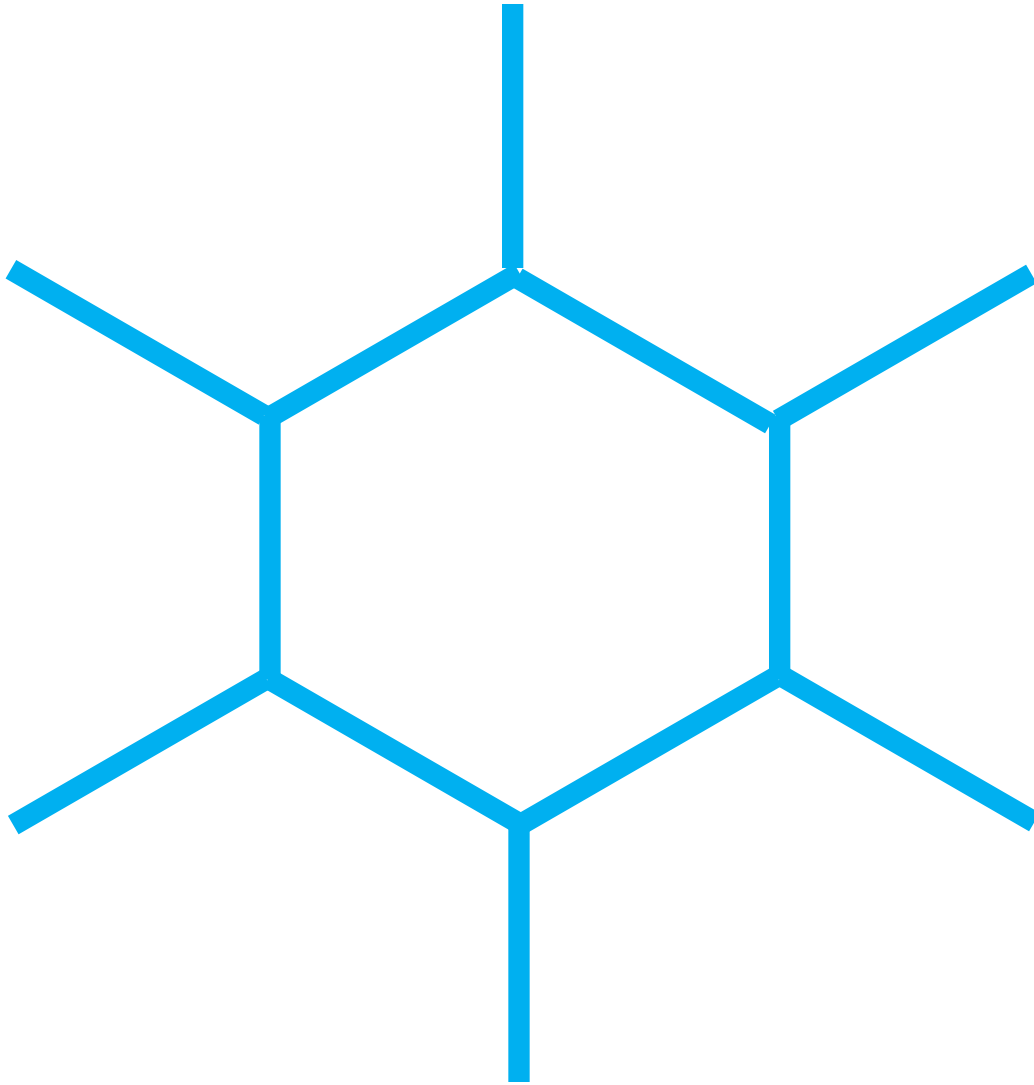


Quantum Circuit for Measuring Q_v

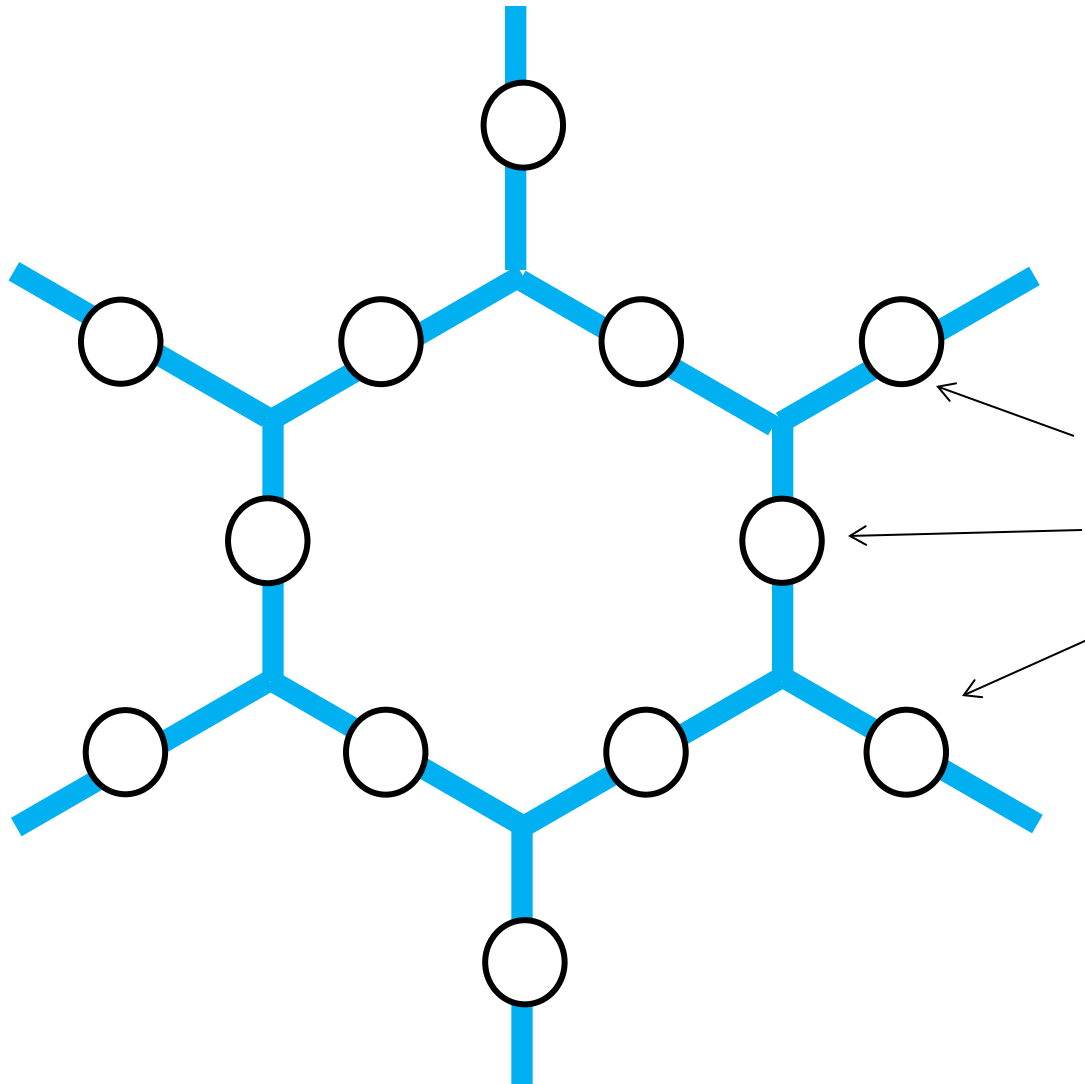


That was easy! What about B_p ?

6-sided Plaquette: B_p is a 12-qubit Operator



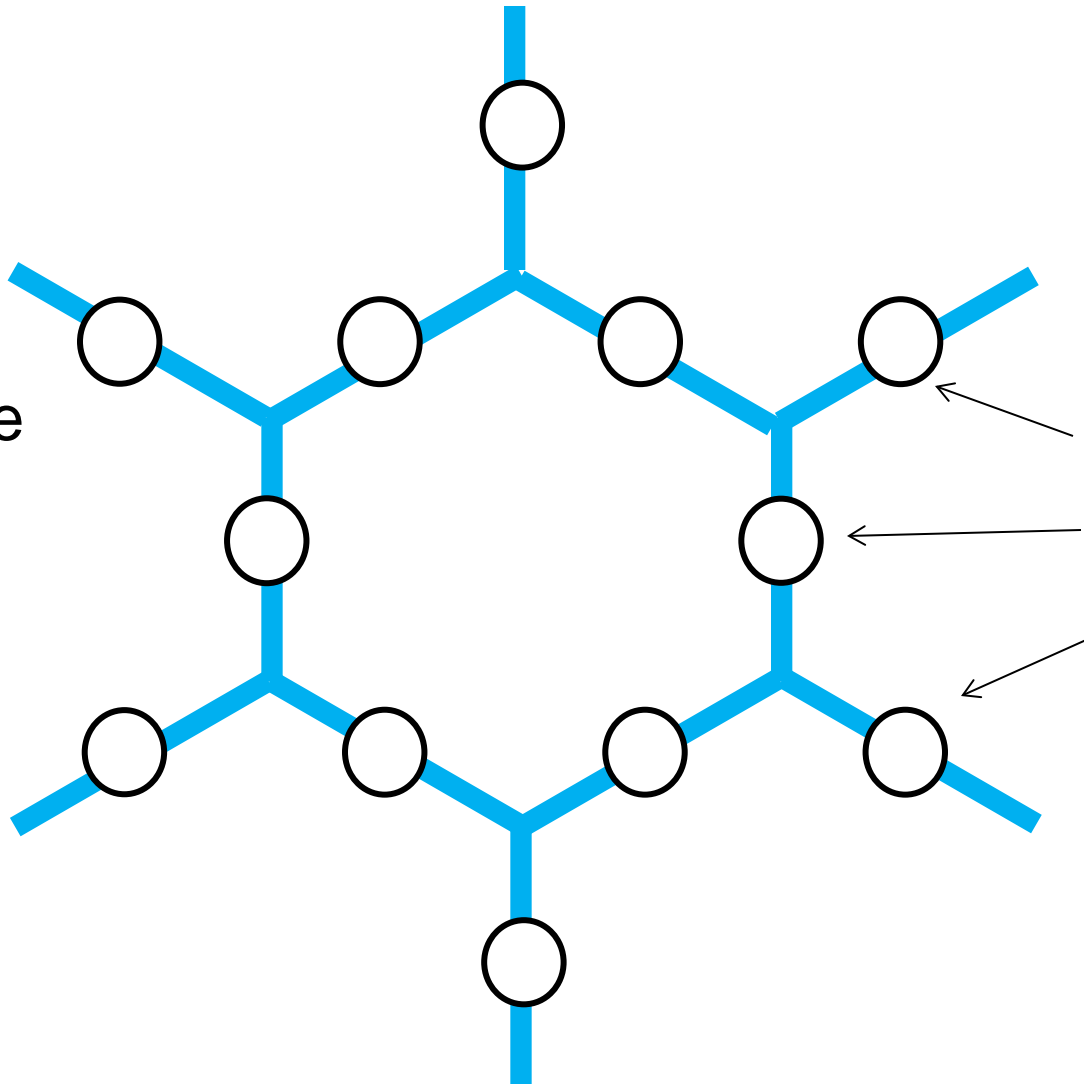
6-sided Plaquette: B_p is a 12-qubit Operator



Physical
qubits are
fixed in
space

6-sided Plaquette: B_p is a 12-qubit Operator

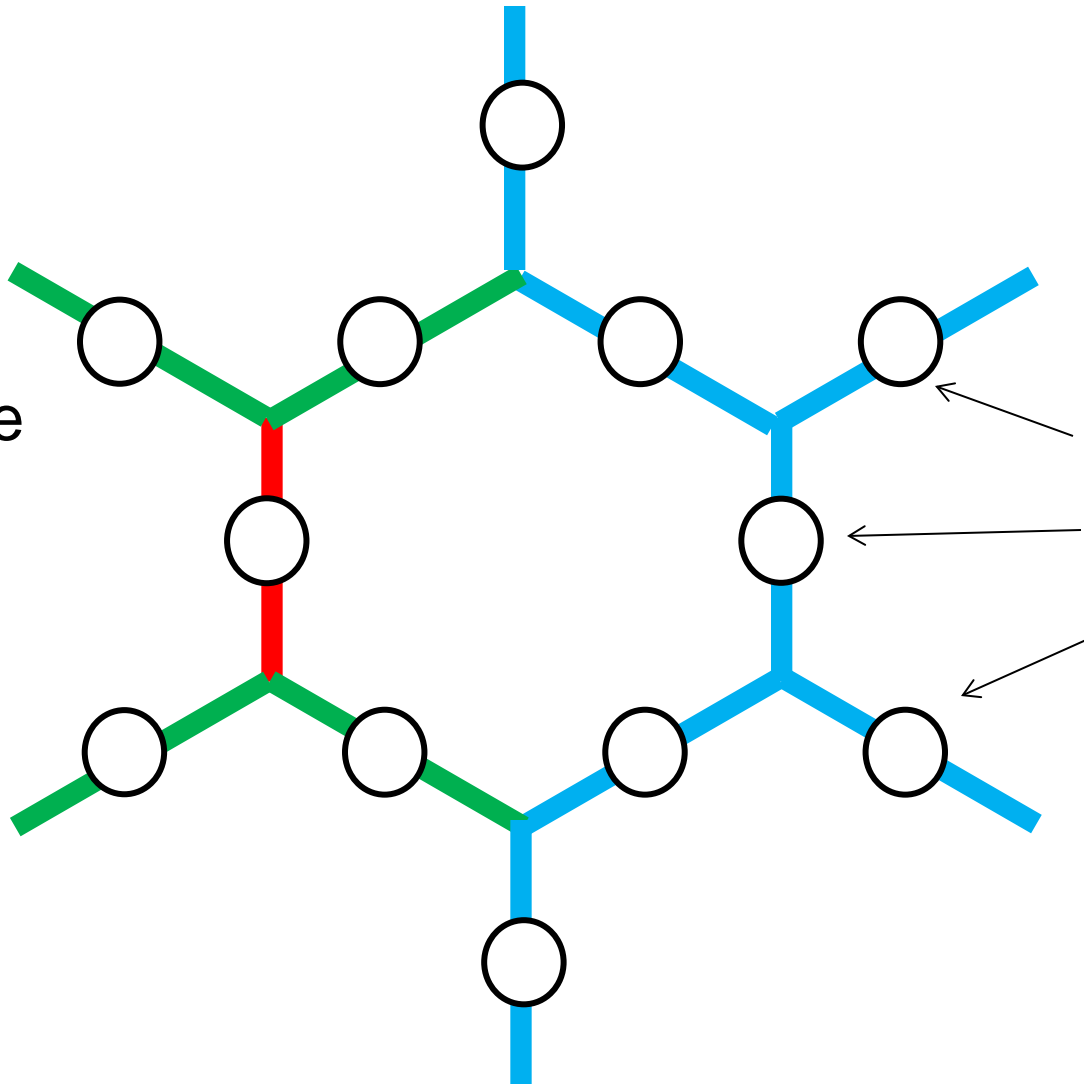
Trivalent lattice
is abstract
and can be
“redrawn”



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6-sided Plaquette: B_p is a 12-qubit Operator

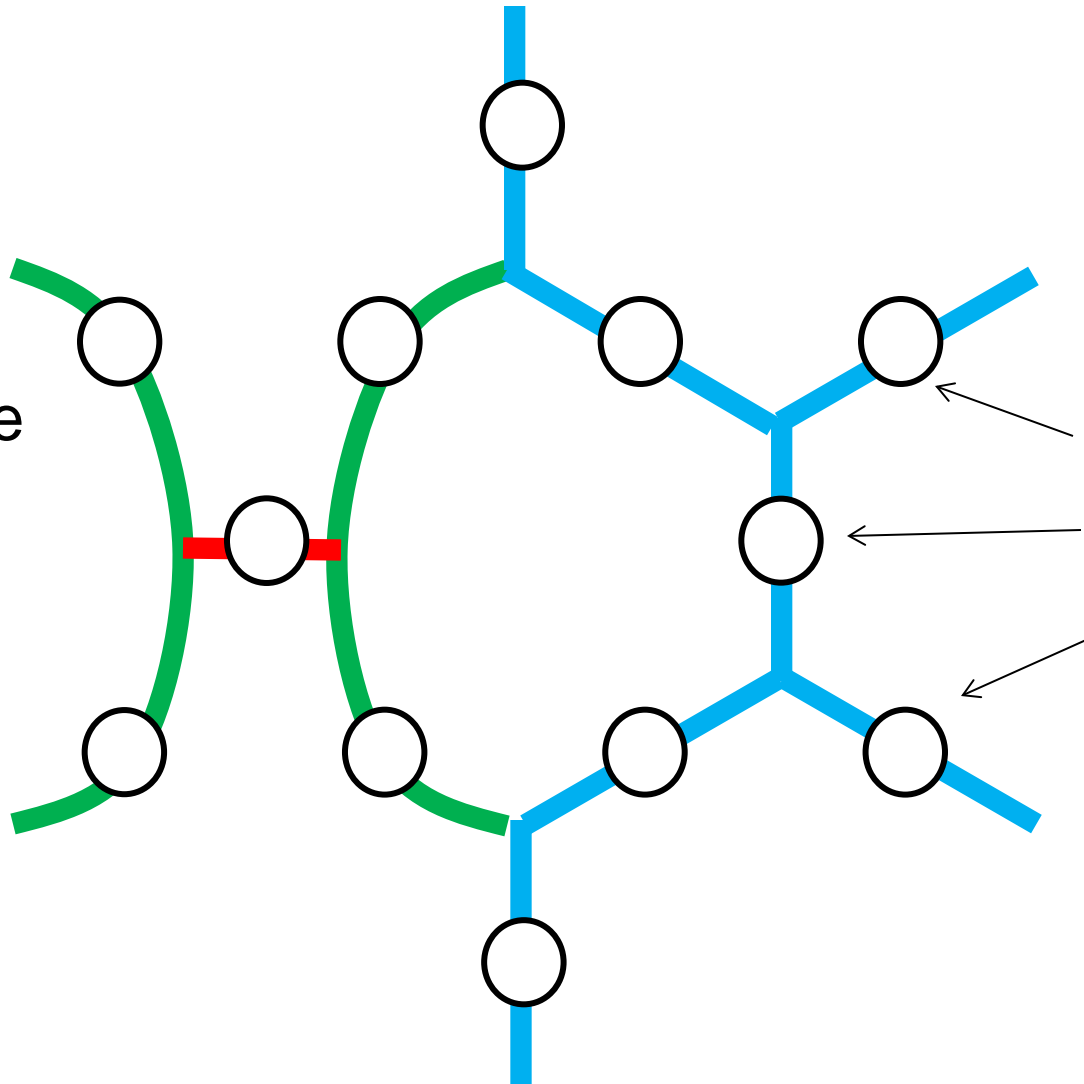
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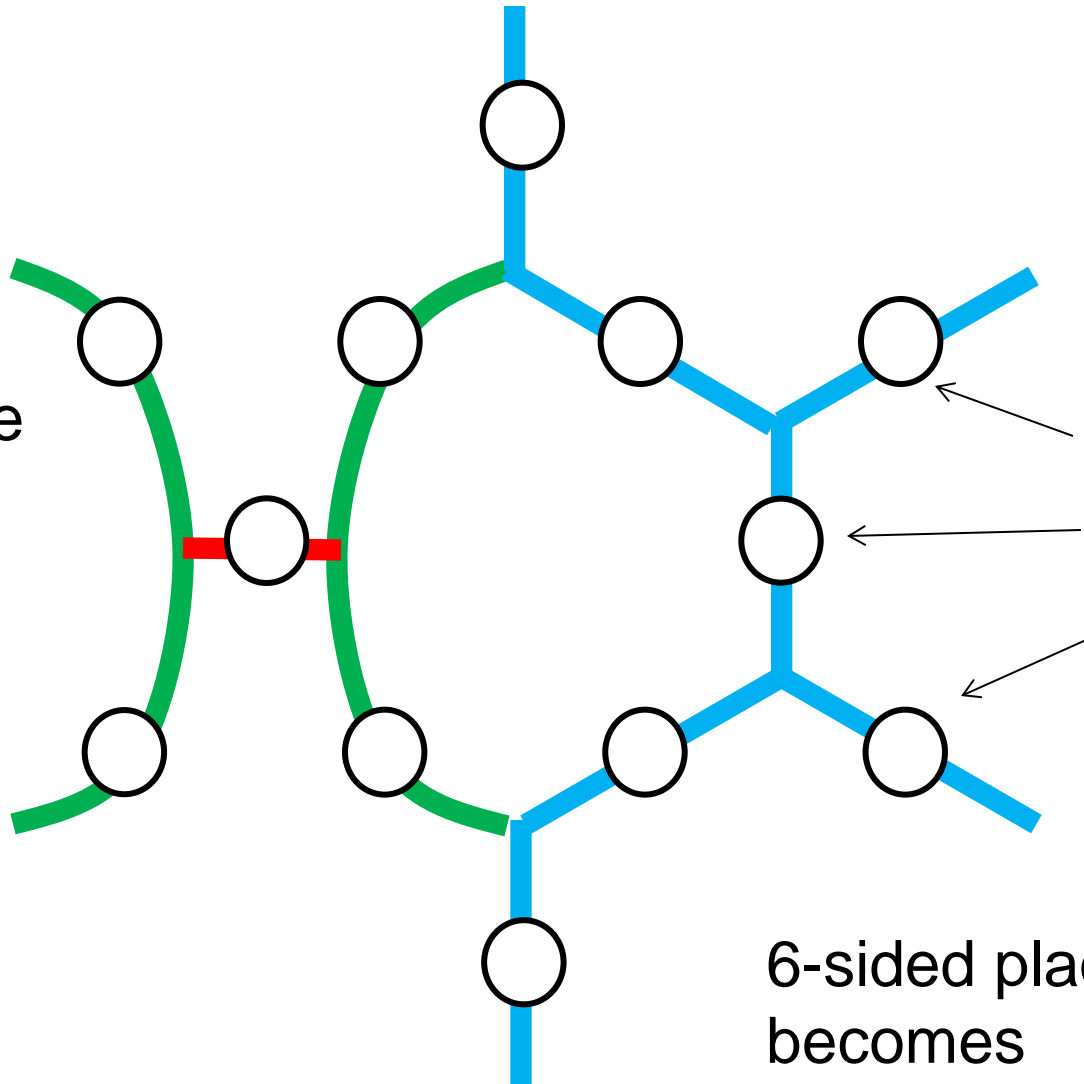
Trivalent lattice
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6-sided Plaquette: B_p is a 12-qubit Operator

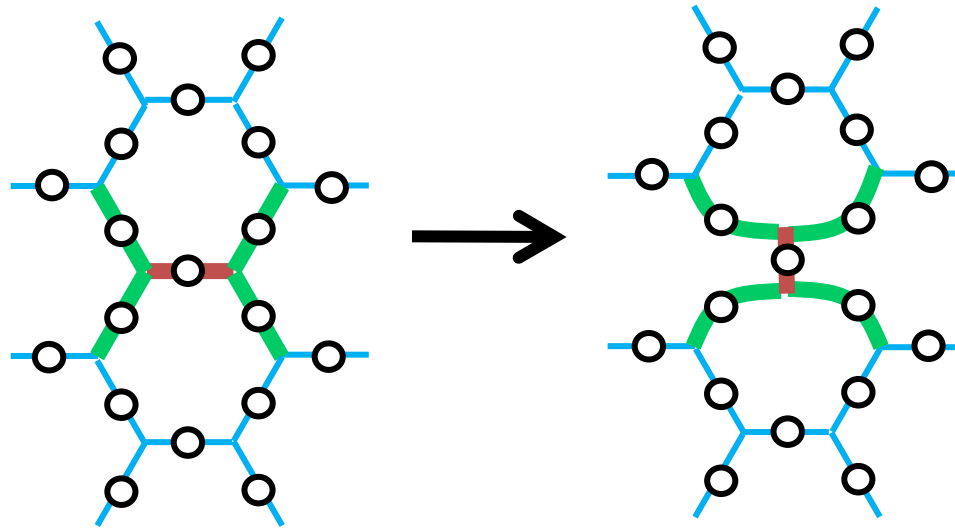
Trivalent lattice
is abstract
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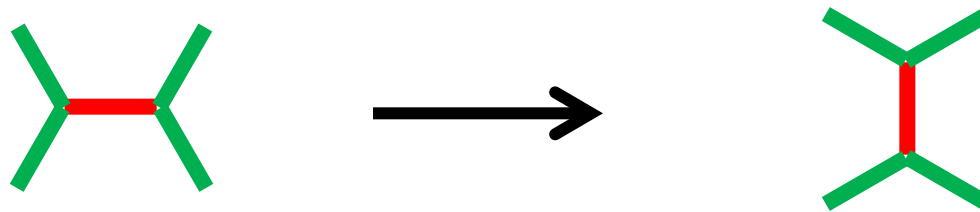
Physical
qubits are
fixed in
space

6-sided plaquette
becomes
5-sided plaquette!

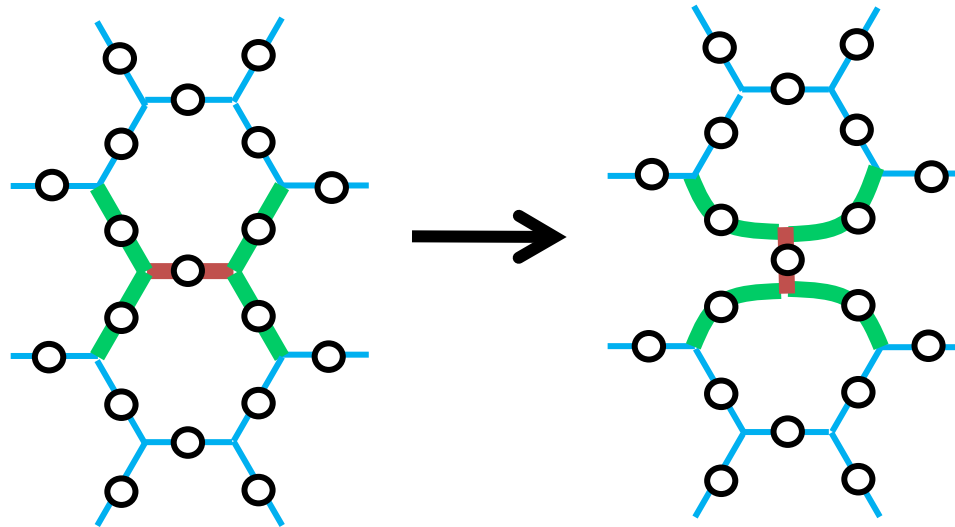
The F-Move



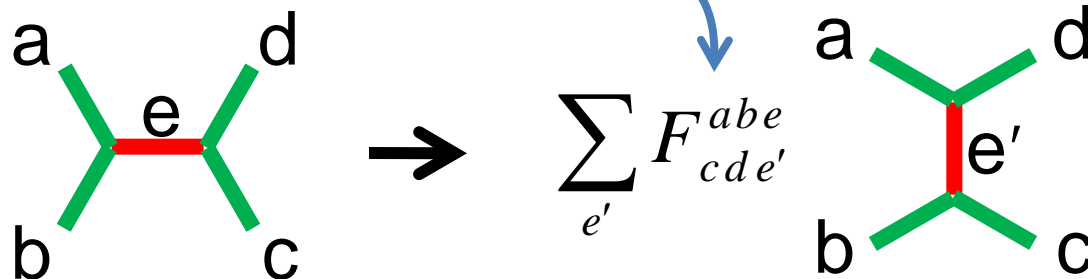
Locally redraw the lattice five qubits at a time.



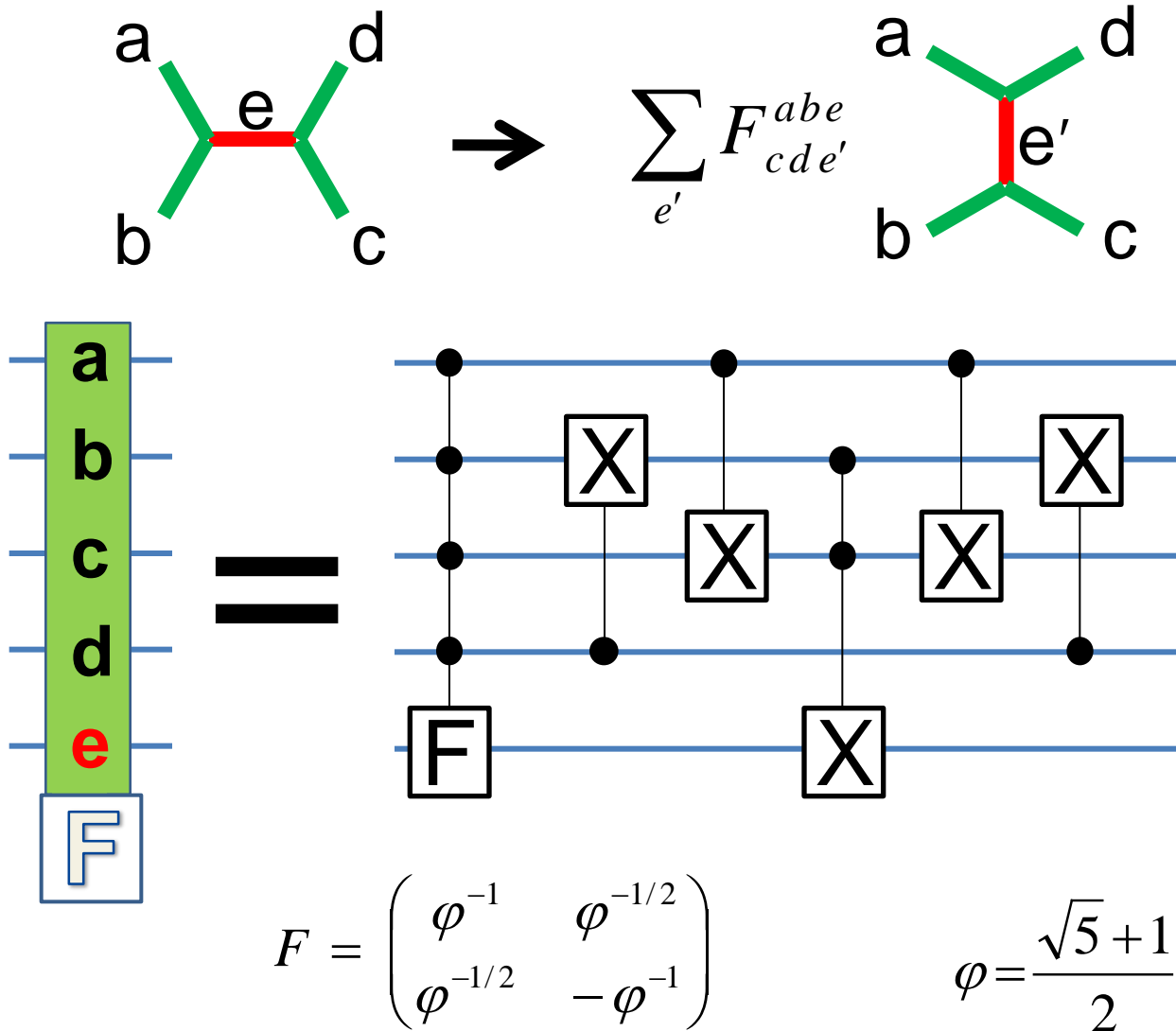
The F-Move



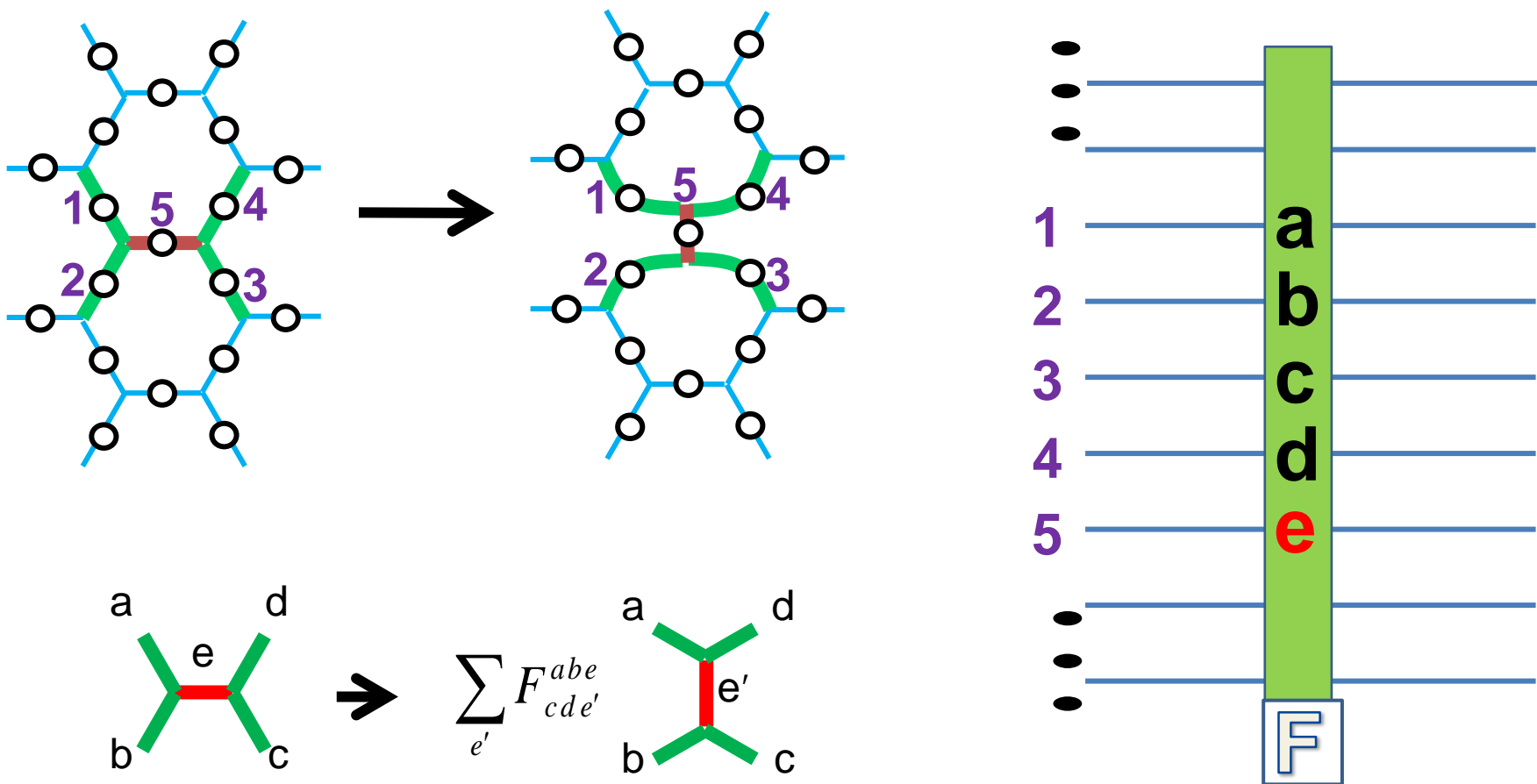
Apply a unitary operation on these five qubits to stay in the Levin-Wen ground state.



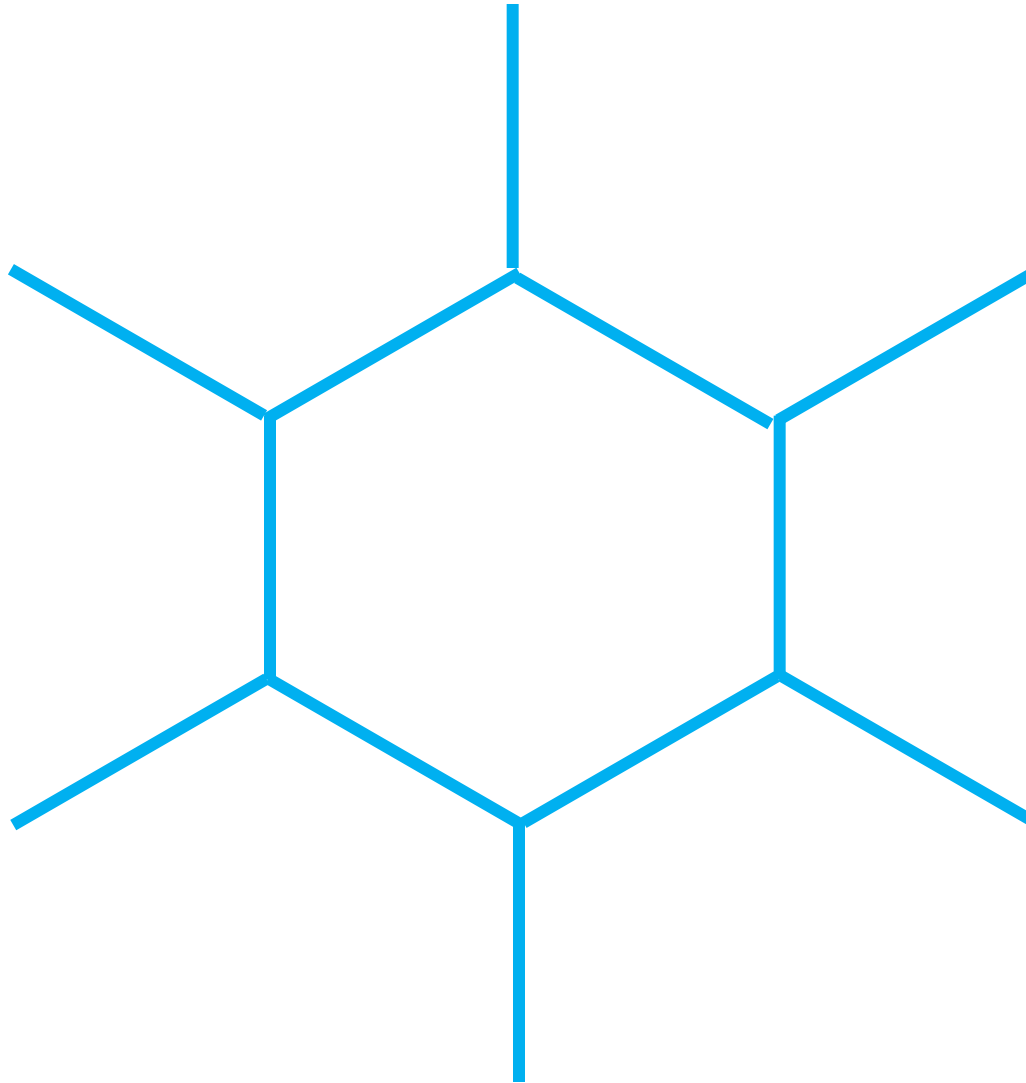
F Quantum Circuit



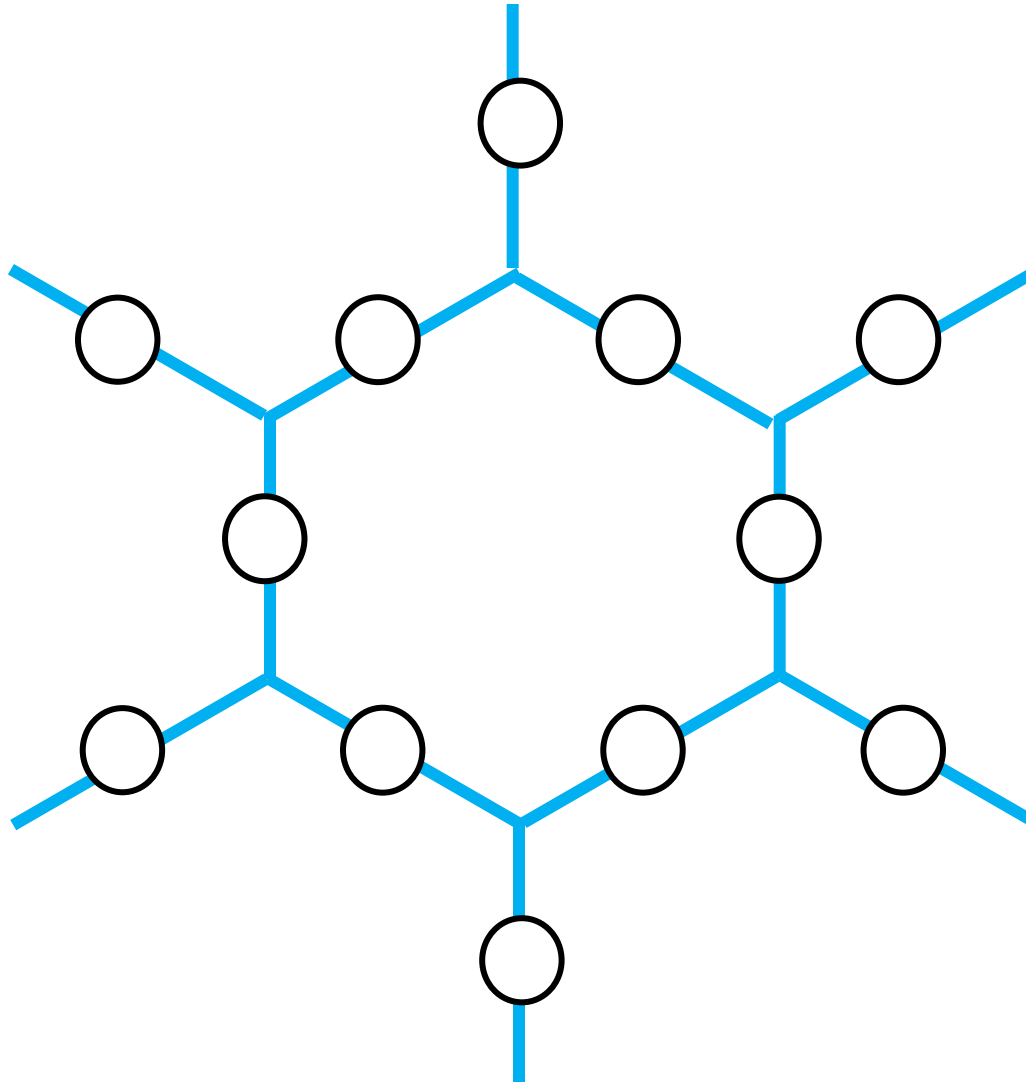
F Quantum Circuit



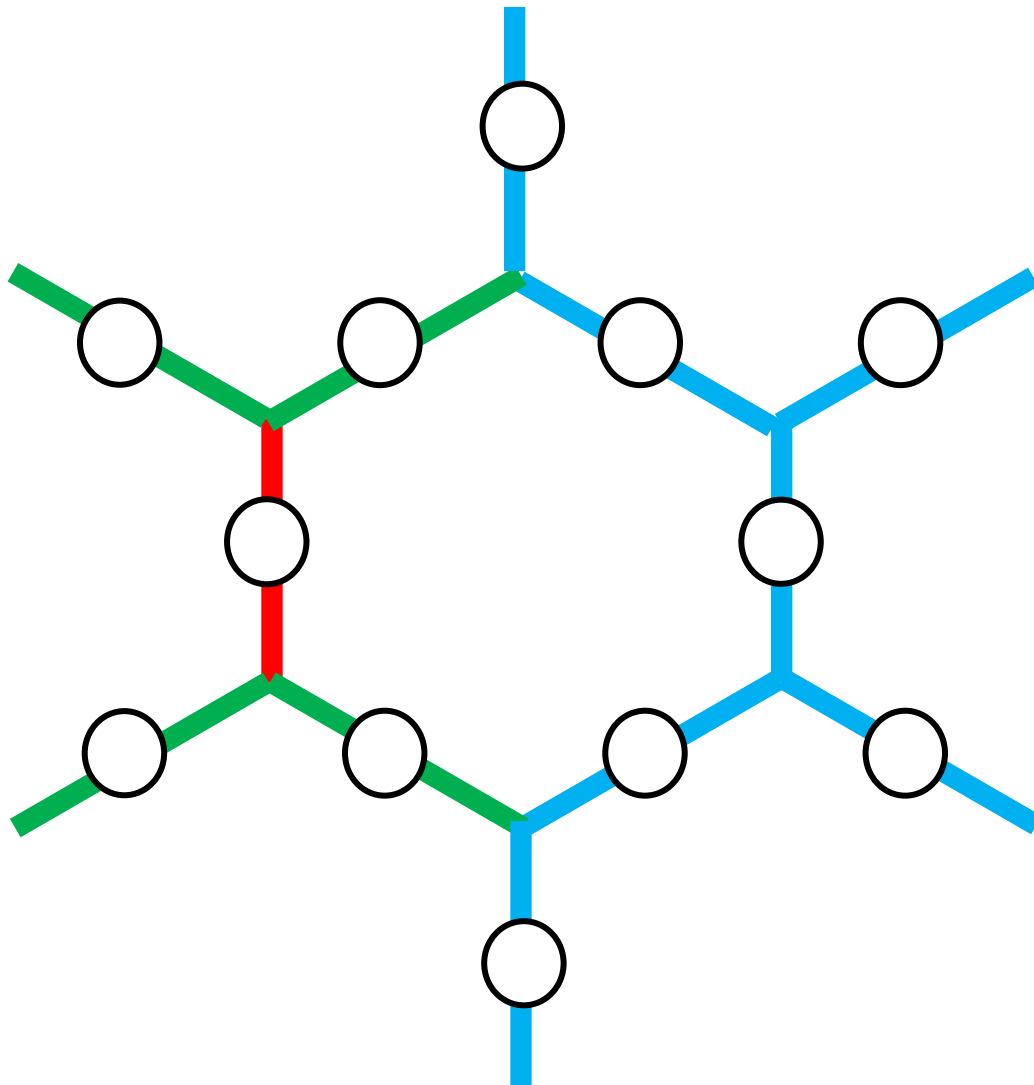
Plaquette Reduction



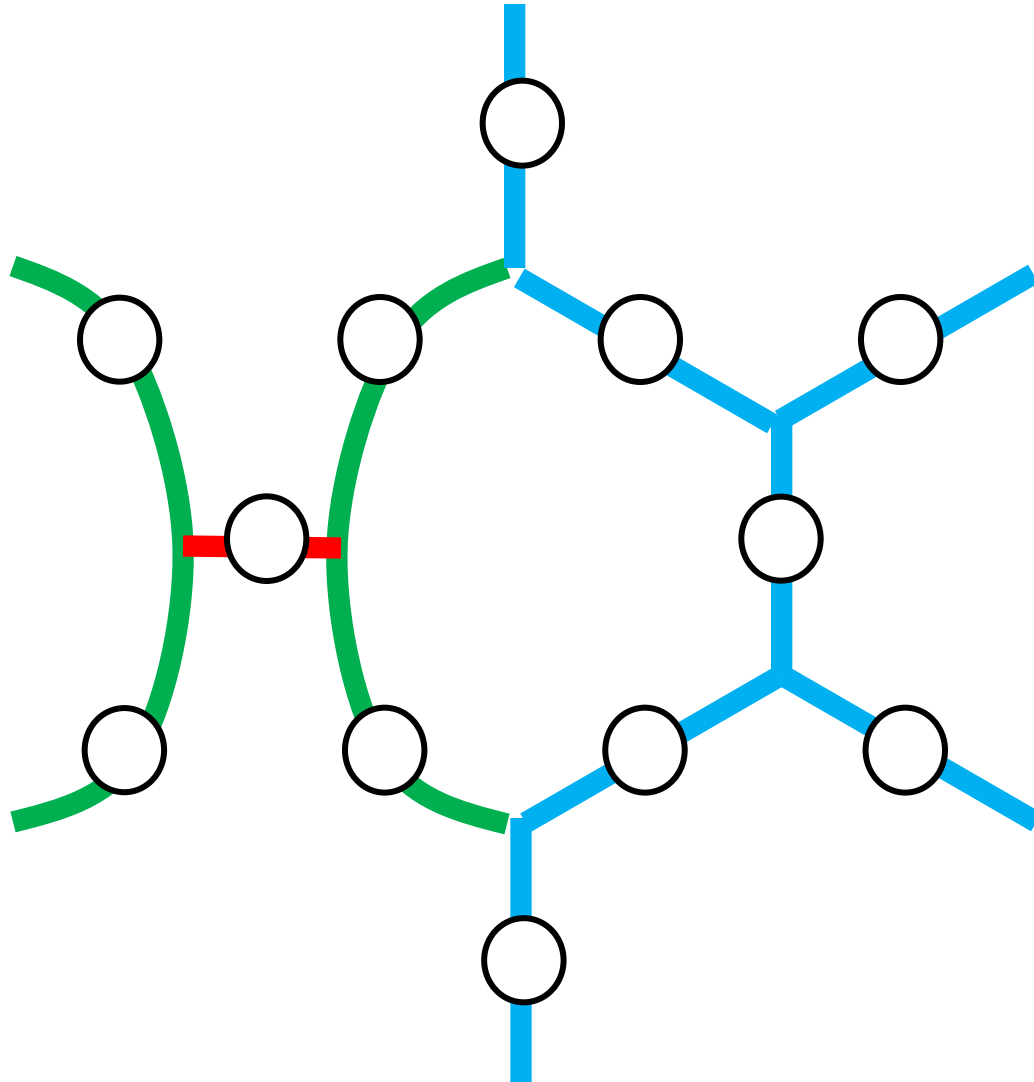
Plaquette Reduction



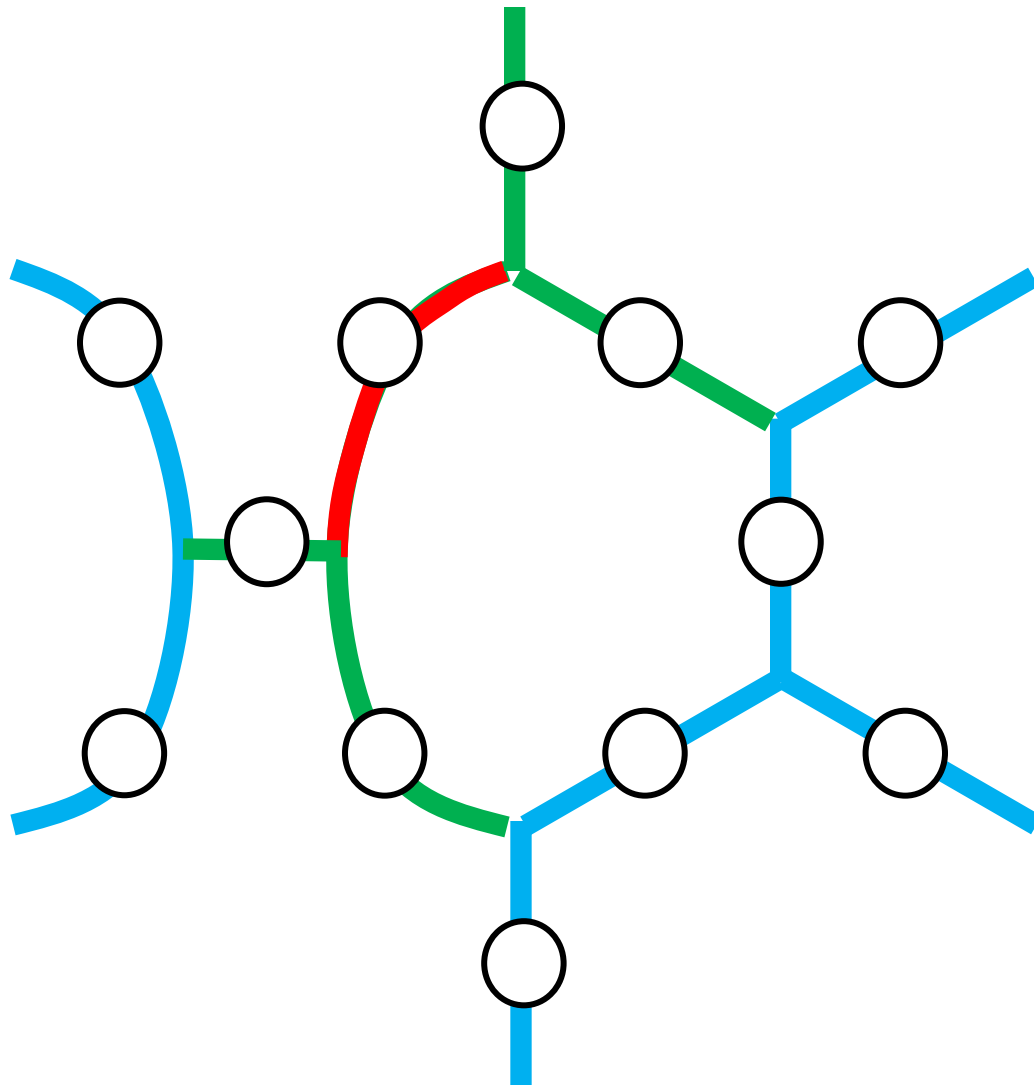
Plaquette Reduction



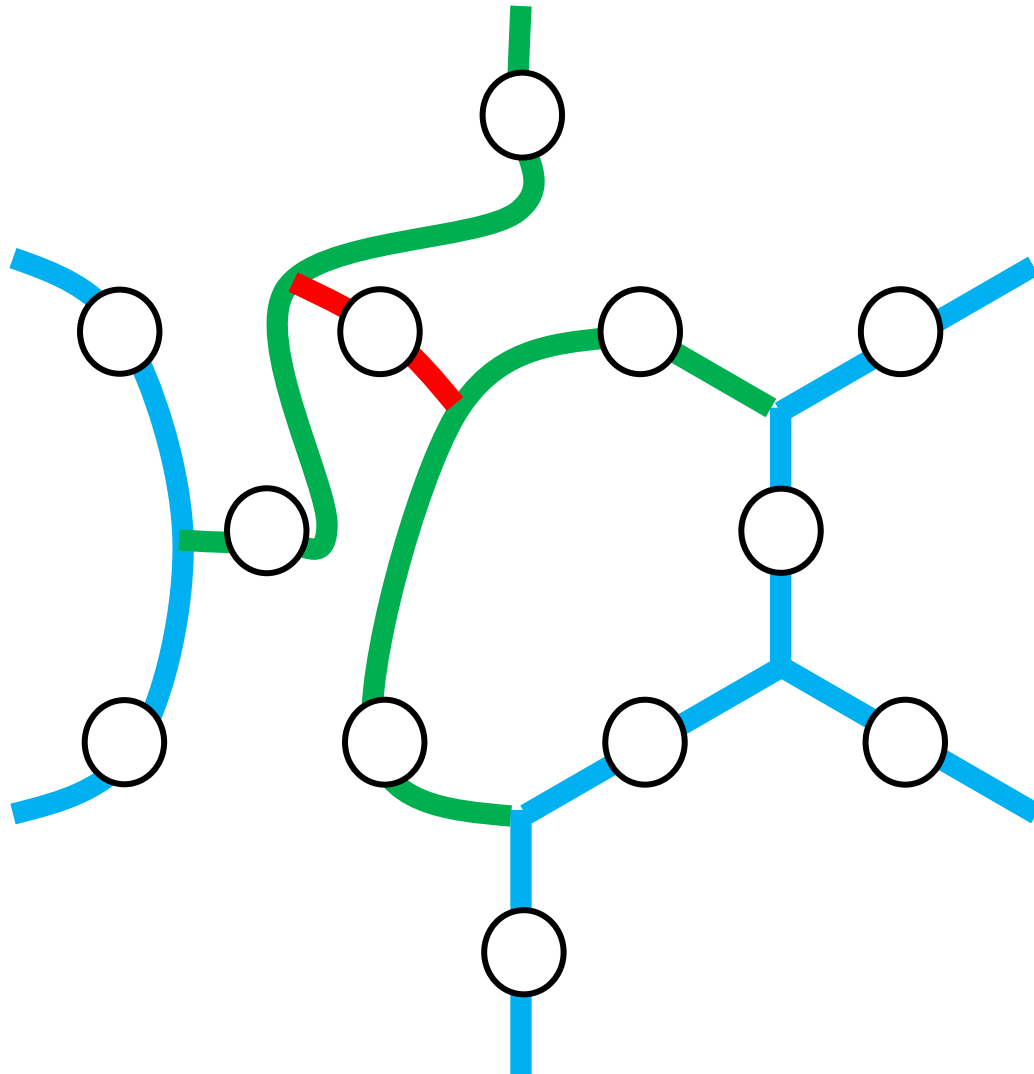
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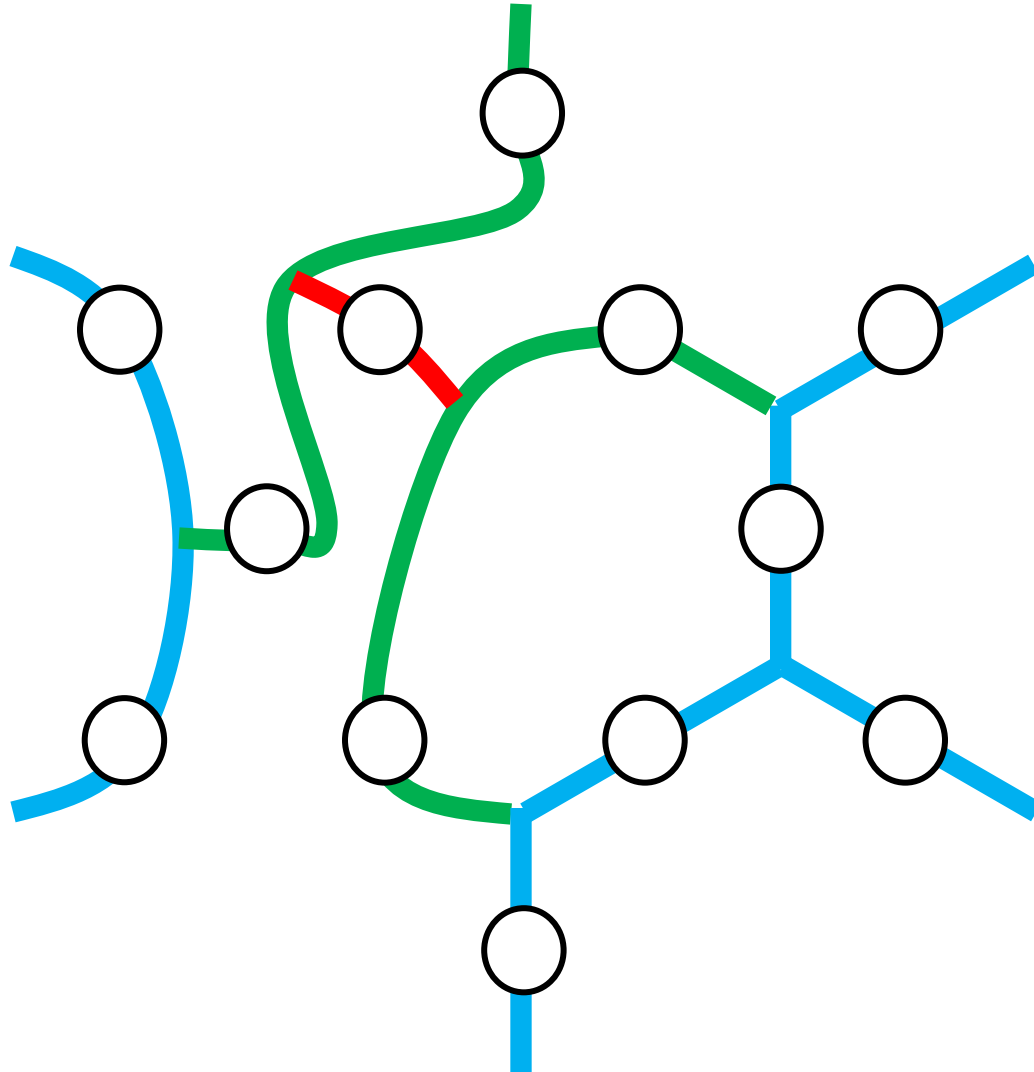
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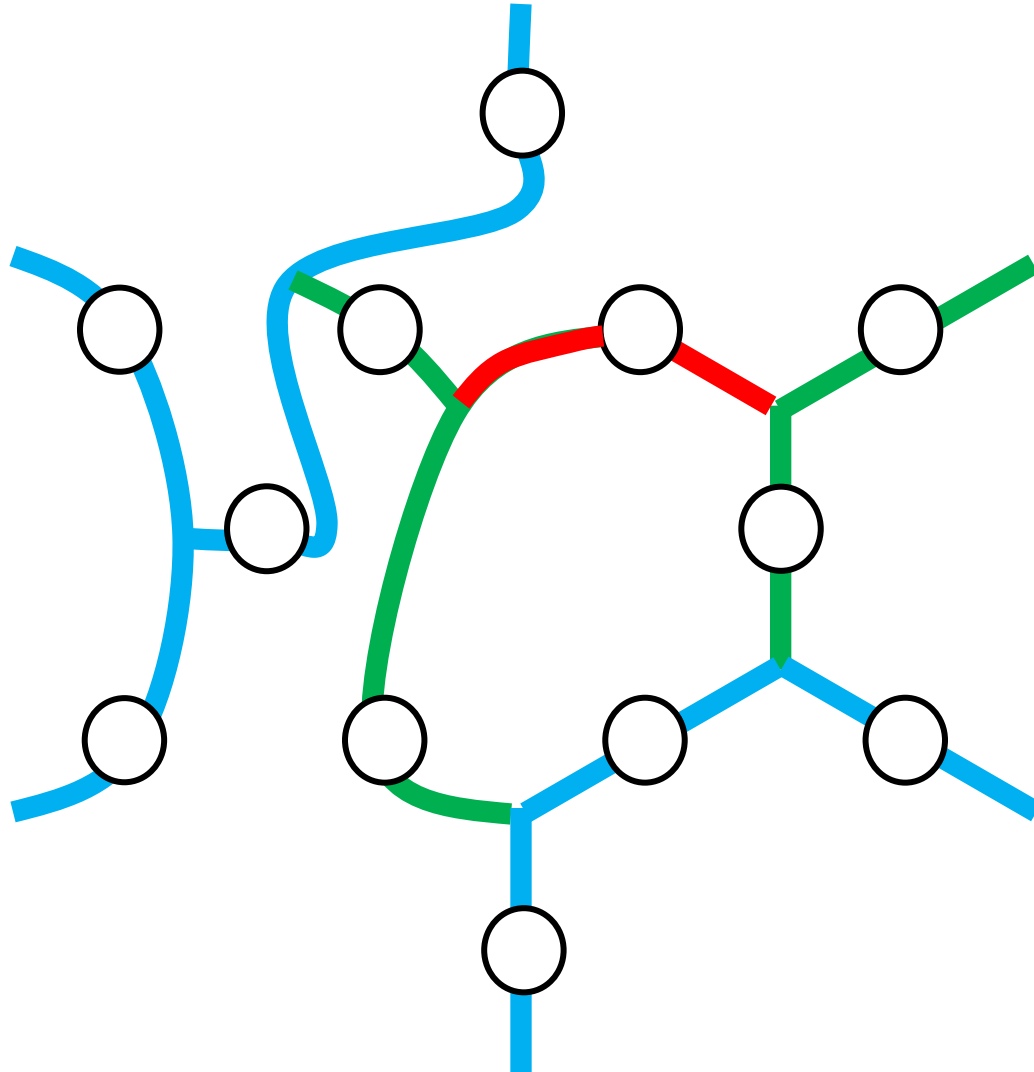
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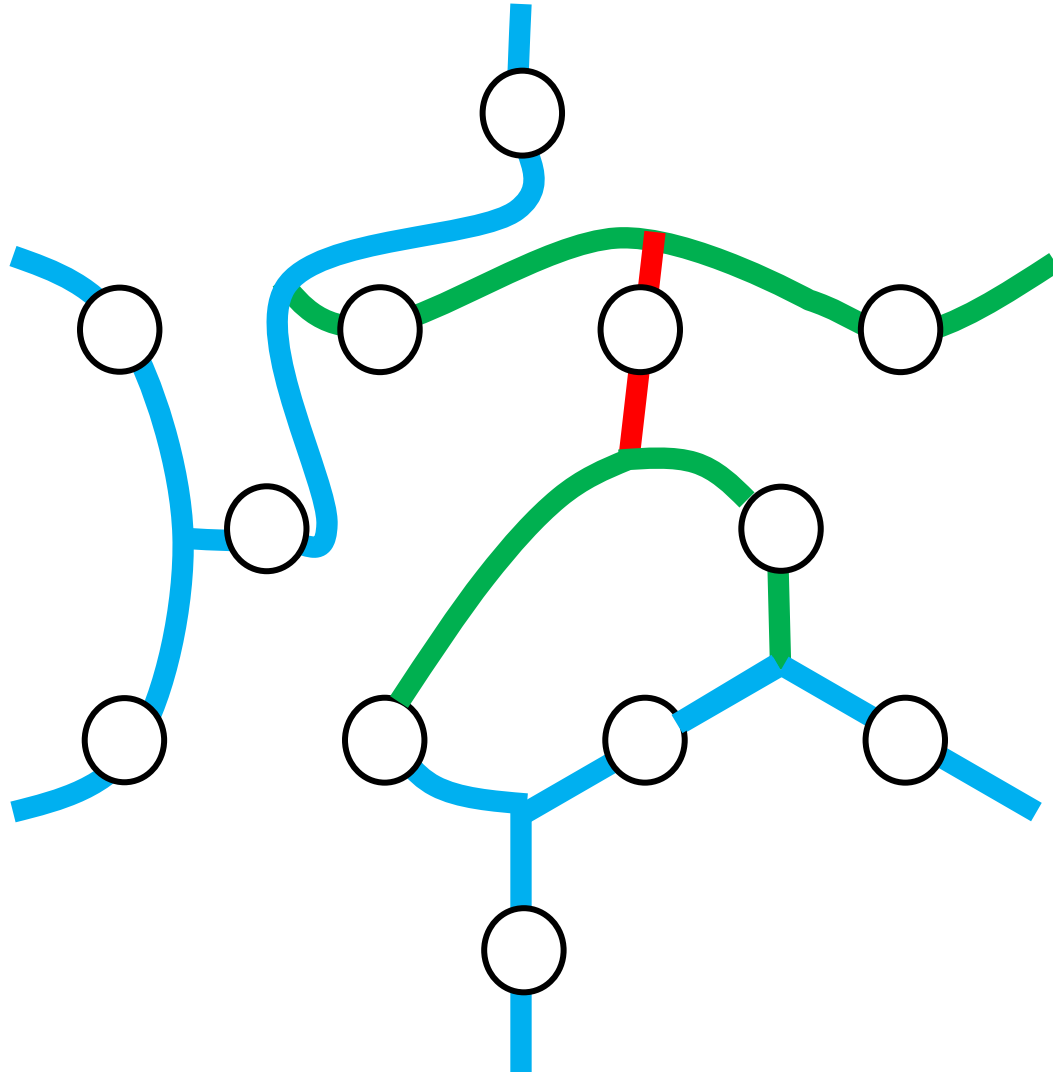
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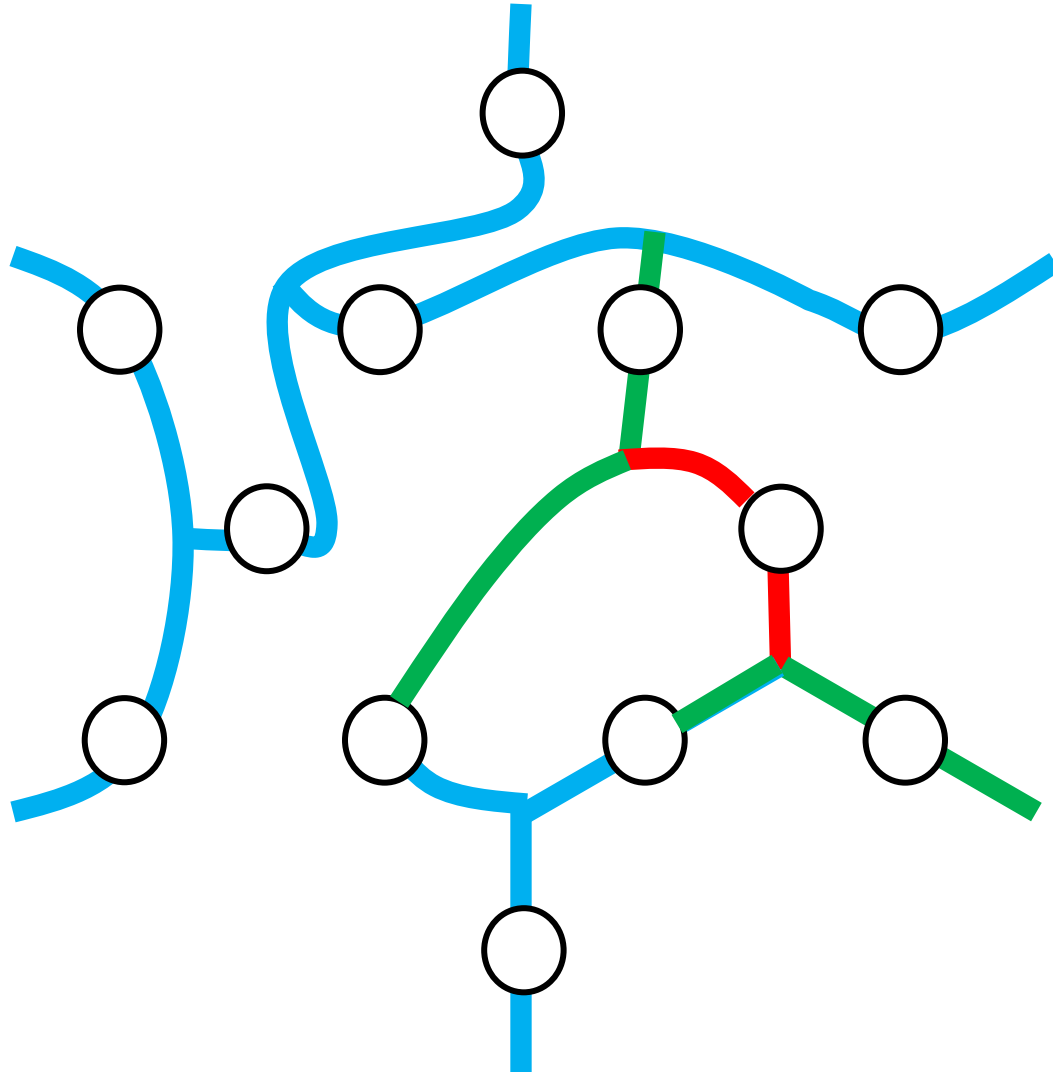
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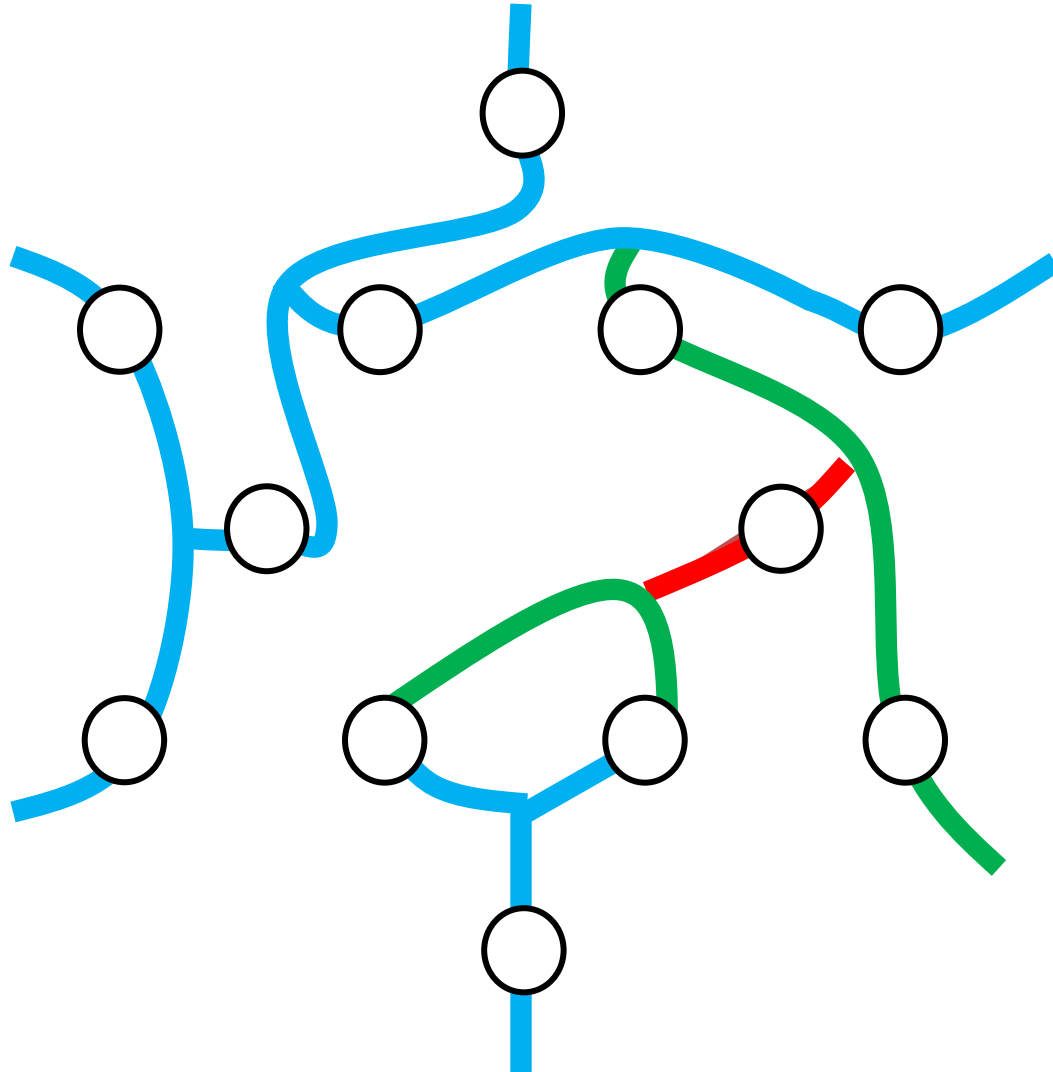
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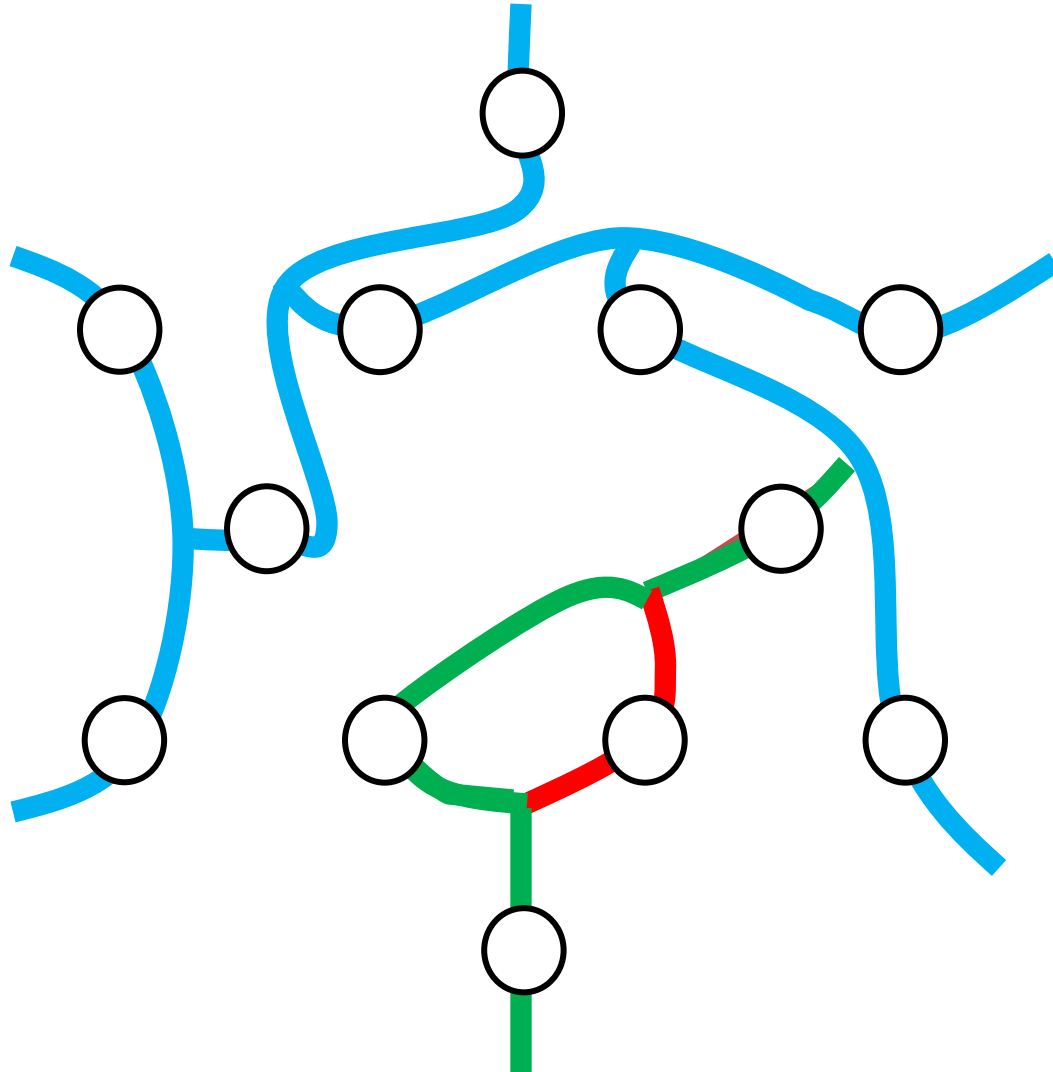
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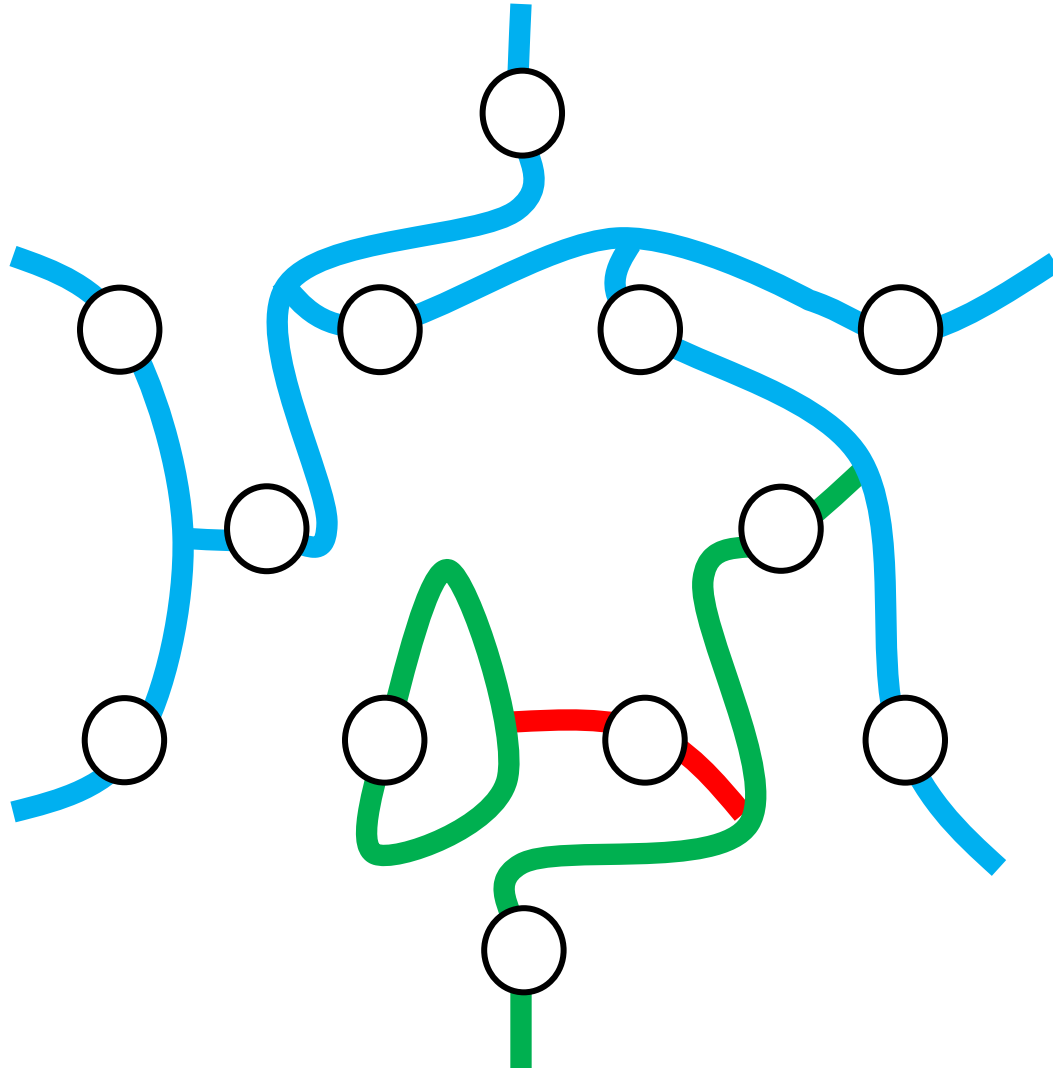
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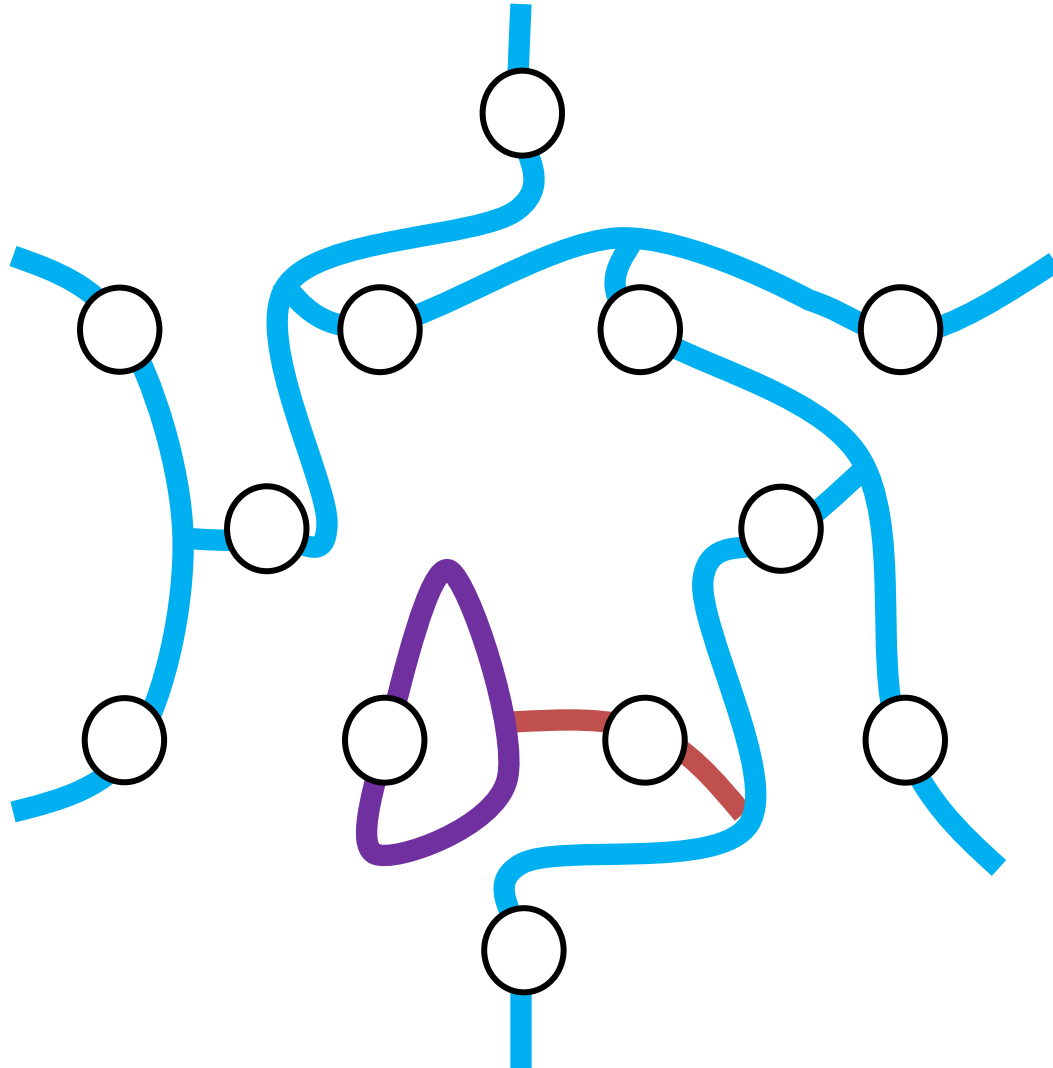
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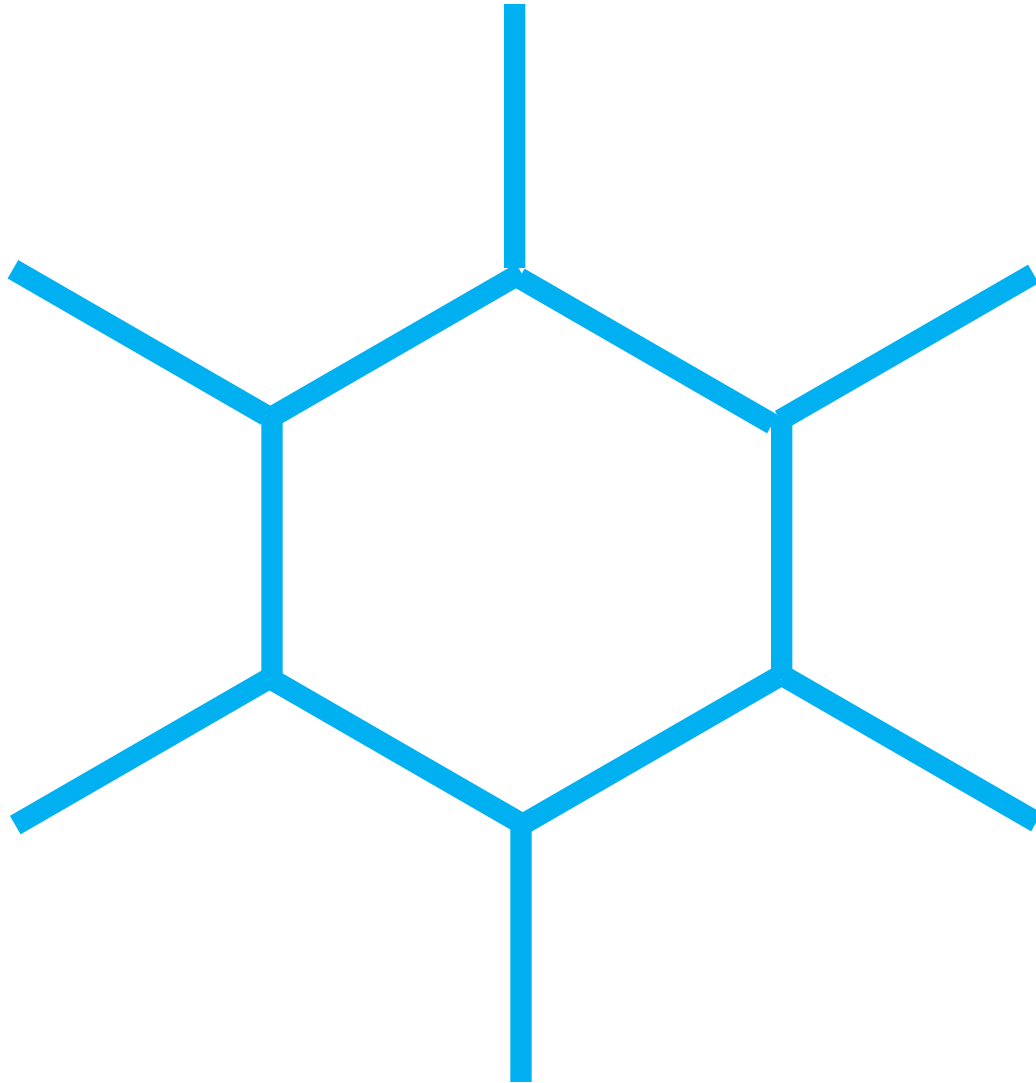
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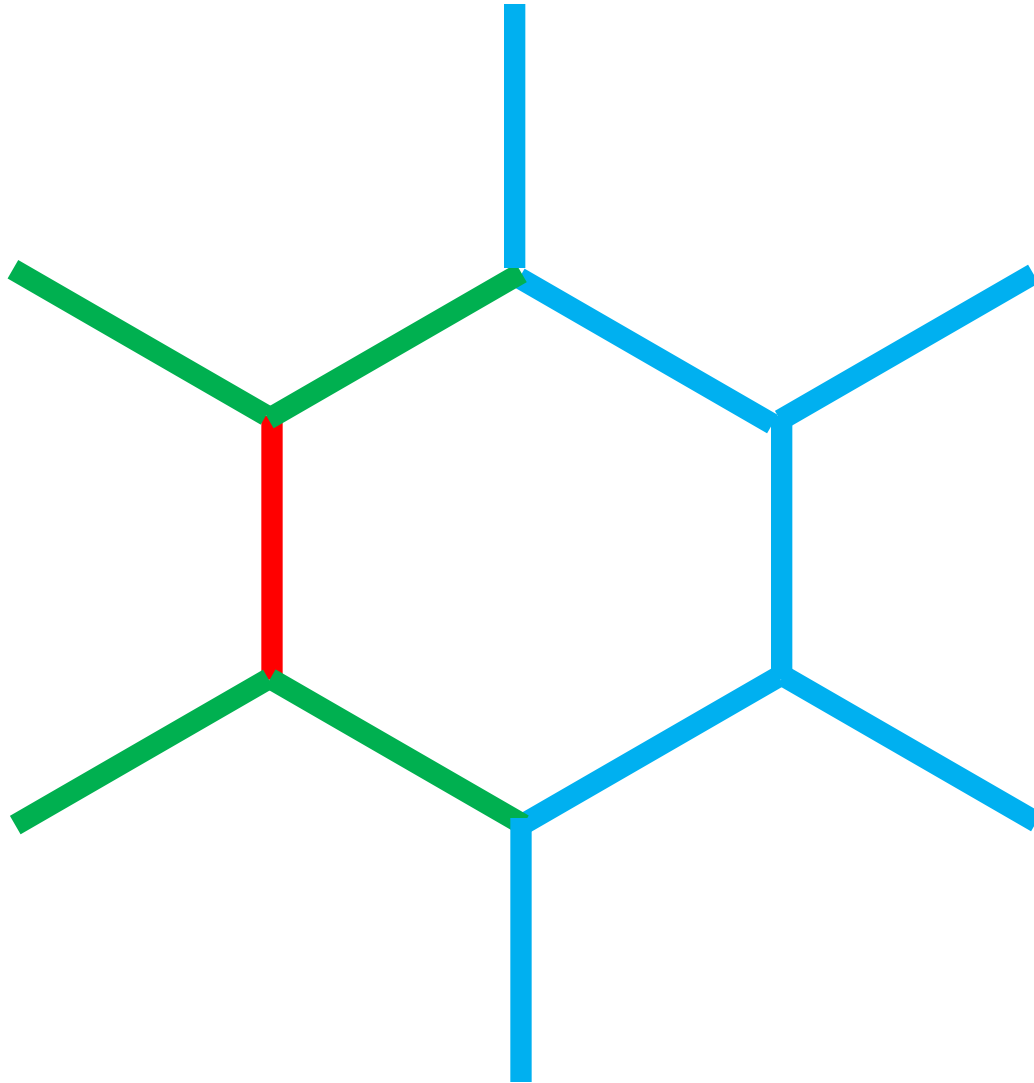
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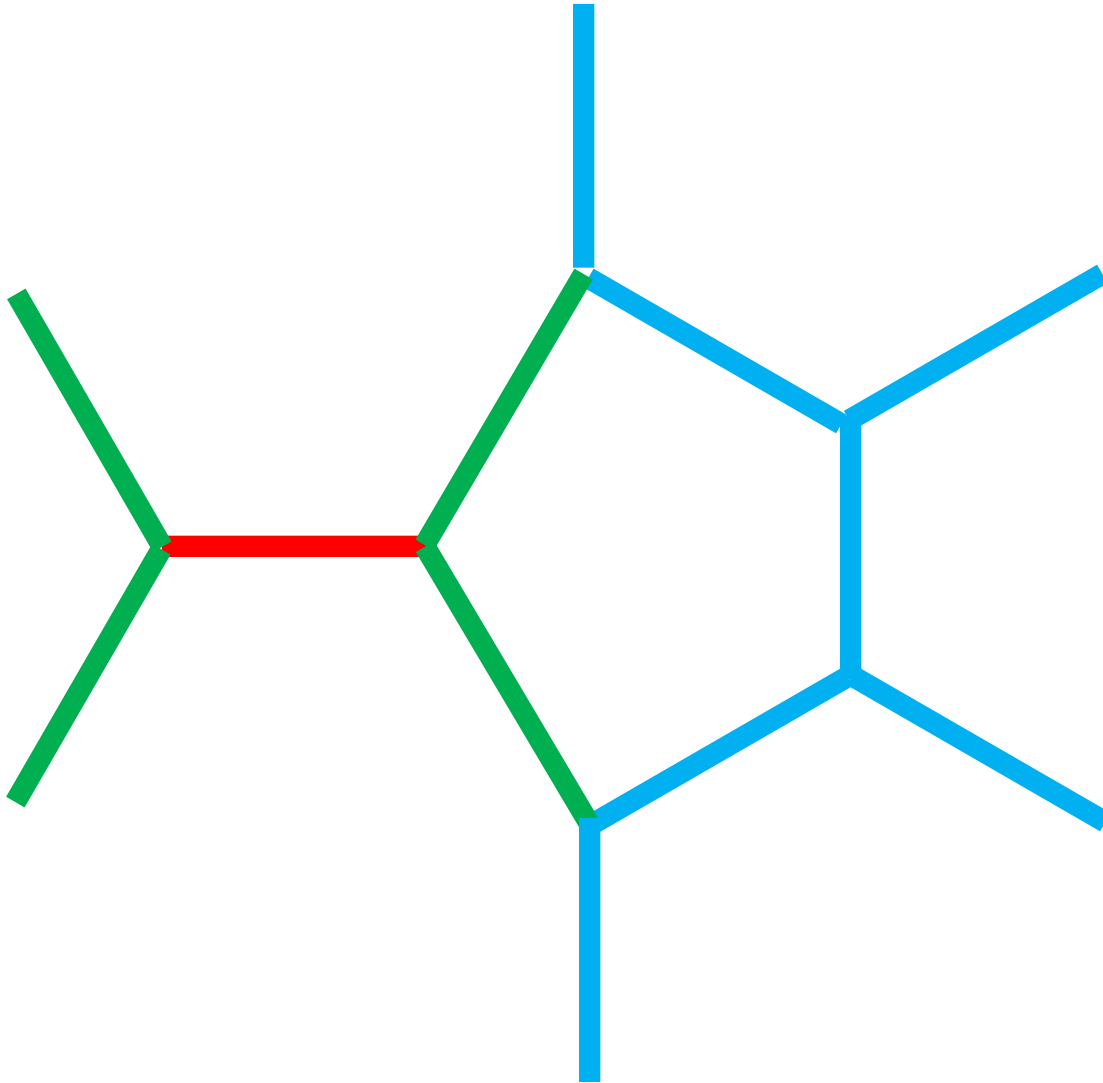
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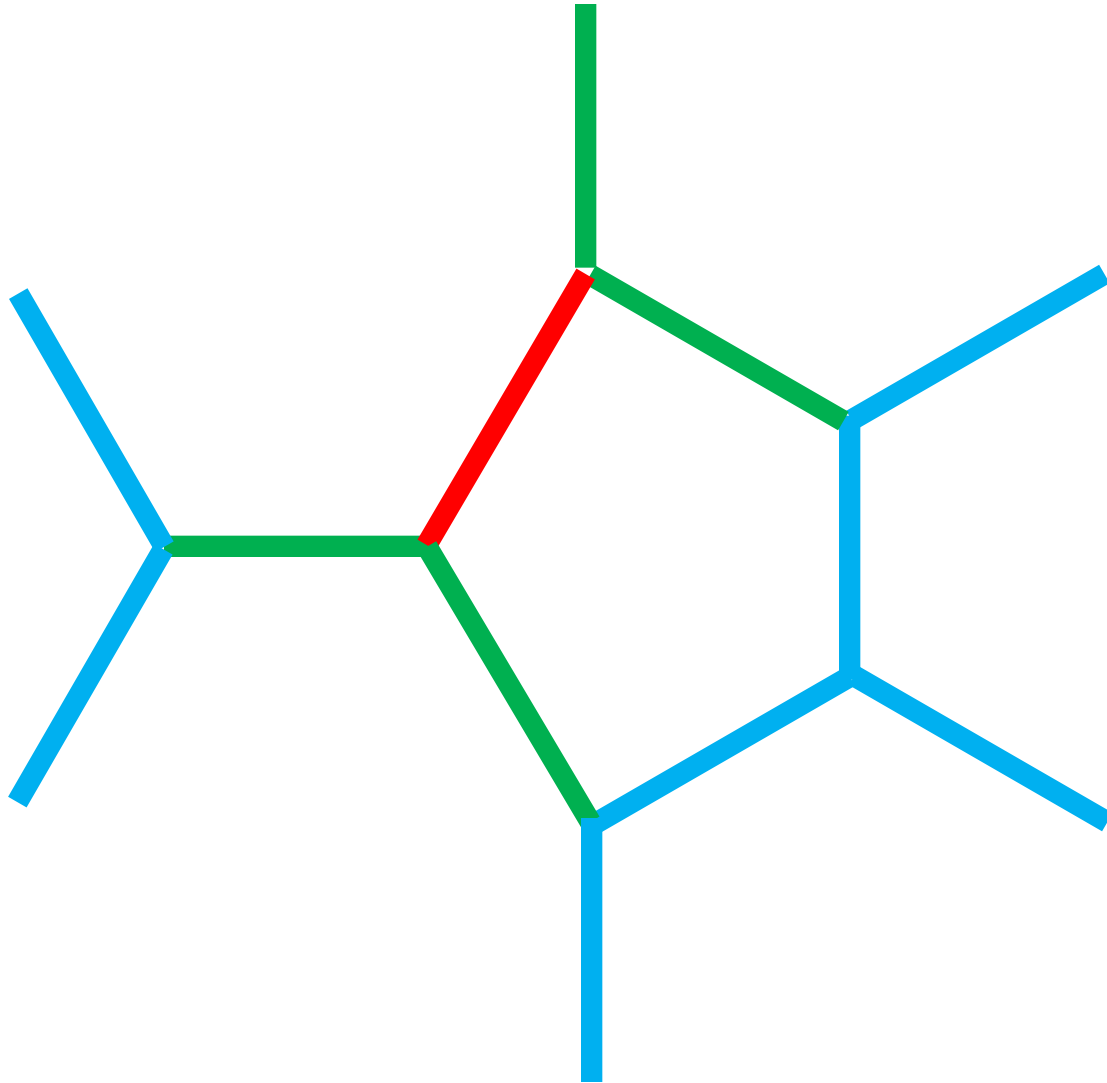
Plaquette Reduction



Plaquette Reduction



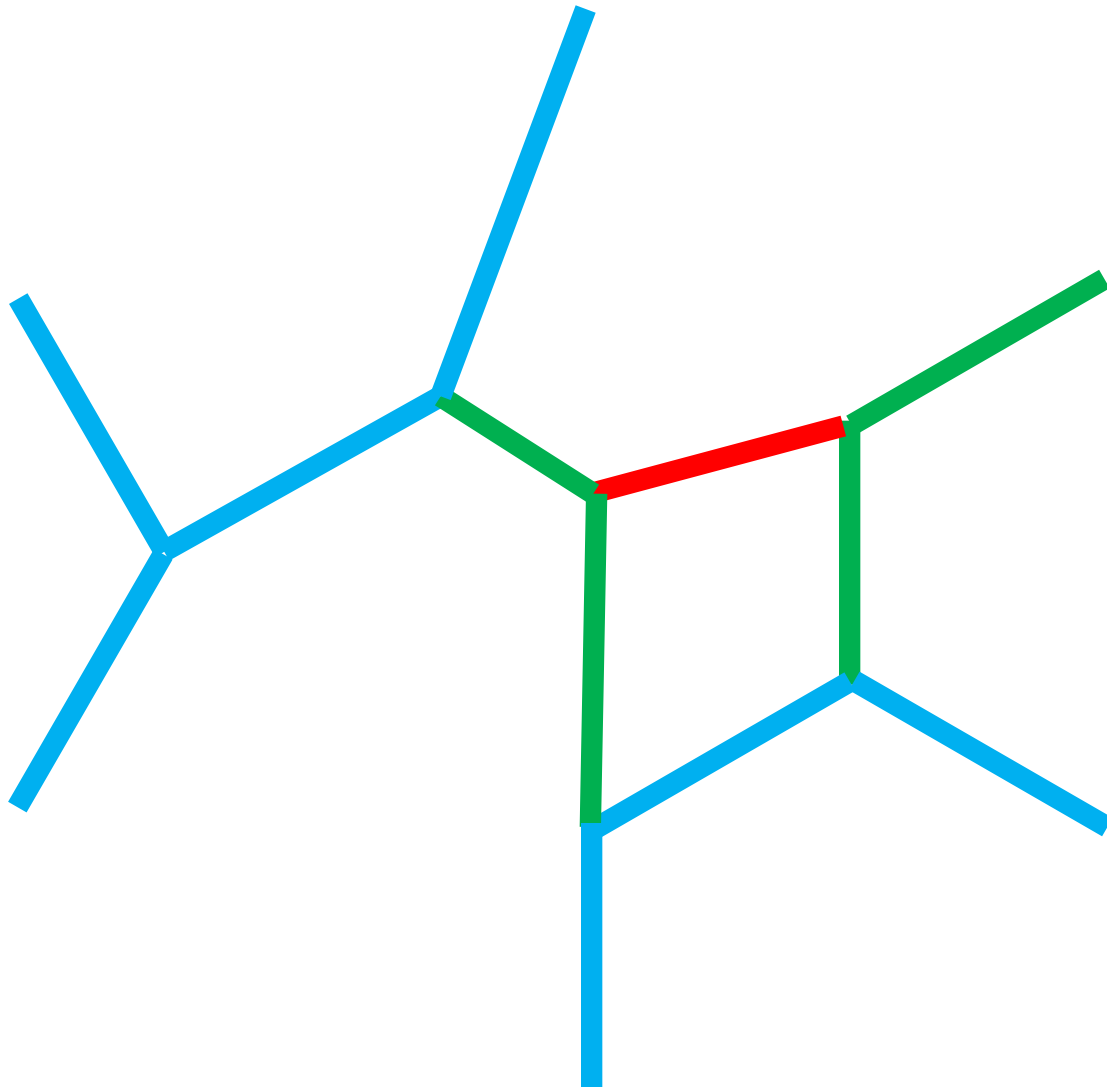
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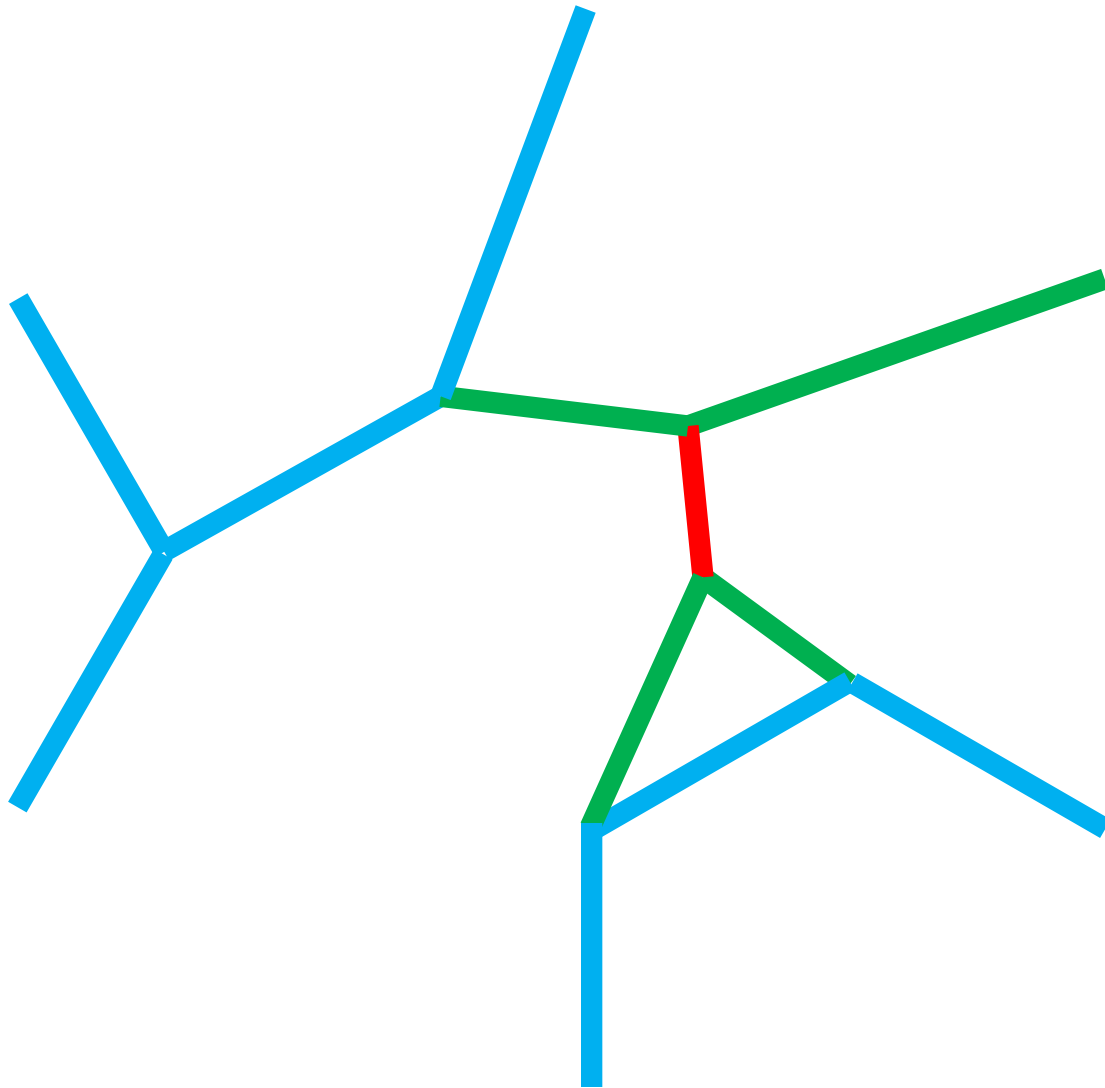
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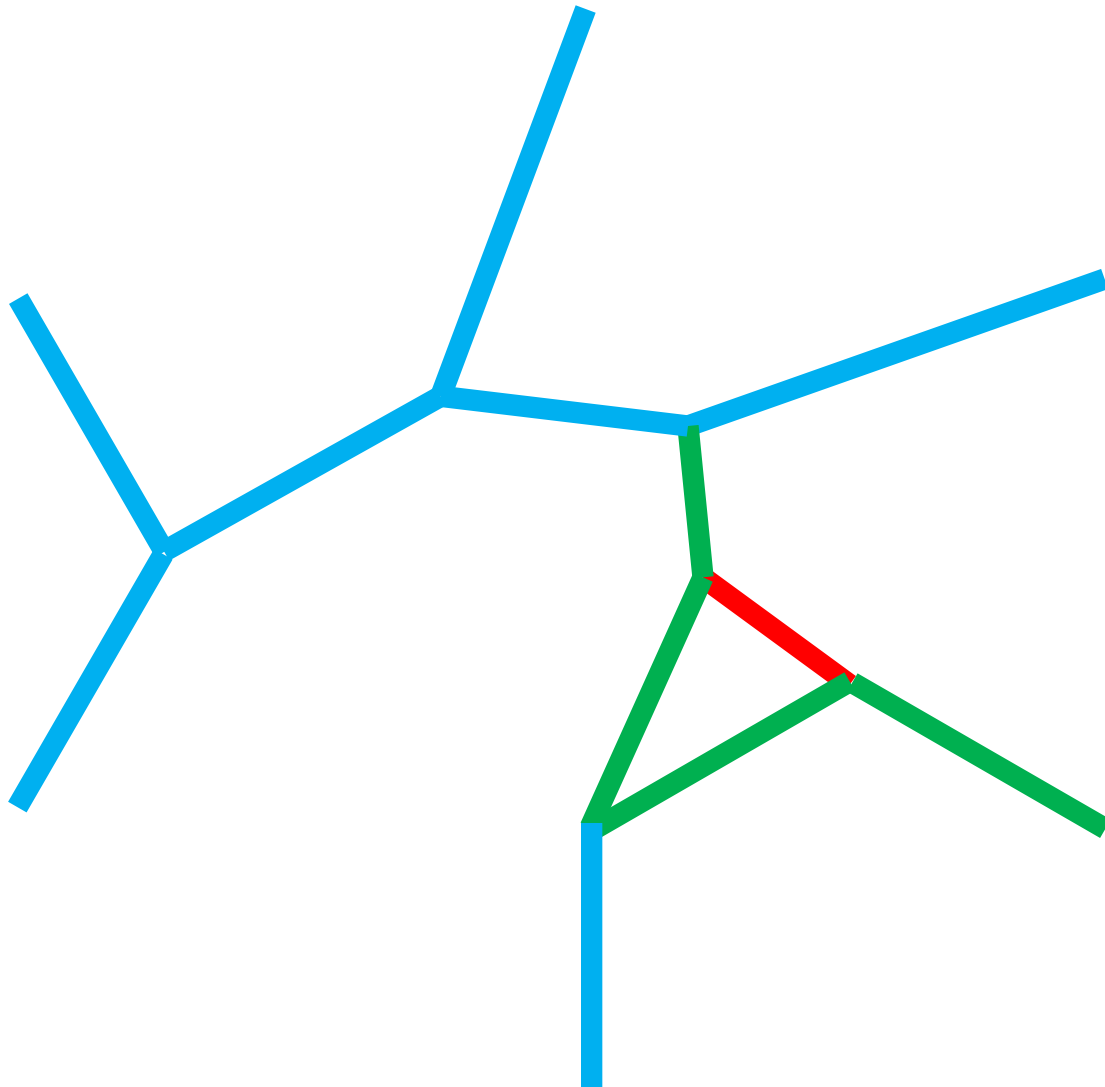
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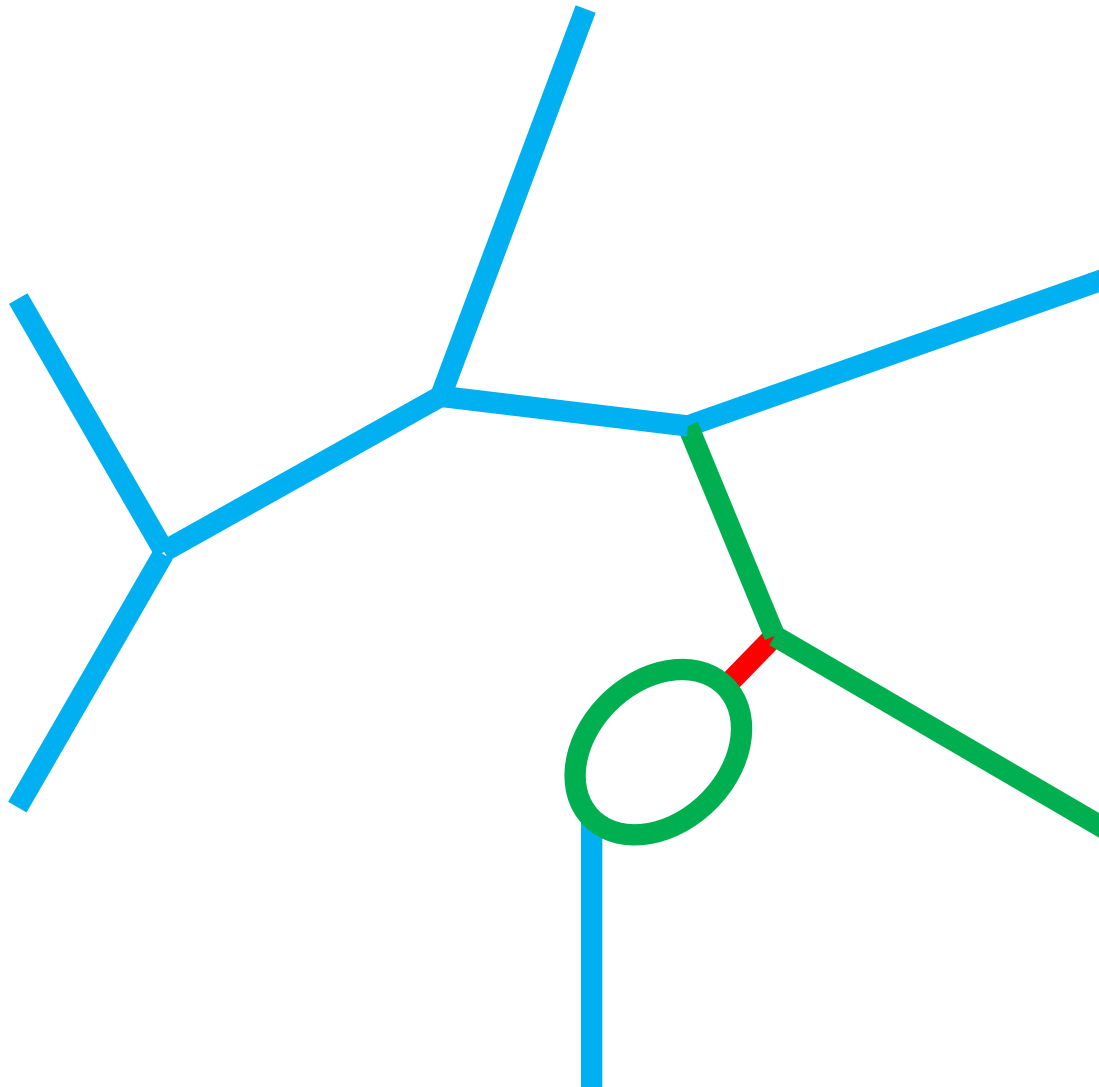
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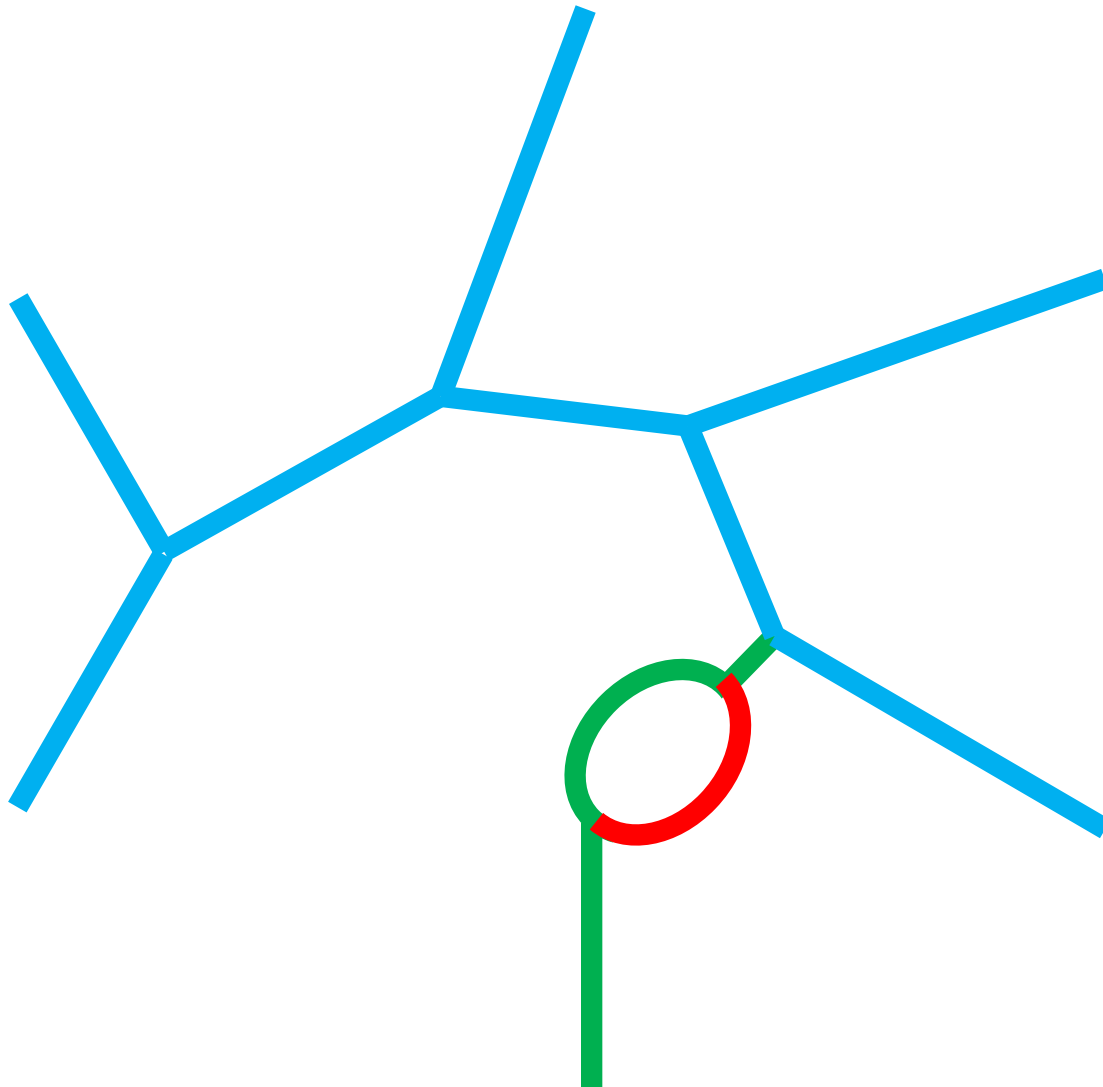
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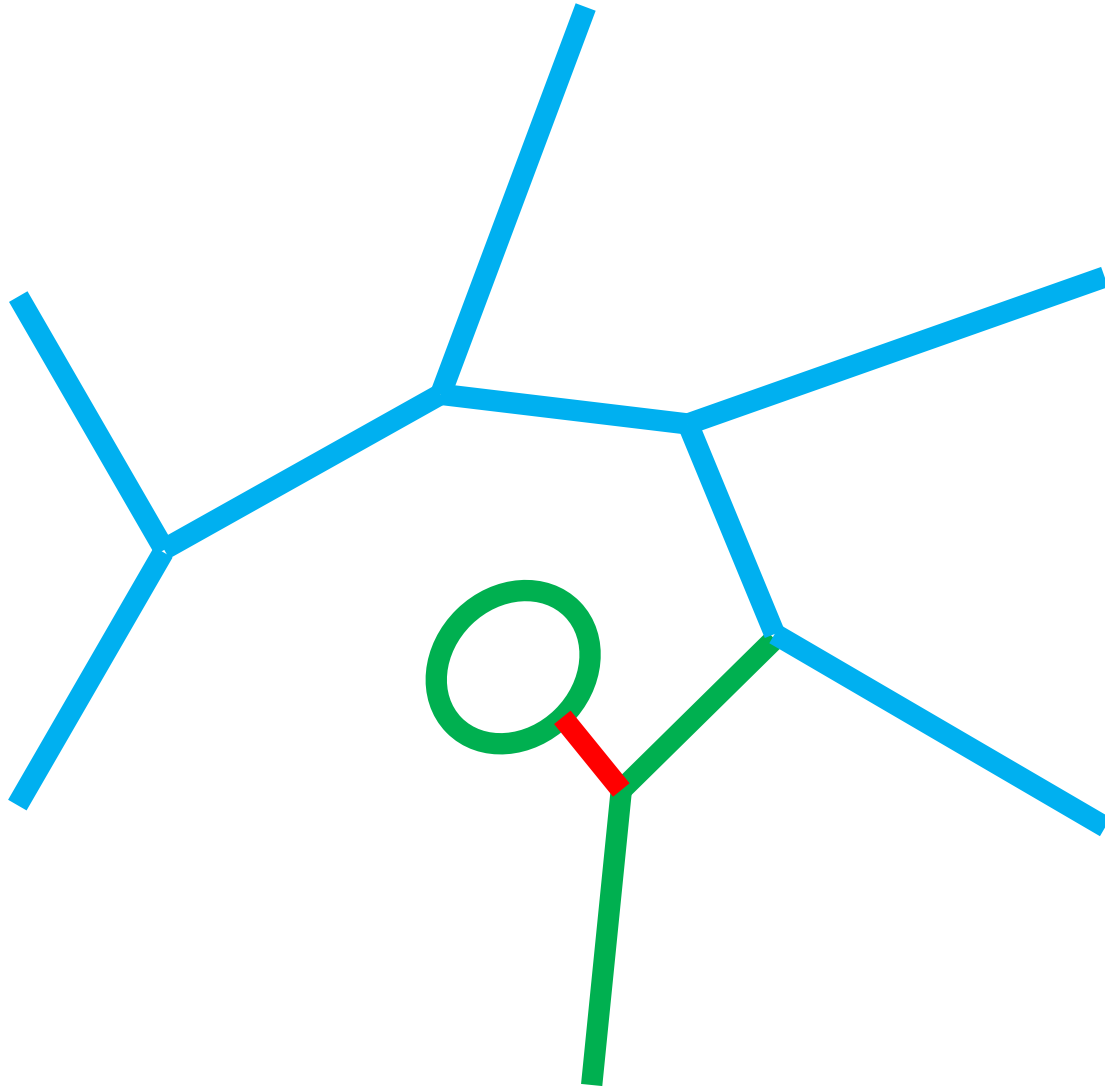
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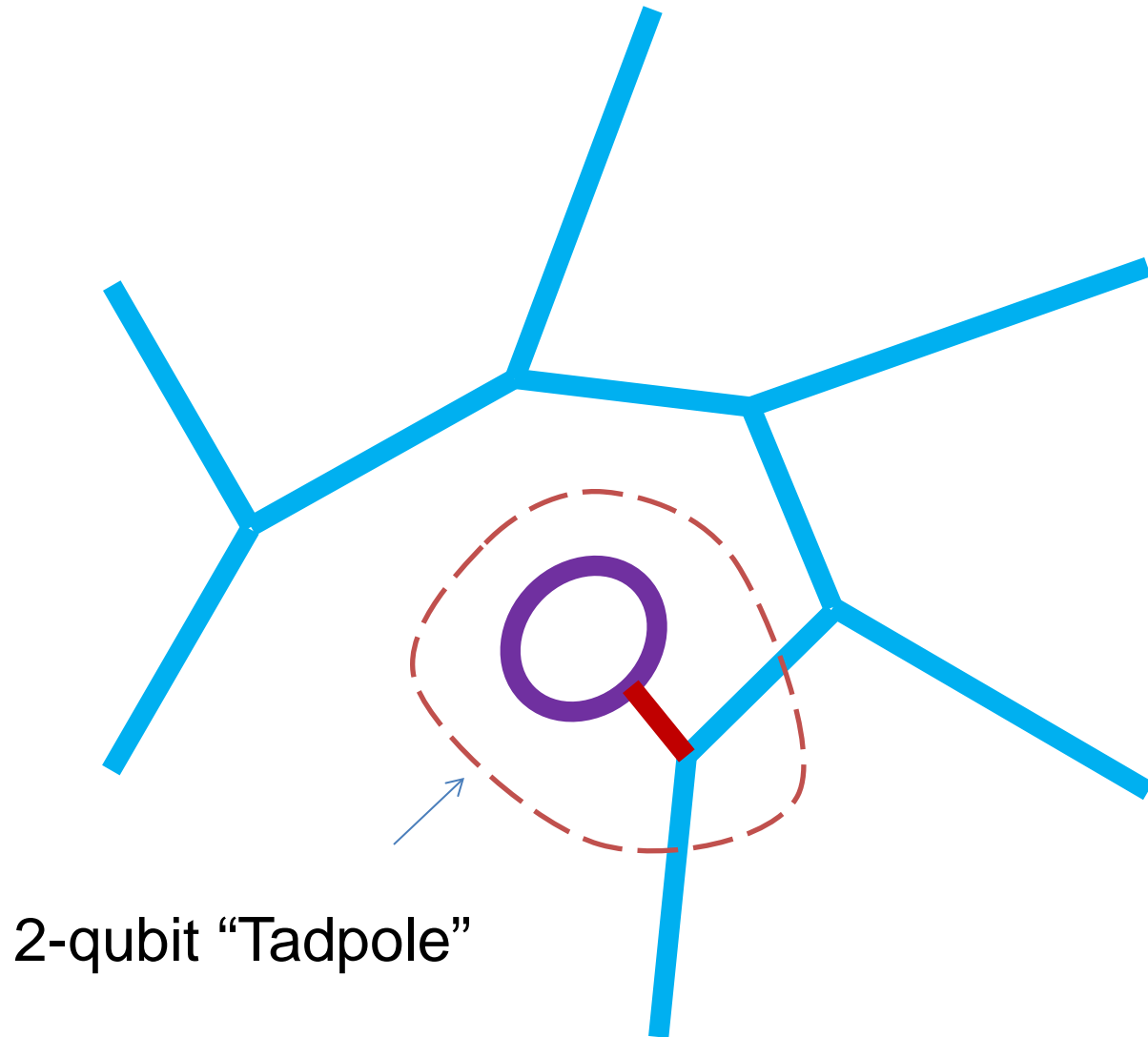
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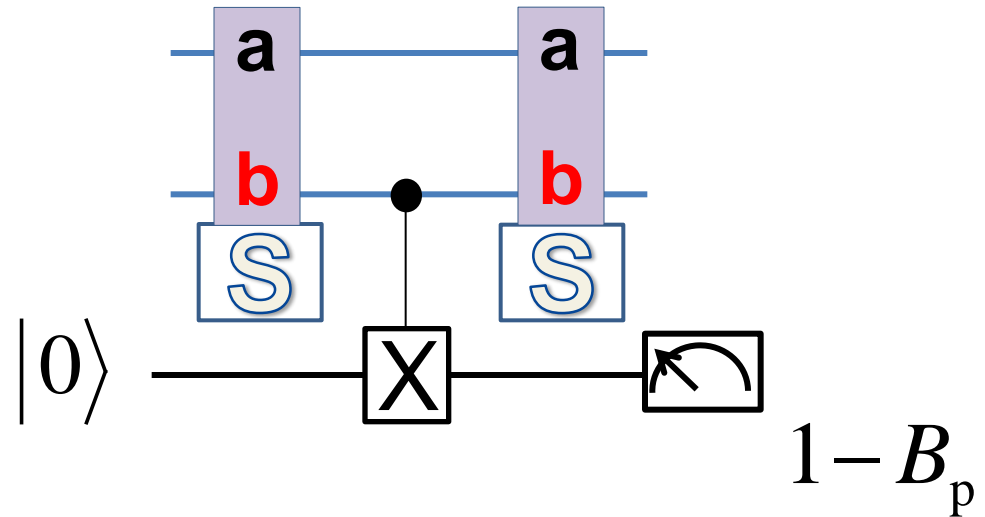
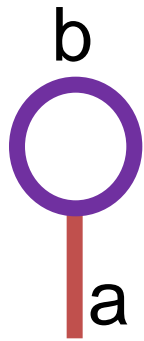
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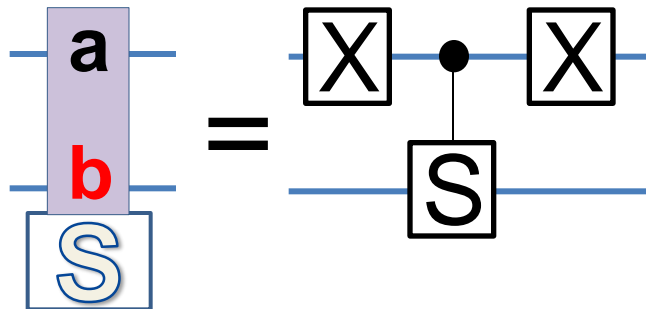
Plaquette Reduction



Measuring B_p for a Tadpole is Easy!

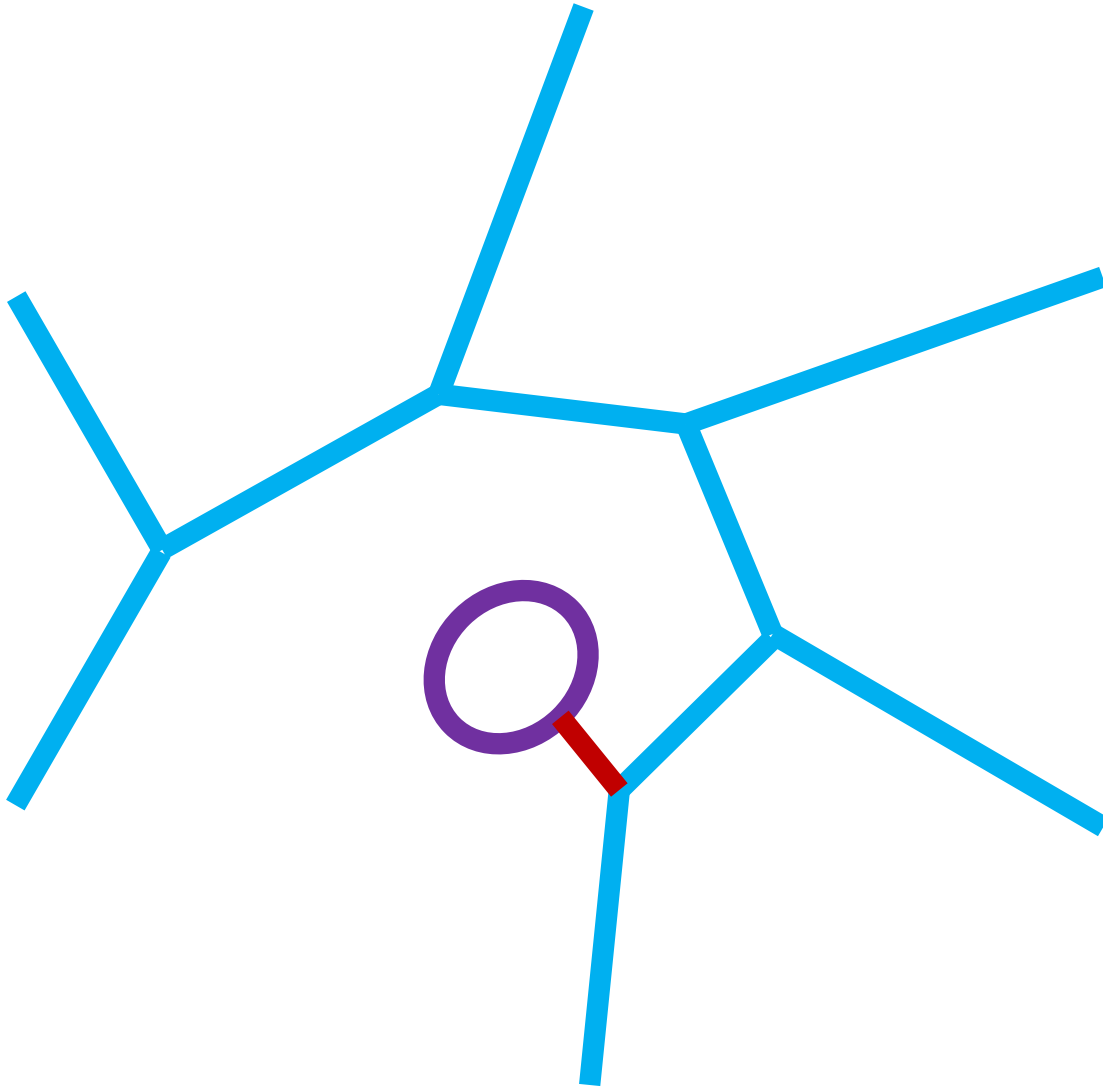


S Quantum Circuit

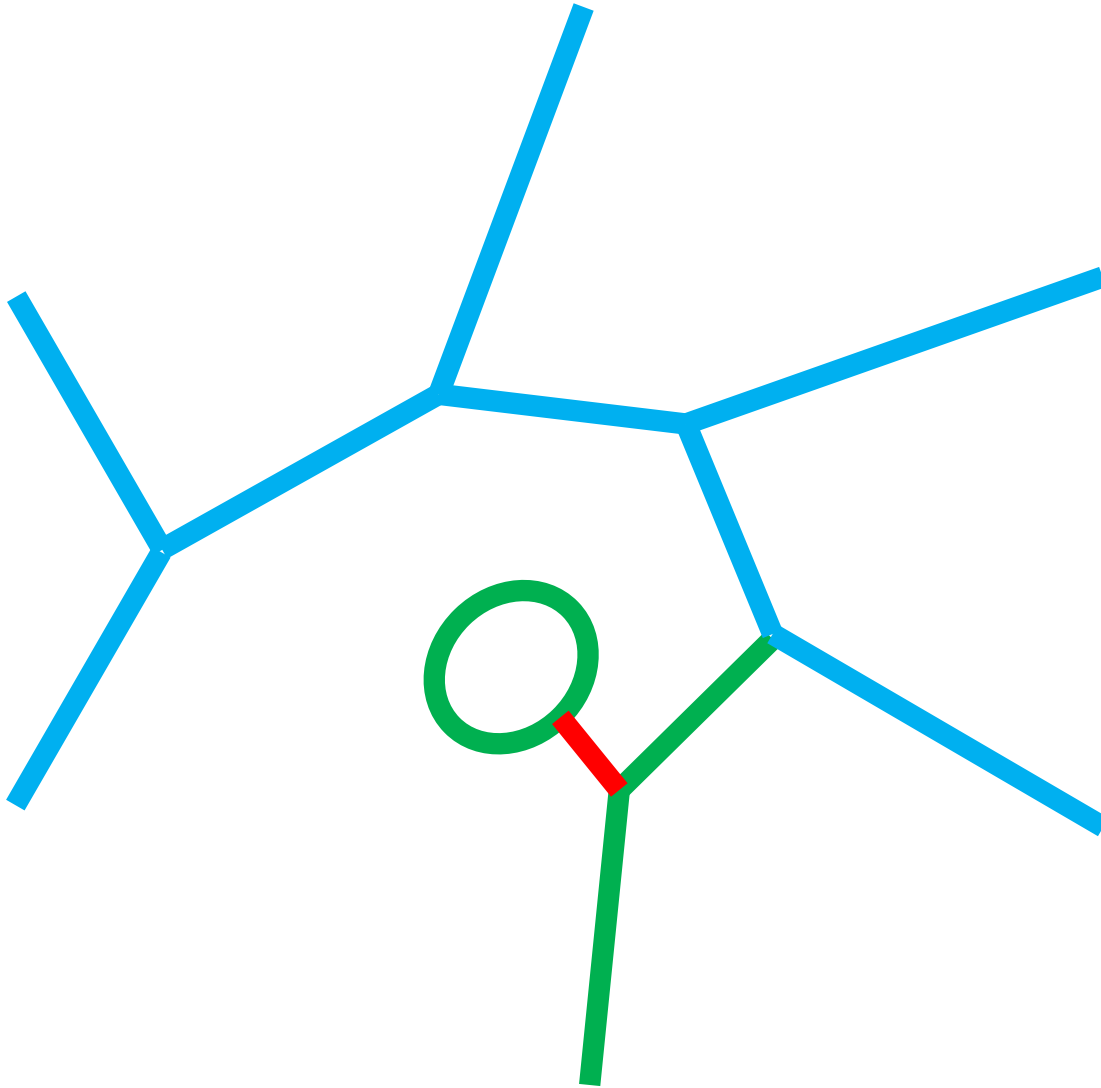


$$S = \frac{1}{\sqrt{1+\varphi^2}} \begin{pmatrix} 1 & \varphi \\ \varphi & -1 \end{pmatrix}$$

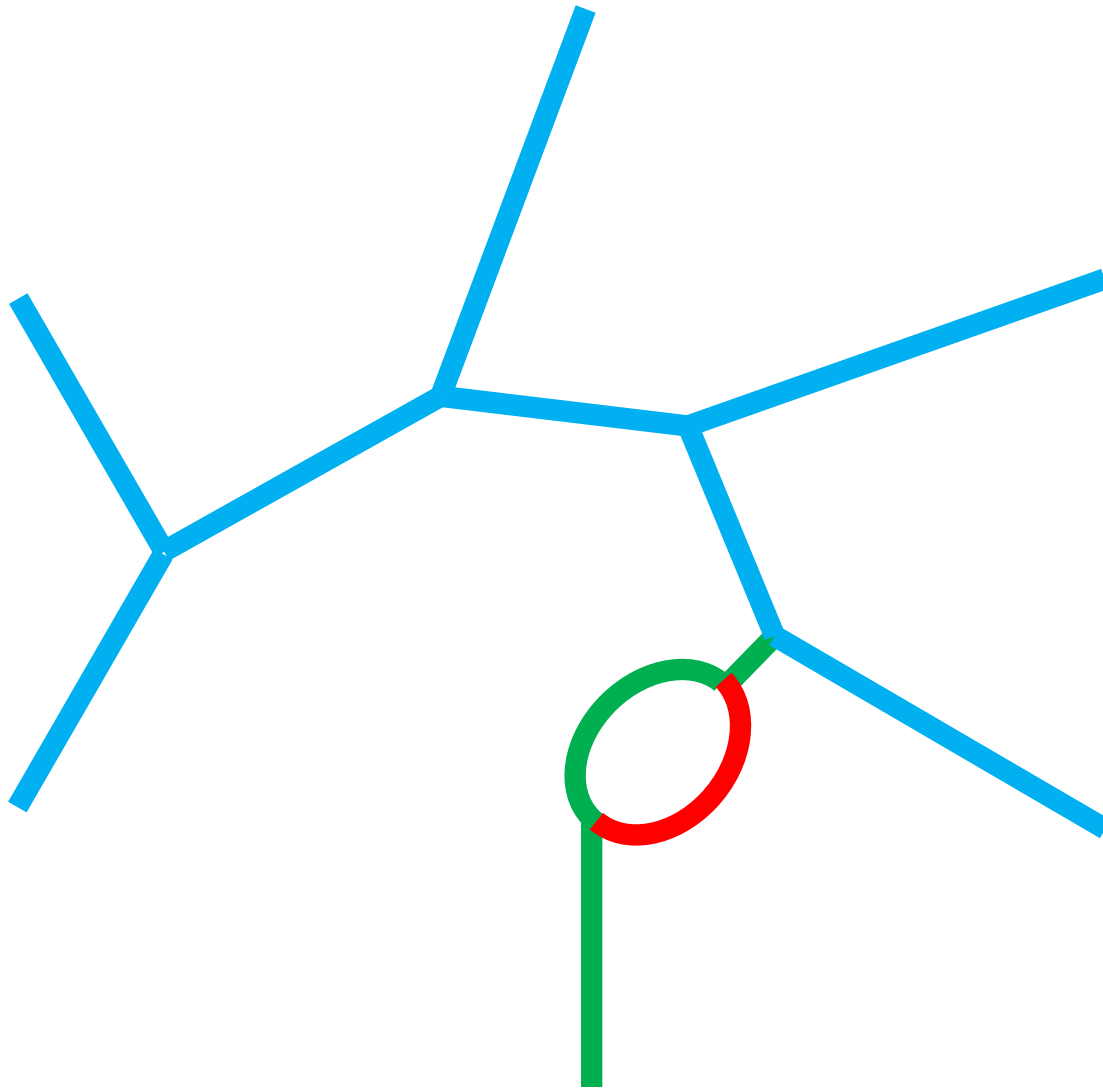
Plaque Restoration



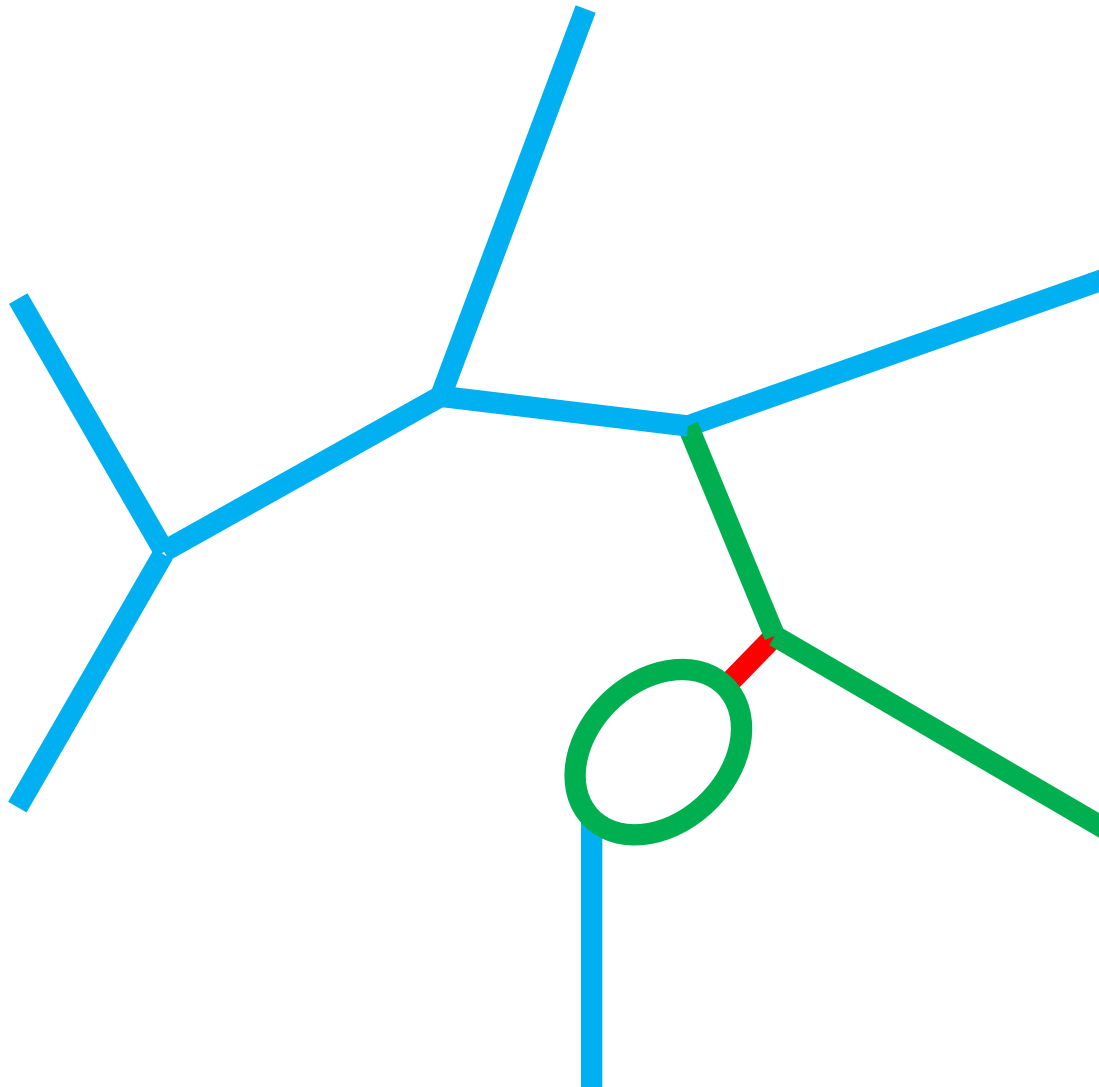
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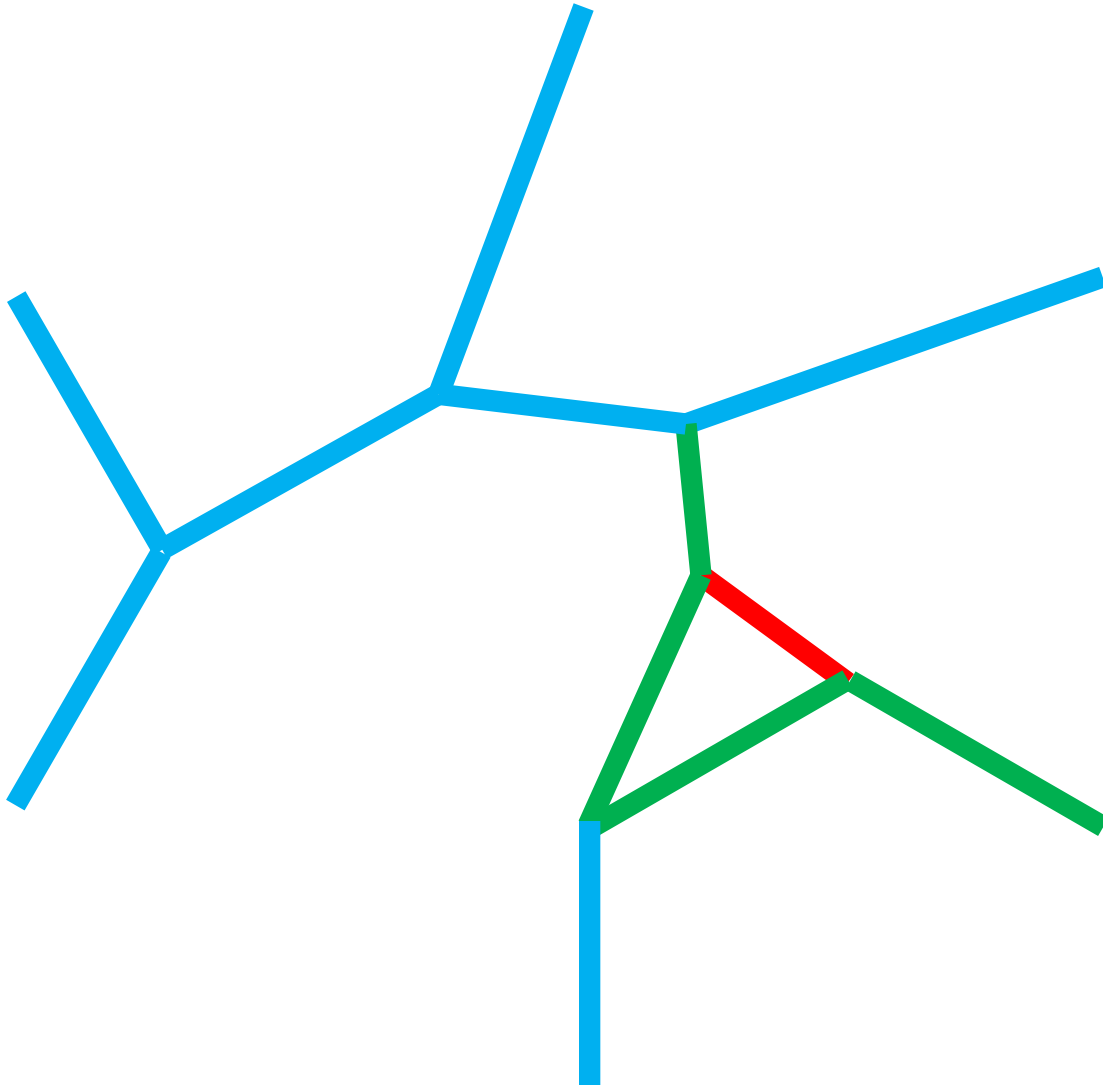
Plaque Restoration



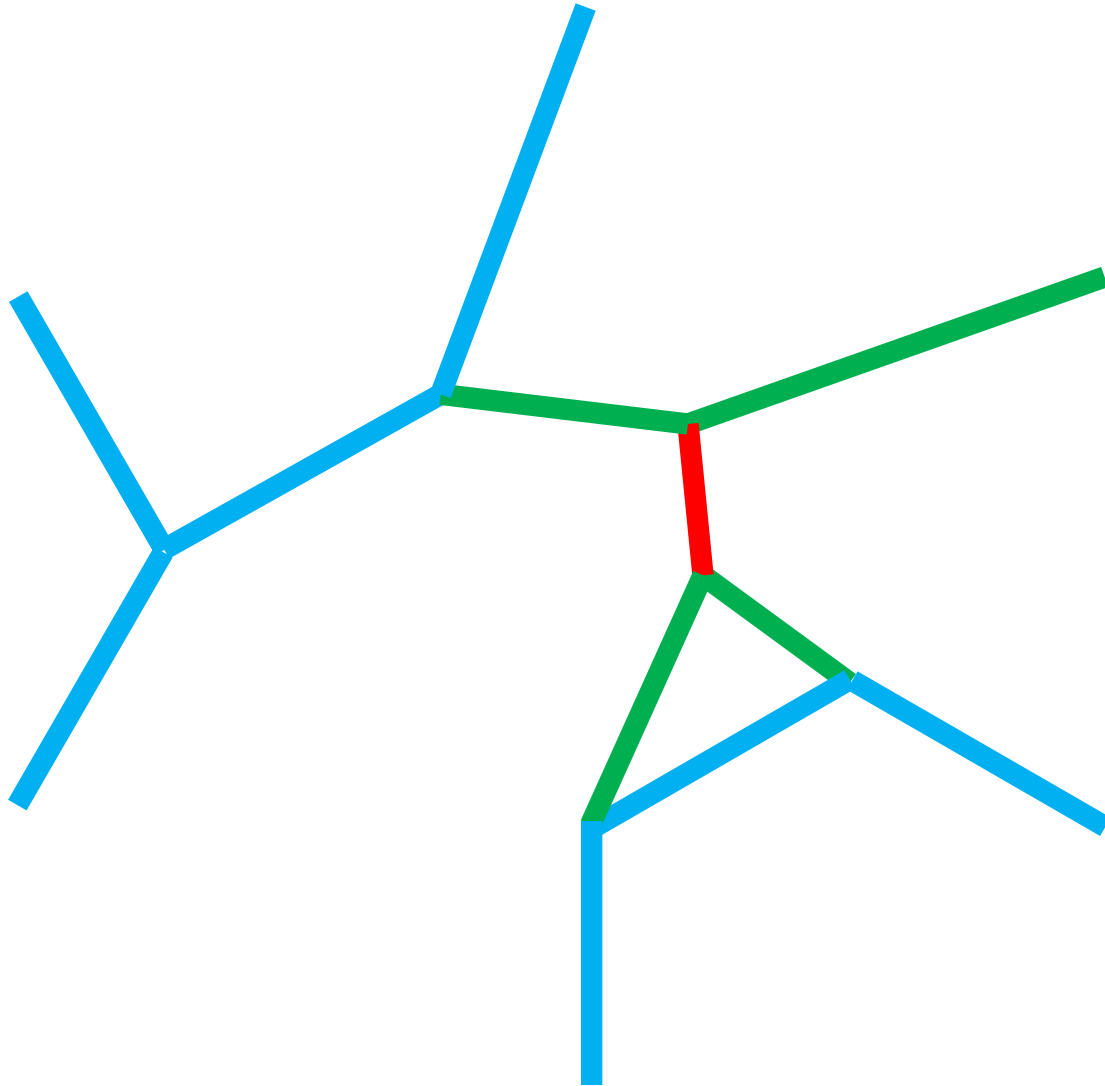
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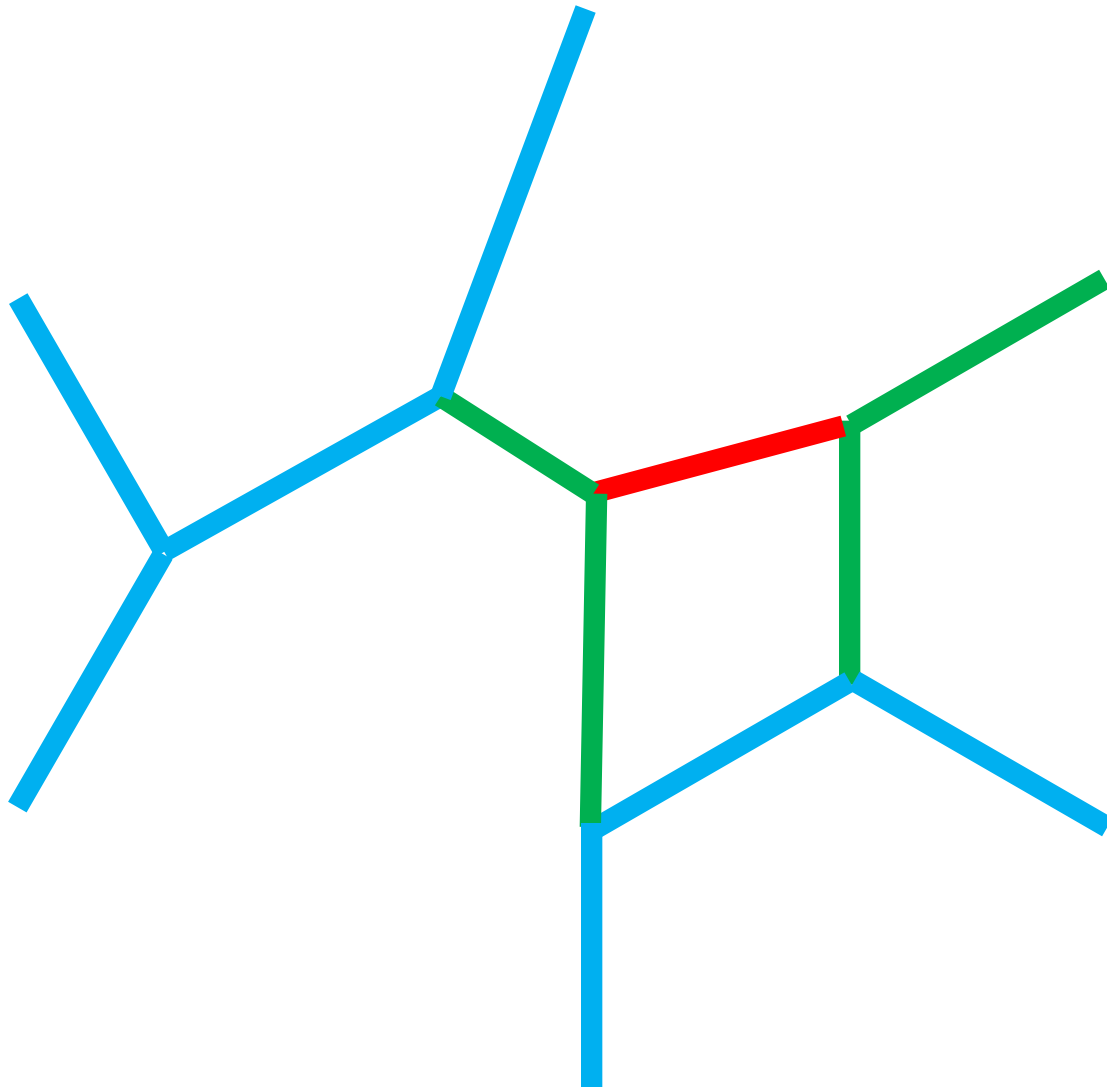
Plaque Restoration



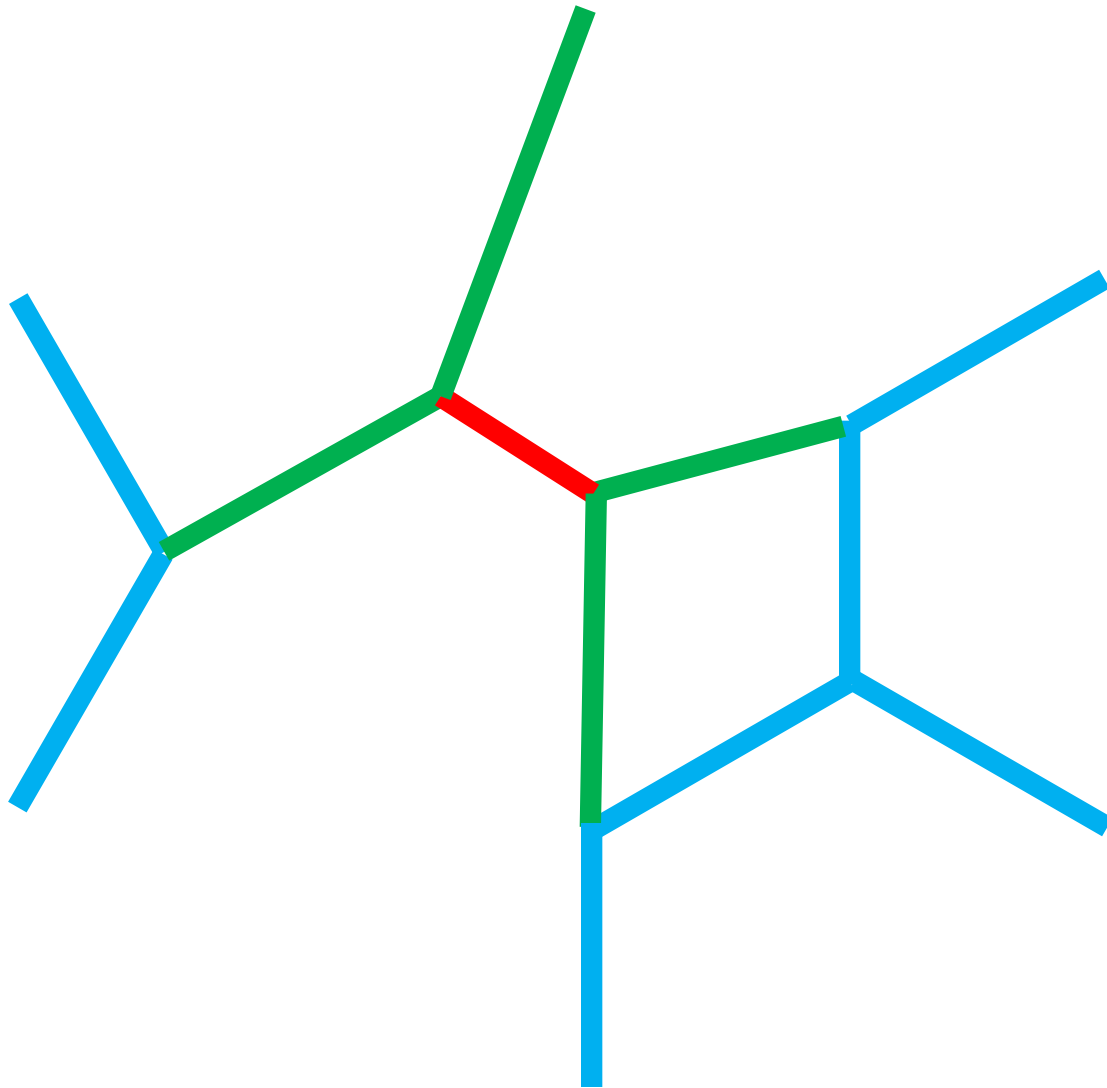
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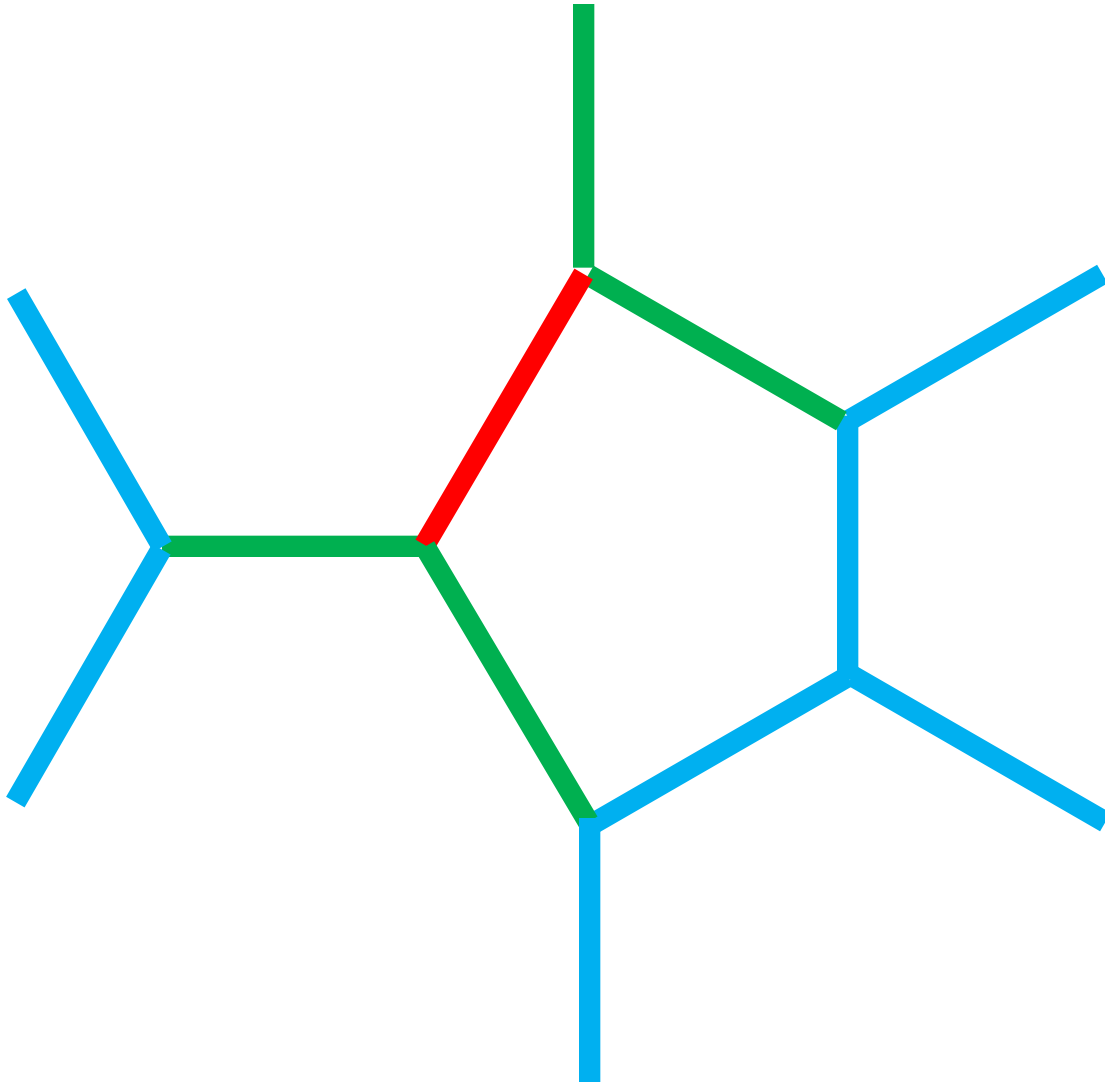
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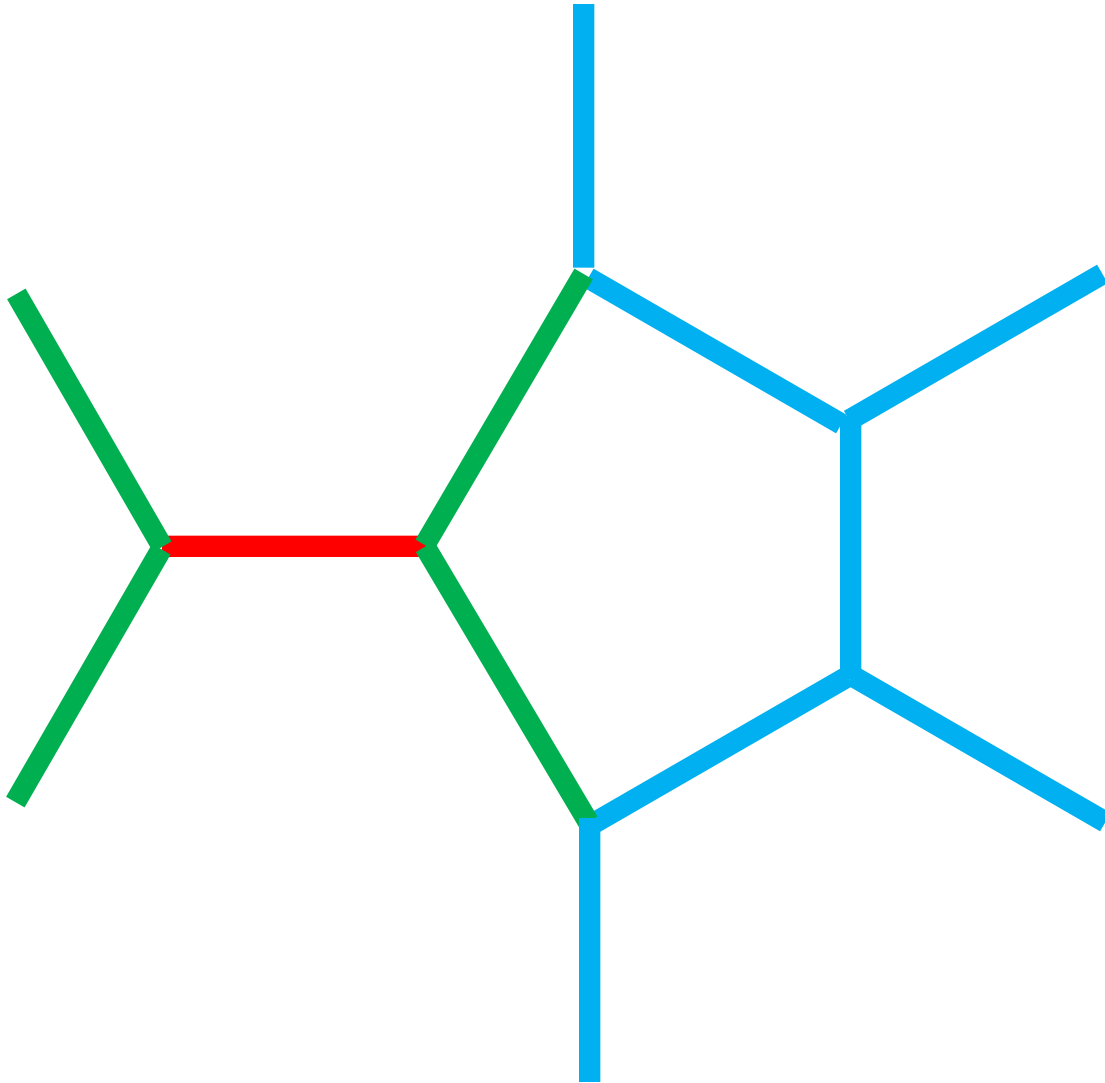
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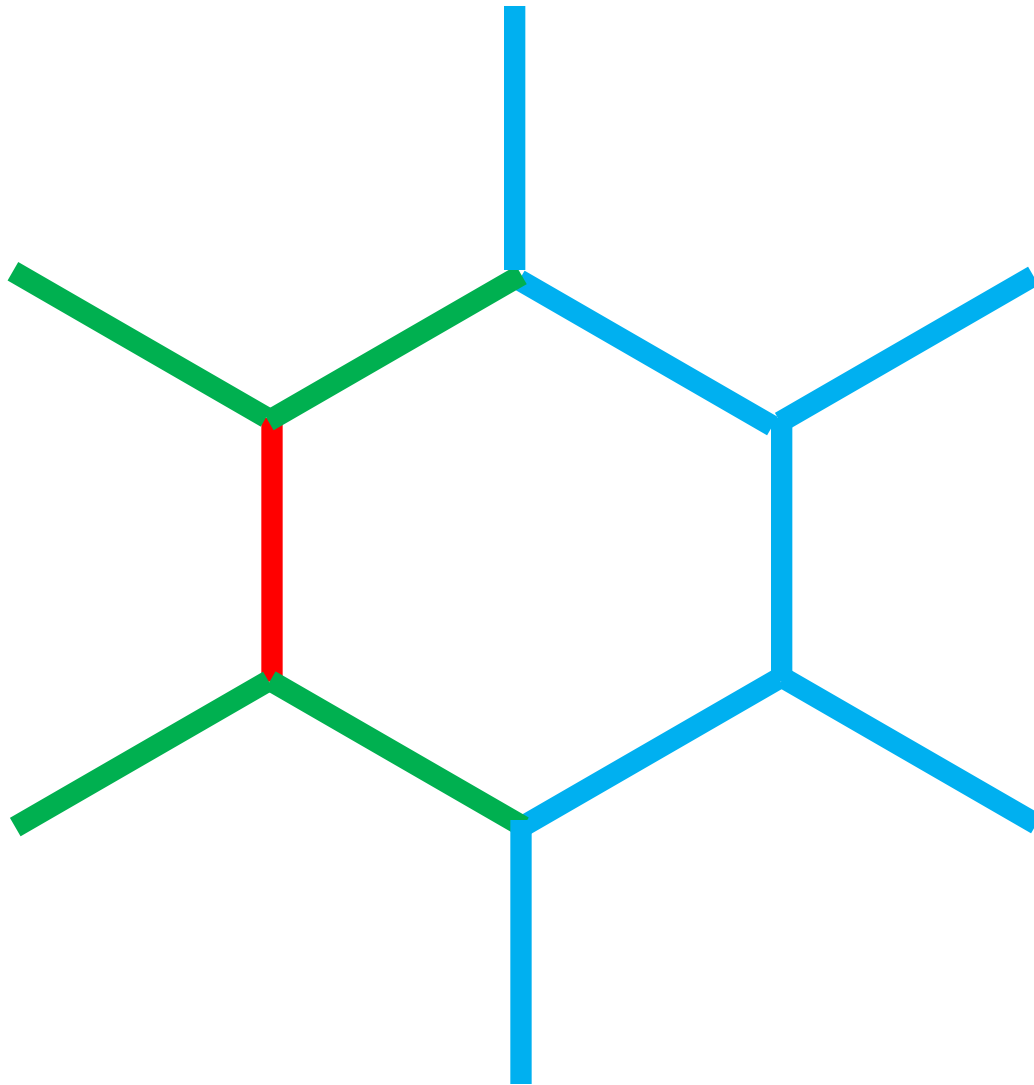
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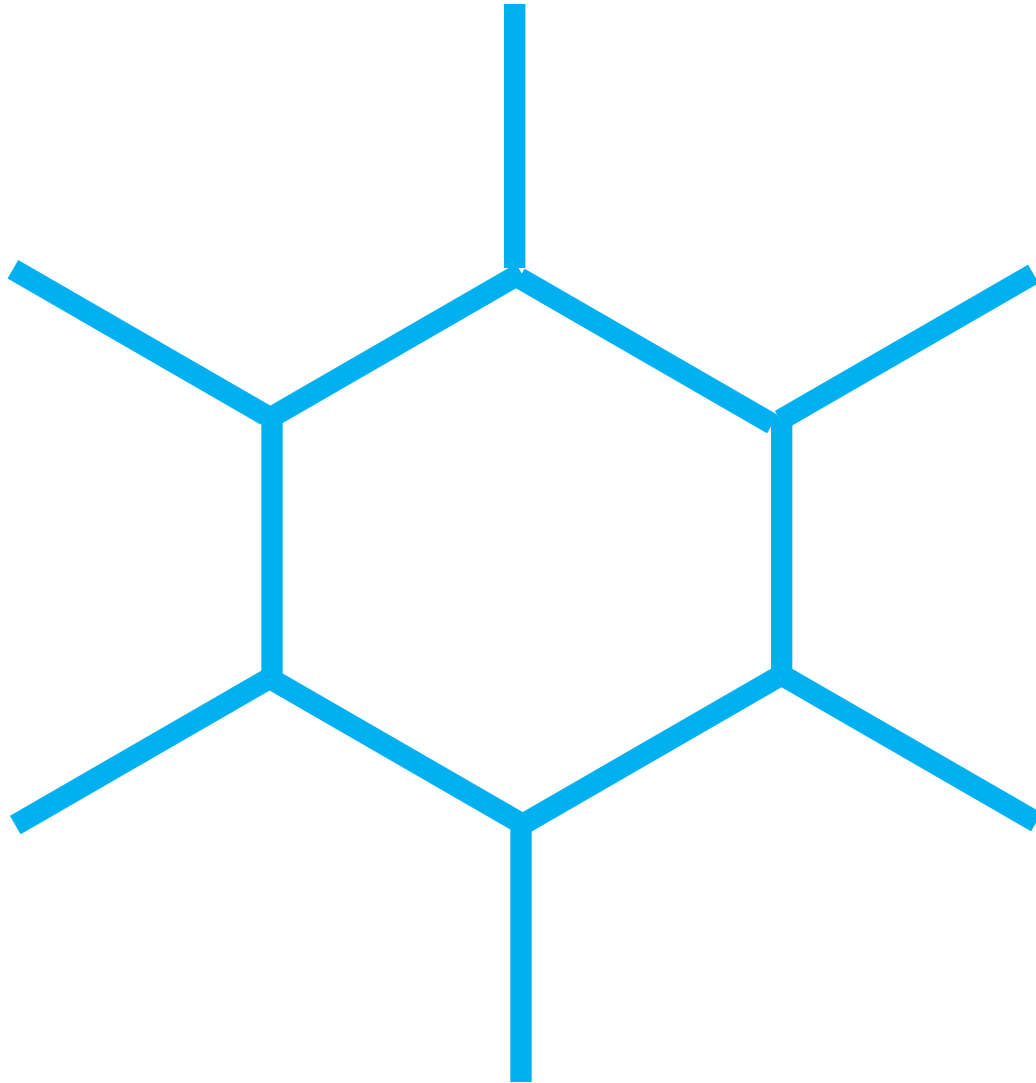
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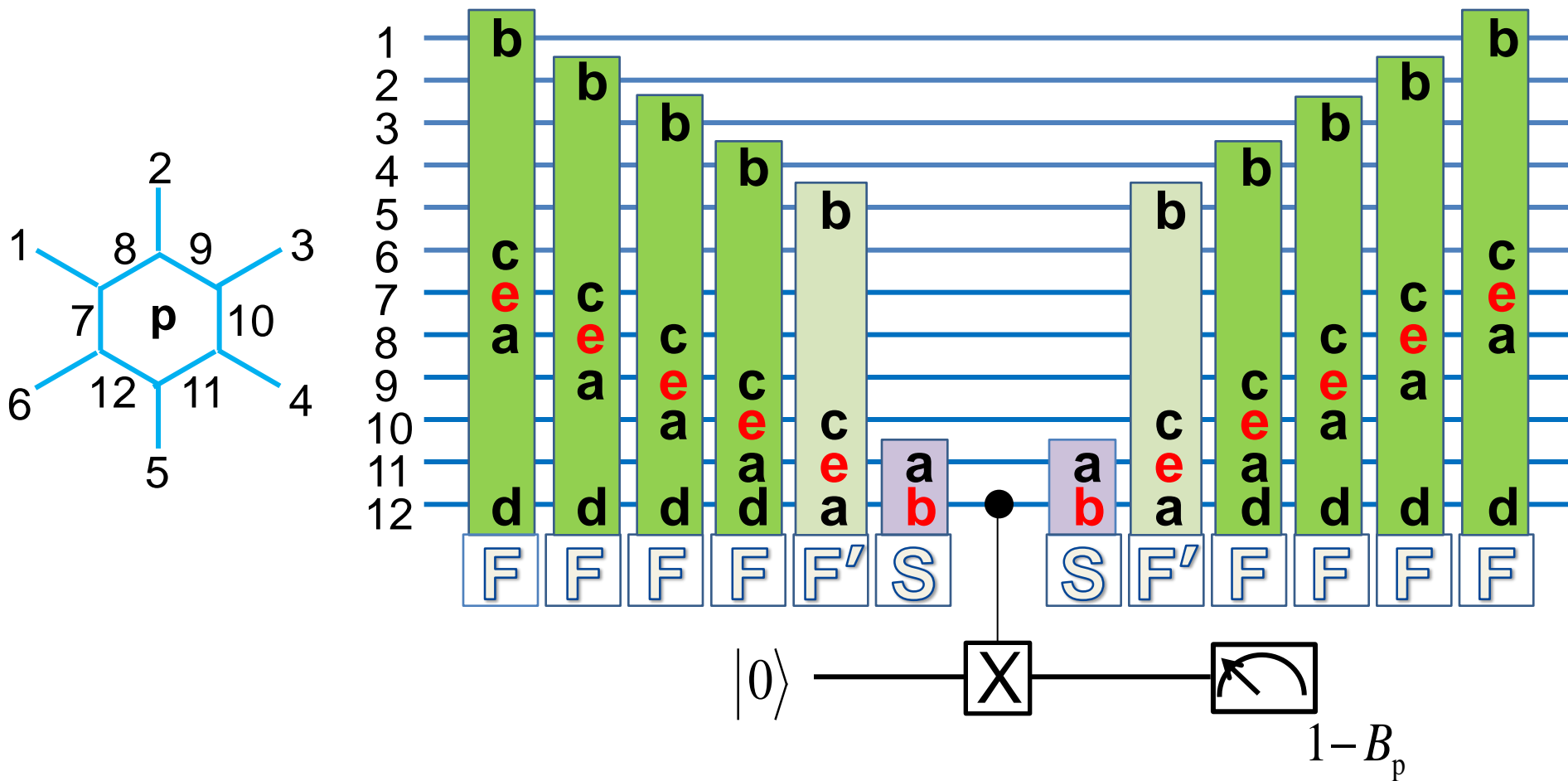
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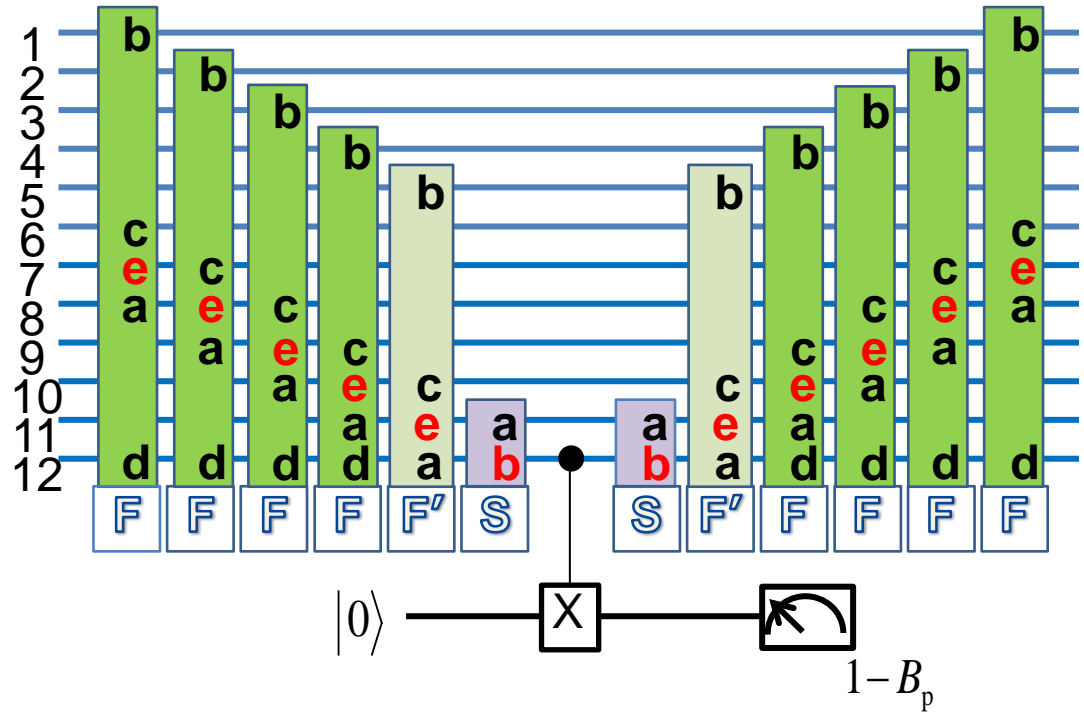
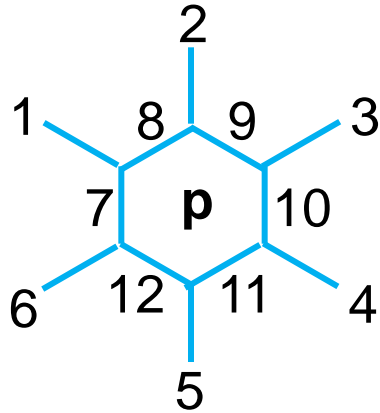
Plaque Restoration



Quantum Circuit for Measuring B_p



Gate Count

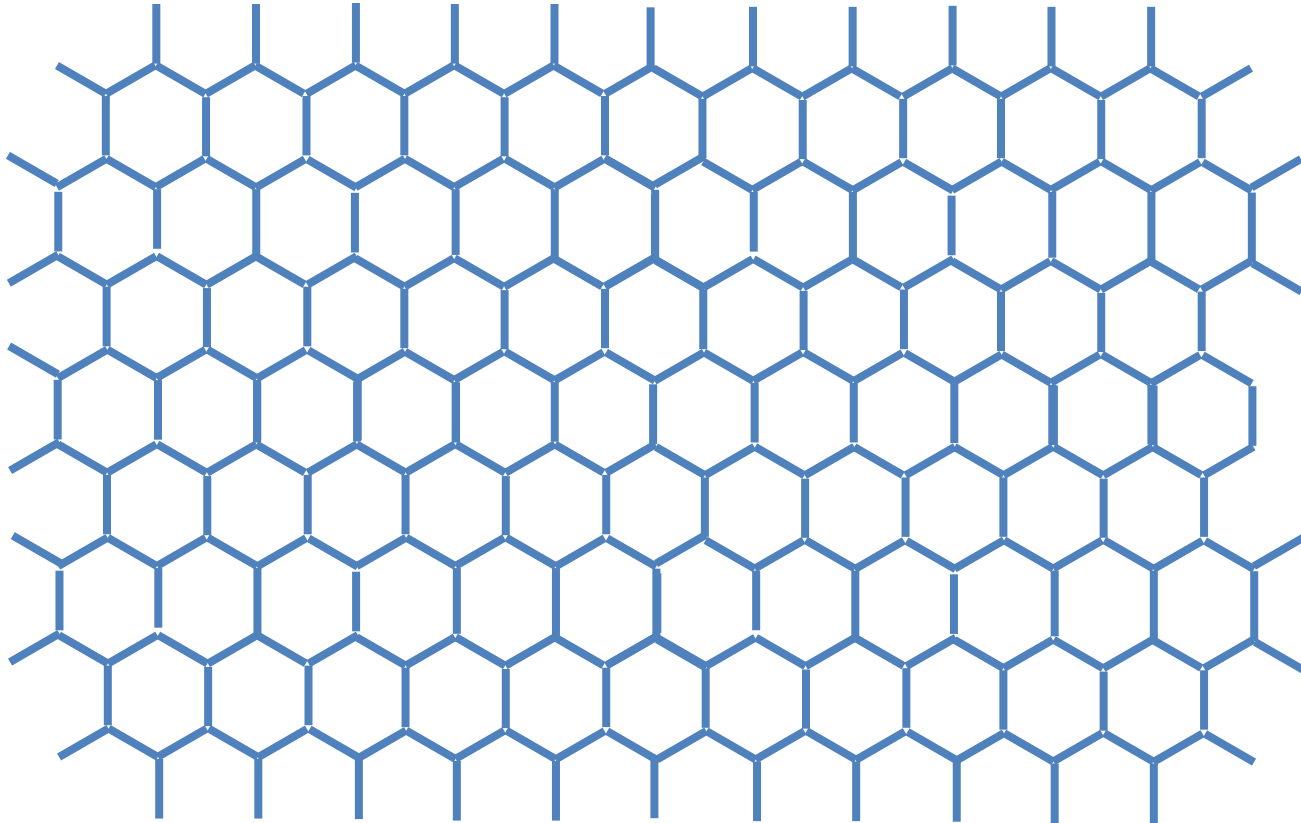


8 5-qubit Toffoli Gates
 2 4-qubit Toffoli Gates
 10 3-qubit Toffoli Gates
 43 CNOT Gates
 24 Single Qubit Gates

371 CNOT Gates
 392 Single Qubit Rotations

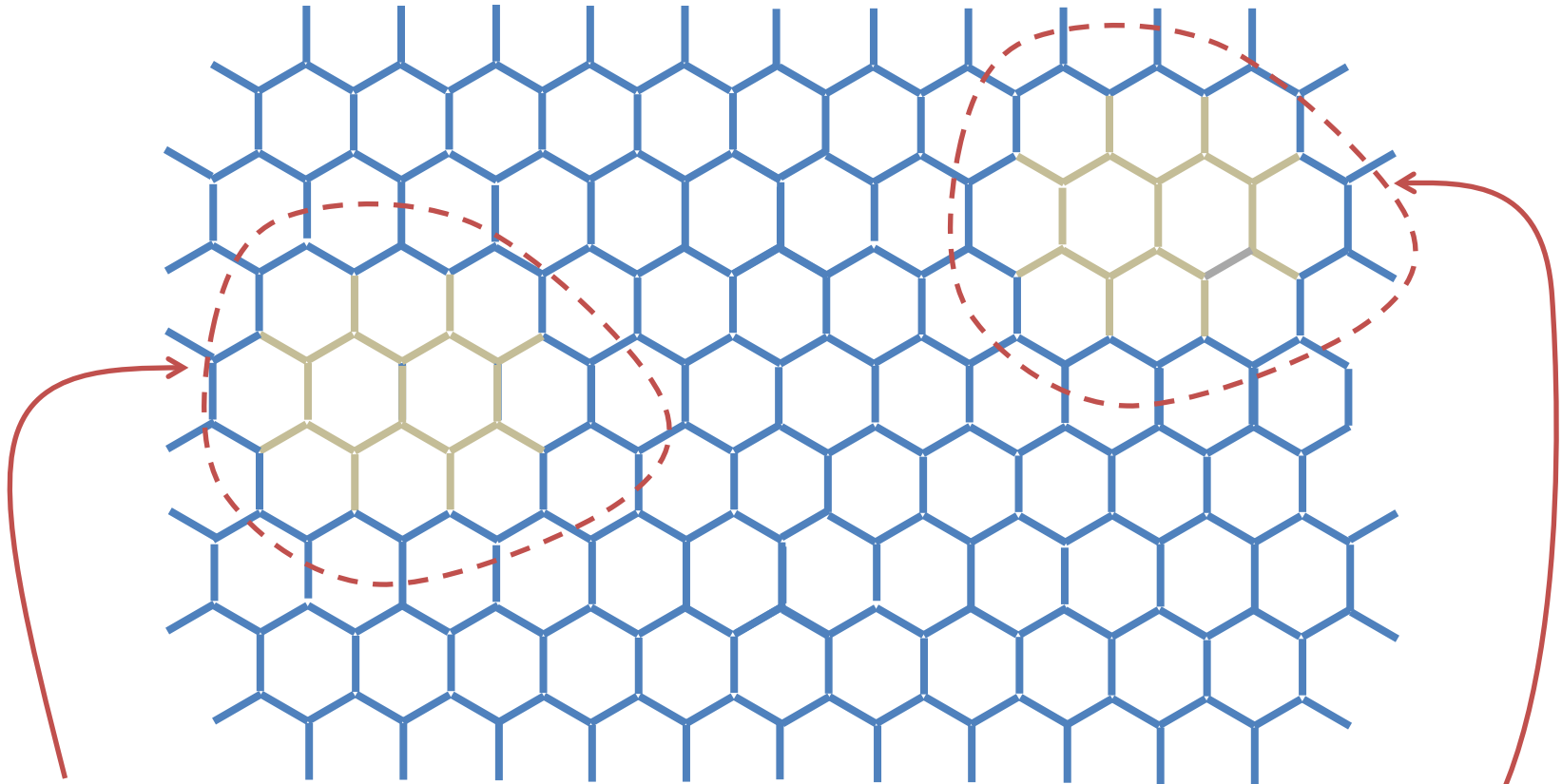
Creating Excited States: Fibonacci Anyons

Konig, Kuperberg, Reichardt, Ann. Phys. 2010



Creating Excited States: Fibonacci Anyons

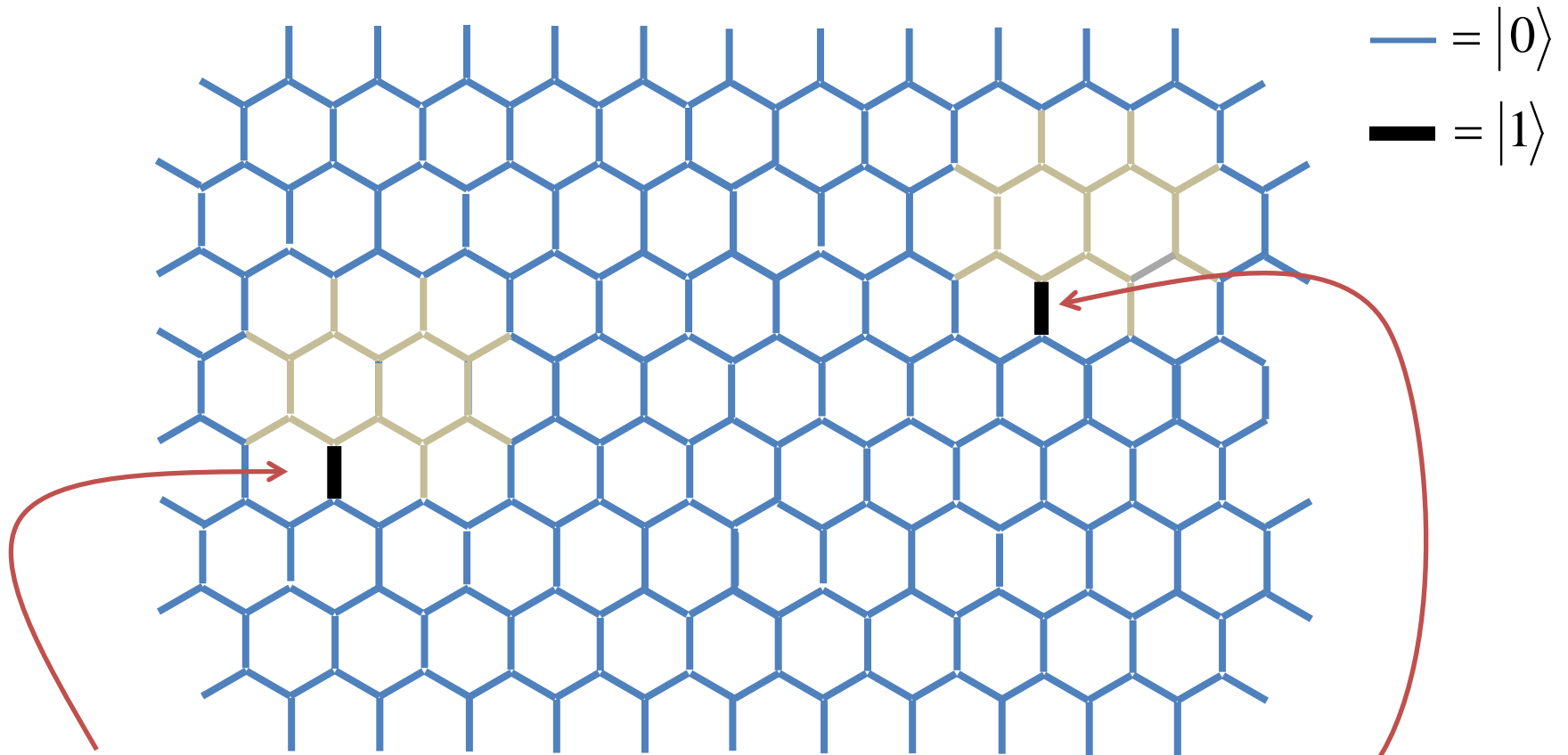
Konig, Kuperberg, Reichardt, Ann. Phys. 2010



Introduce "holes"
(Stop measuring Q_v and B_p
inside hole)

Creating Excited States: Fibonacci Anyons

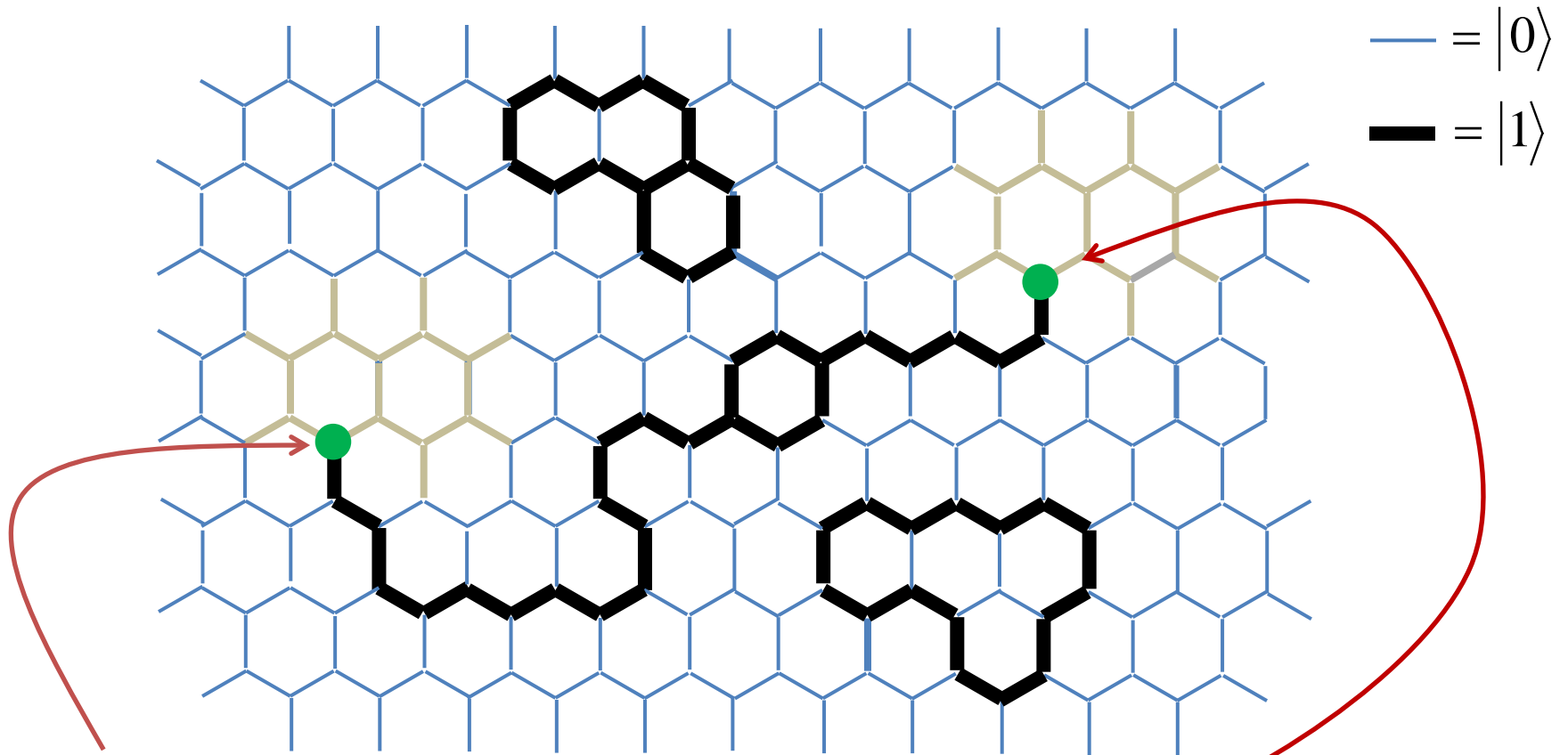
Konig, Kuperberg, Reichardt, Ann. Phys. 2010



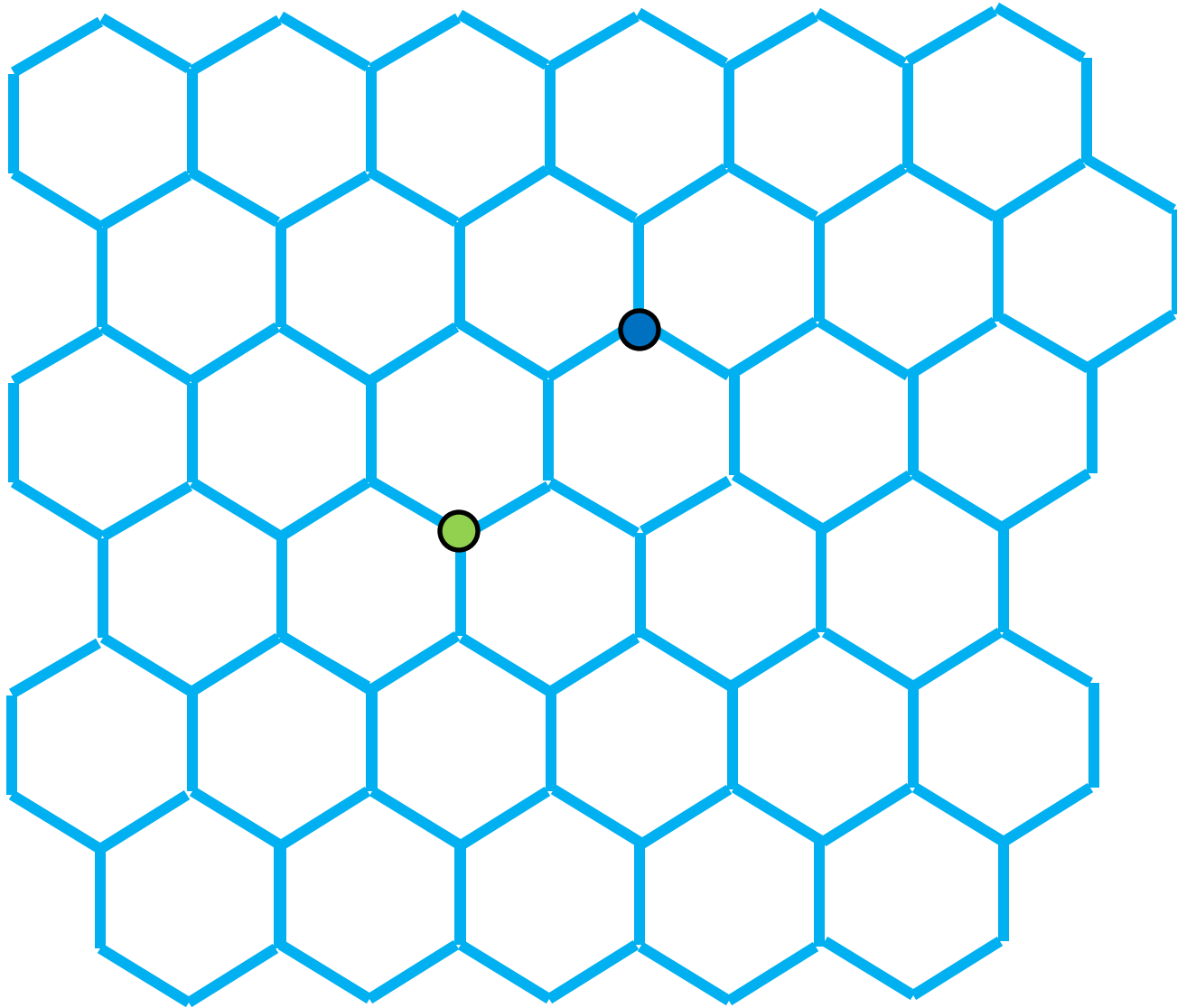
Fix boundary condition by
modifying Q_v operators

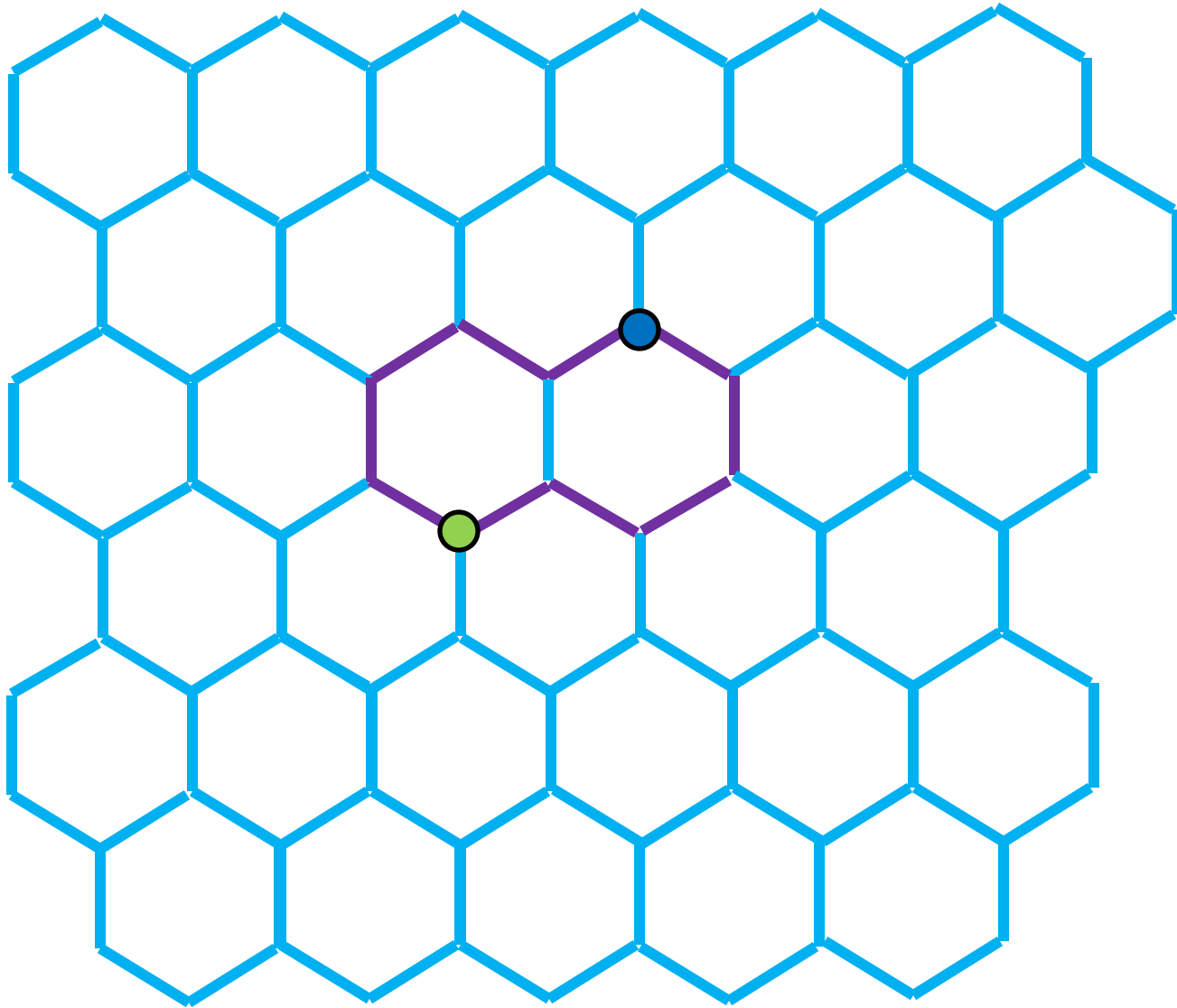
Creating Excited States: Fibonacci Anyons

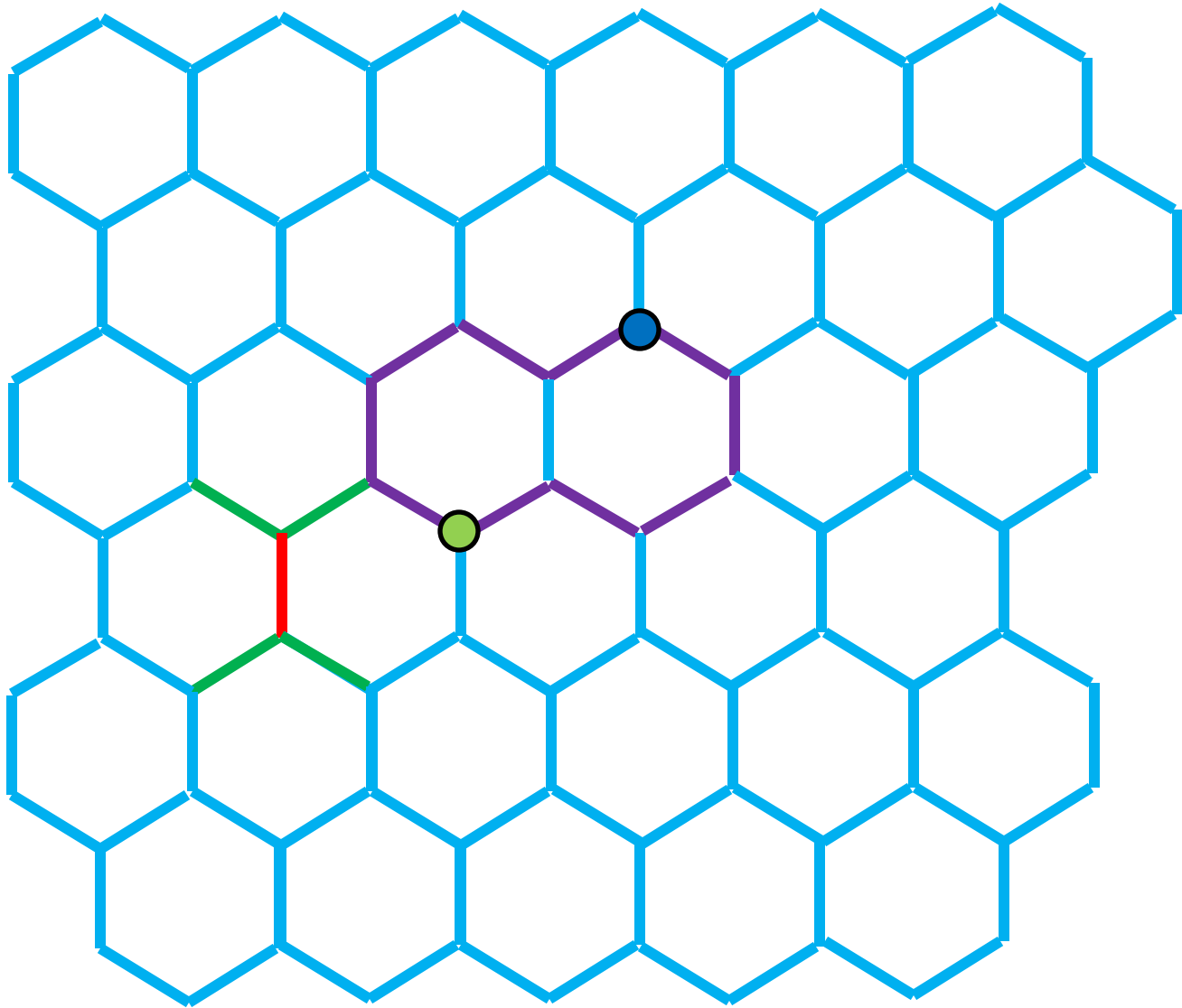
Konig, Kuperberg, Reichardt, Ann. Phys. 2010

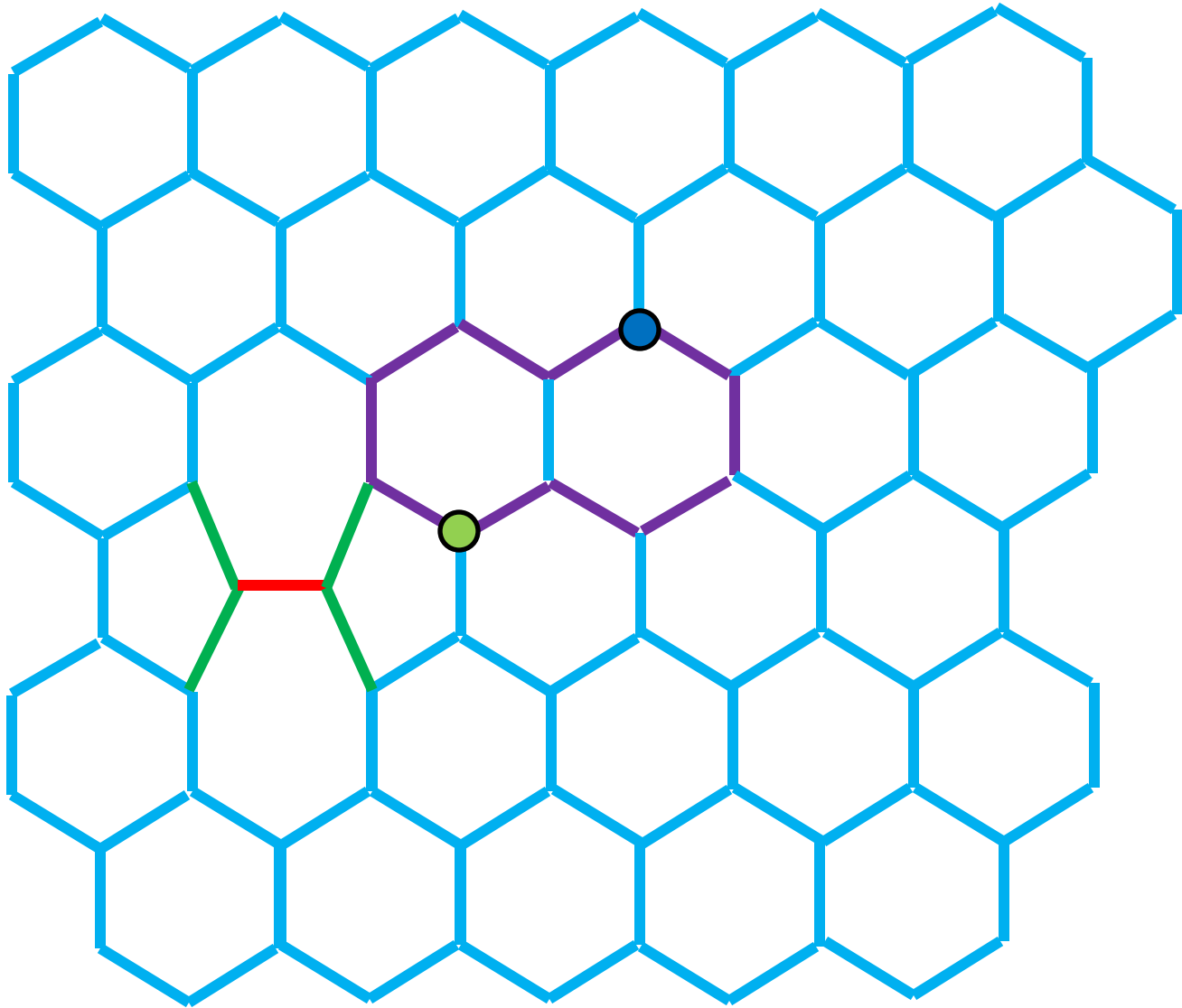


Fibonacci anyons

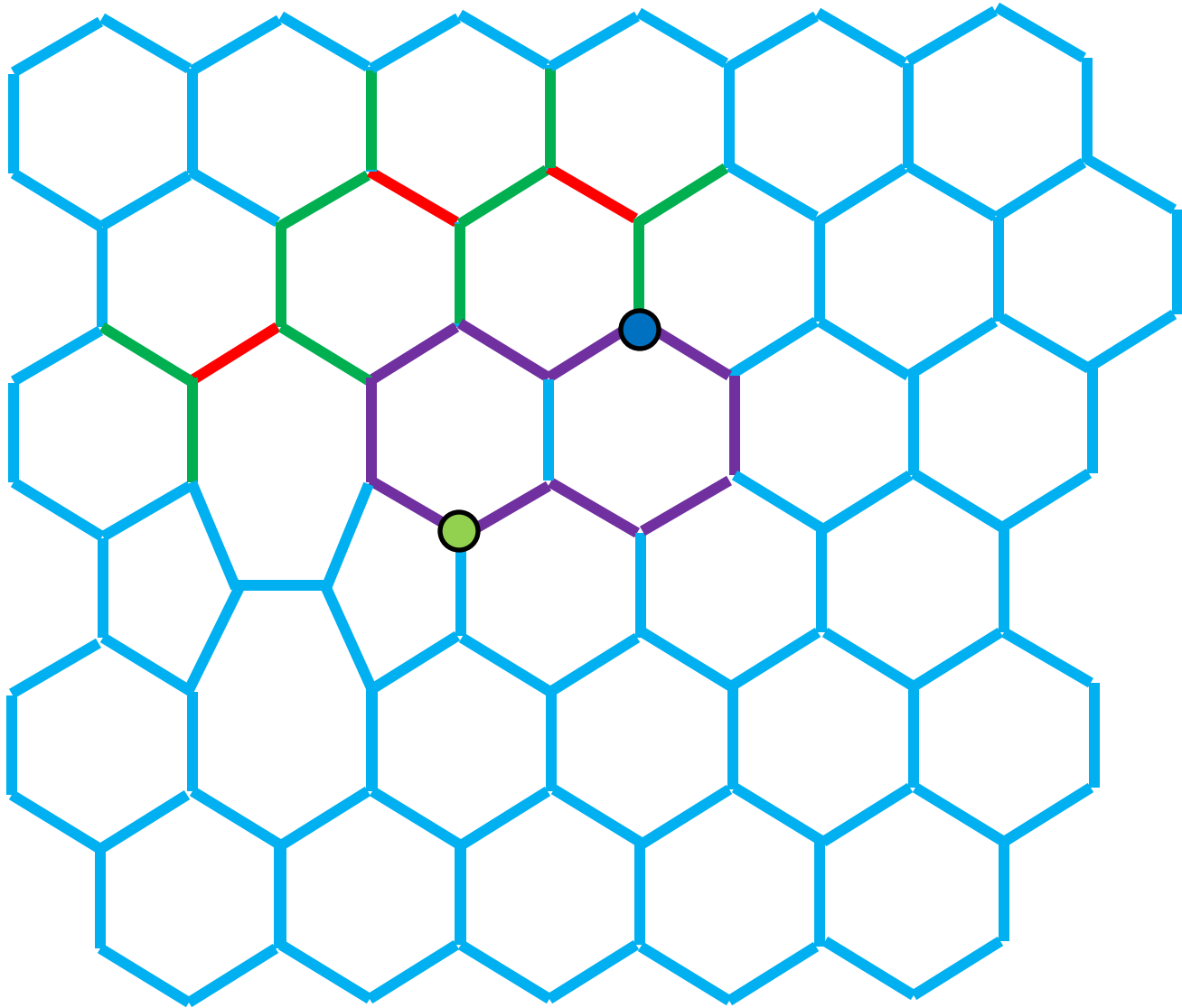




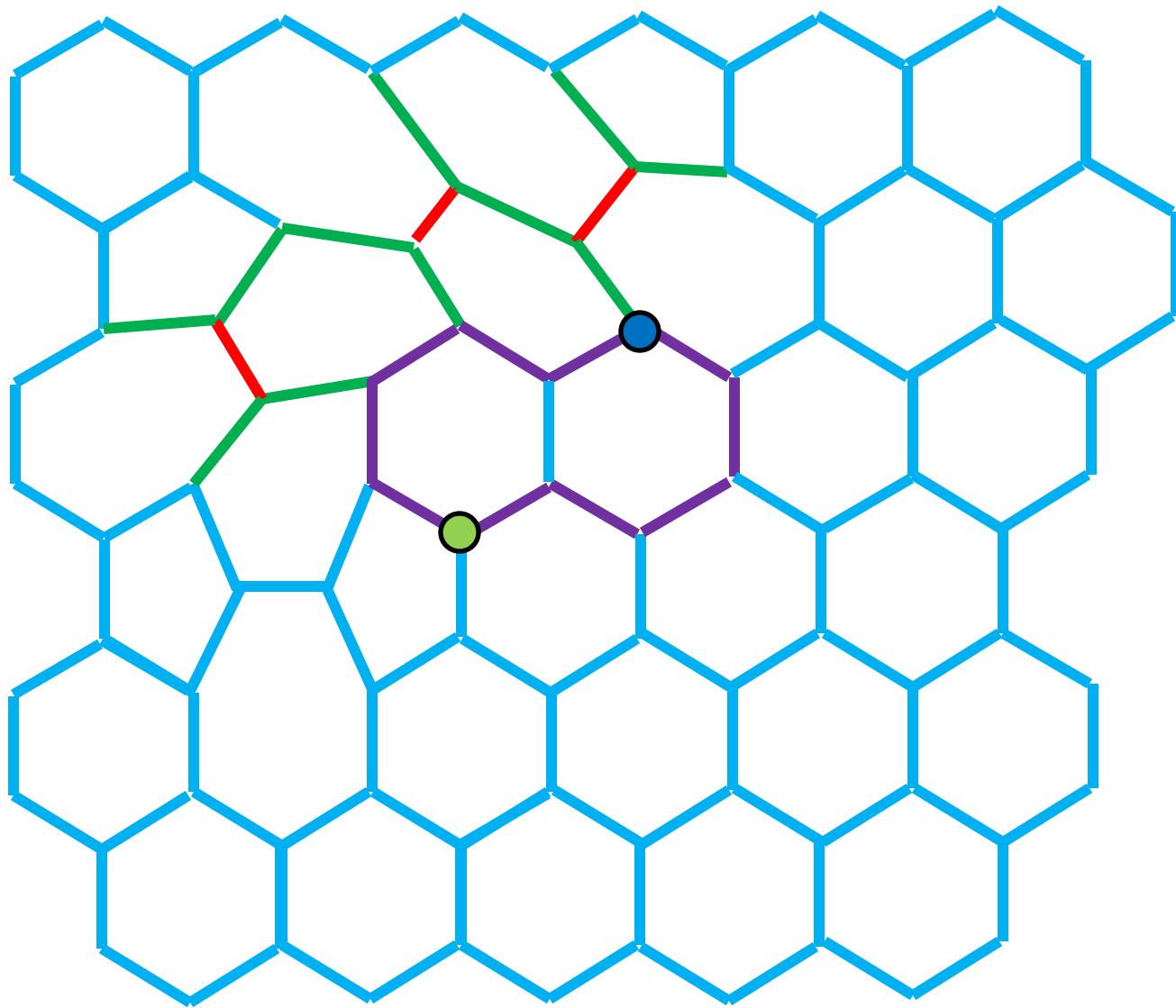




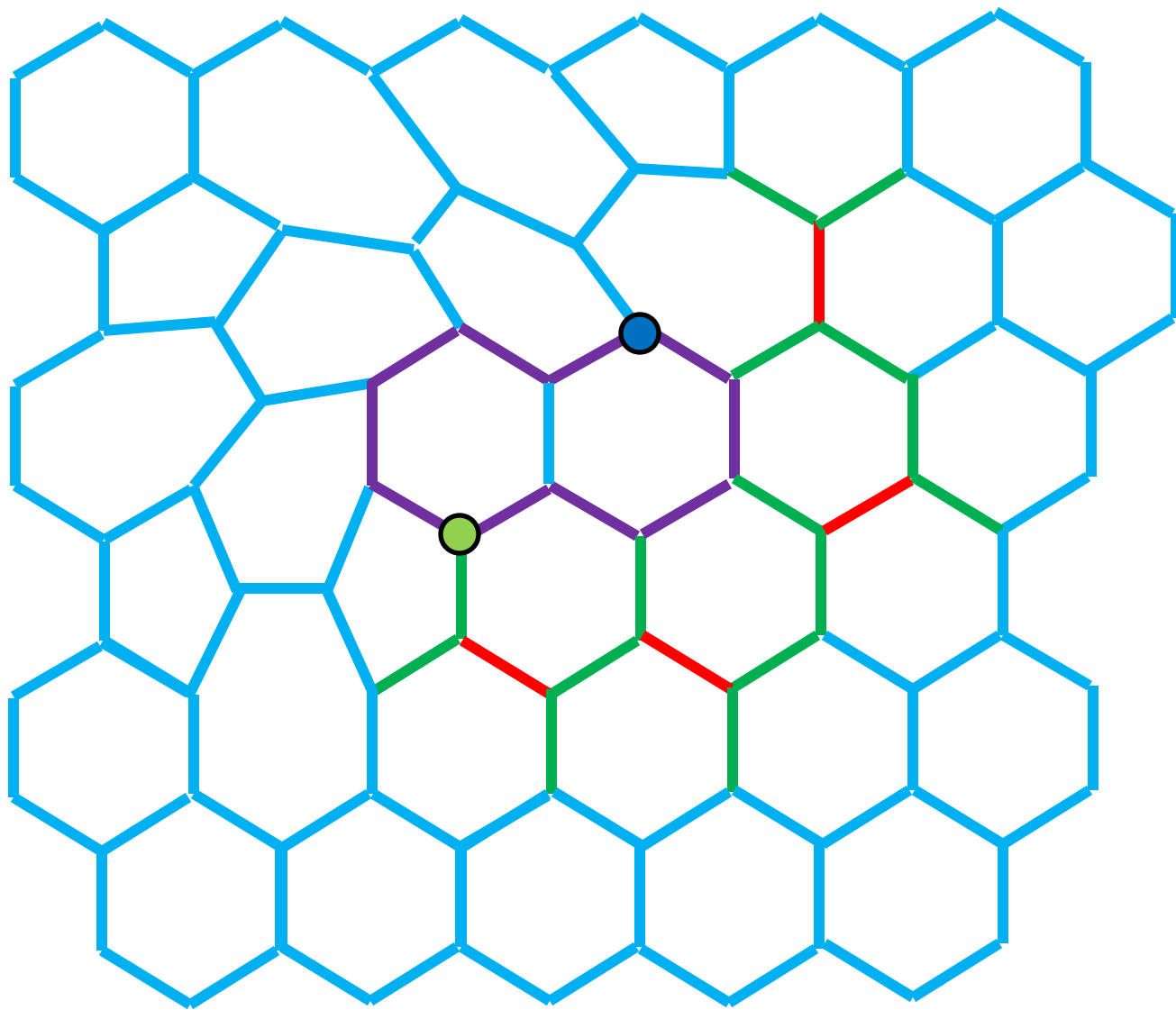
1 F-moves



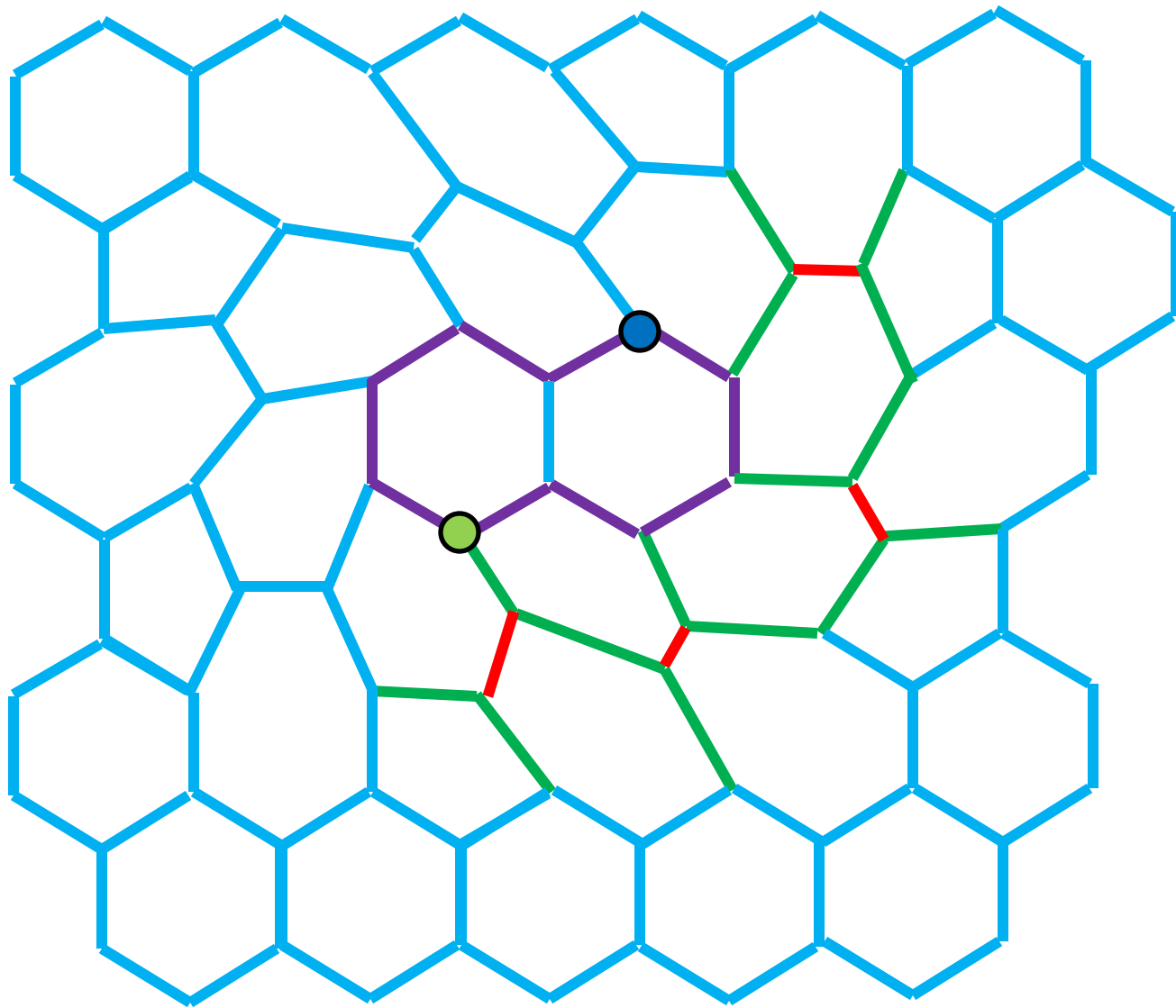
1 F-moves



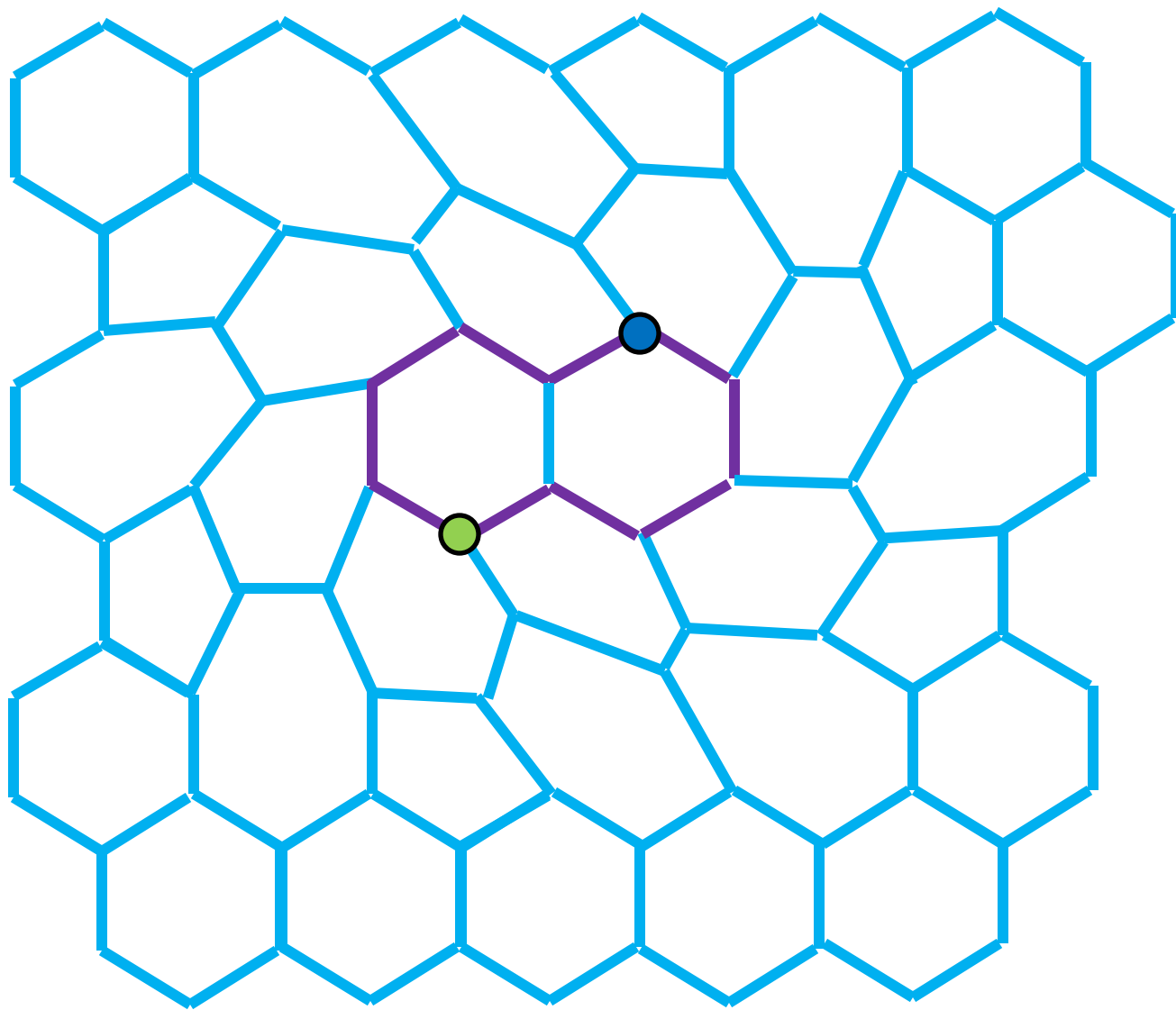
4 F-moves



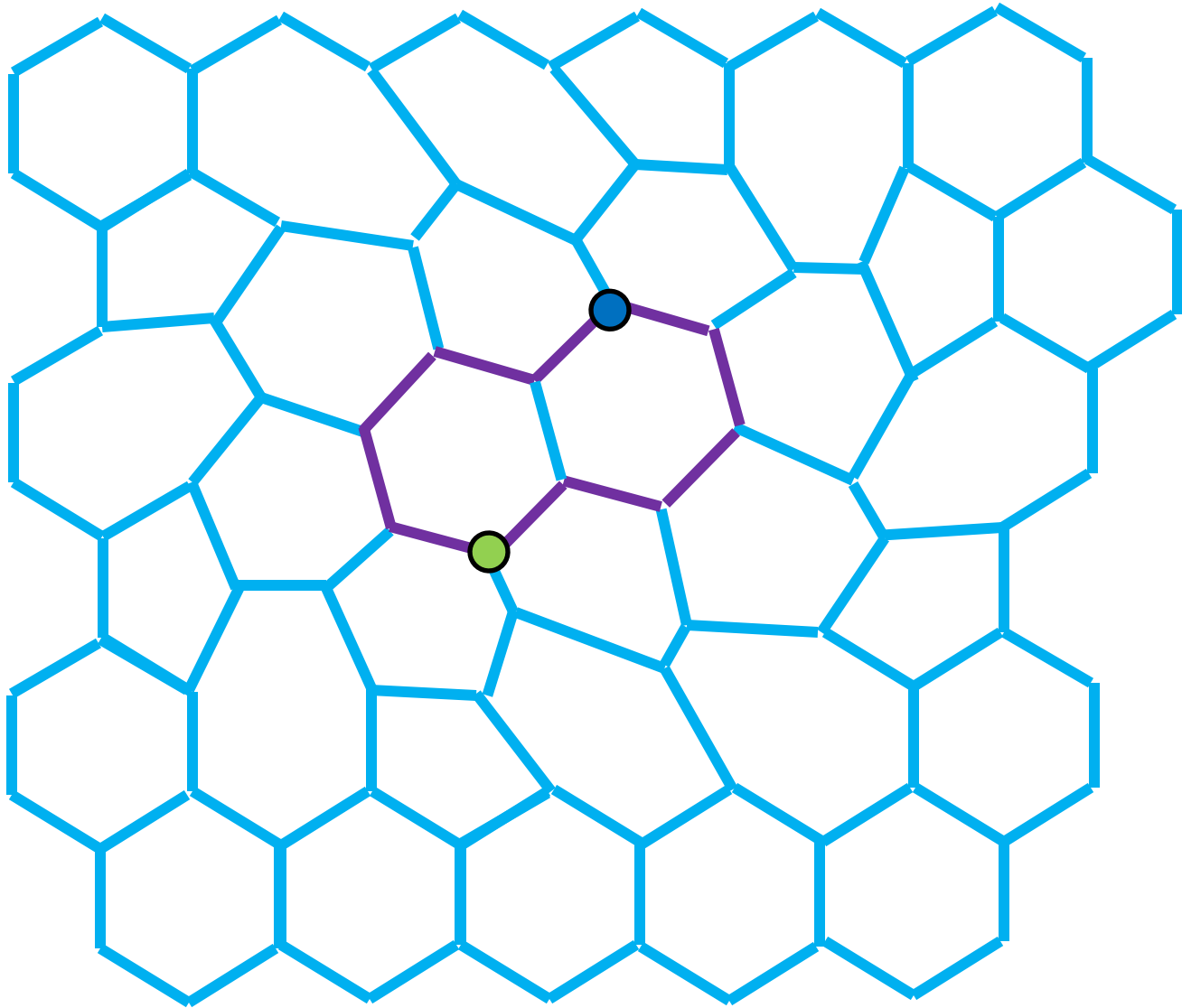
4 F-moves



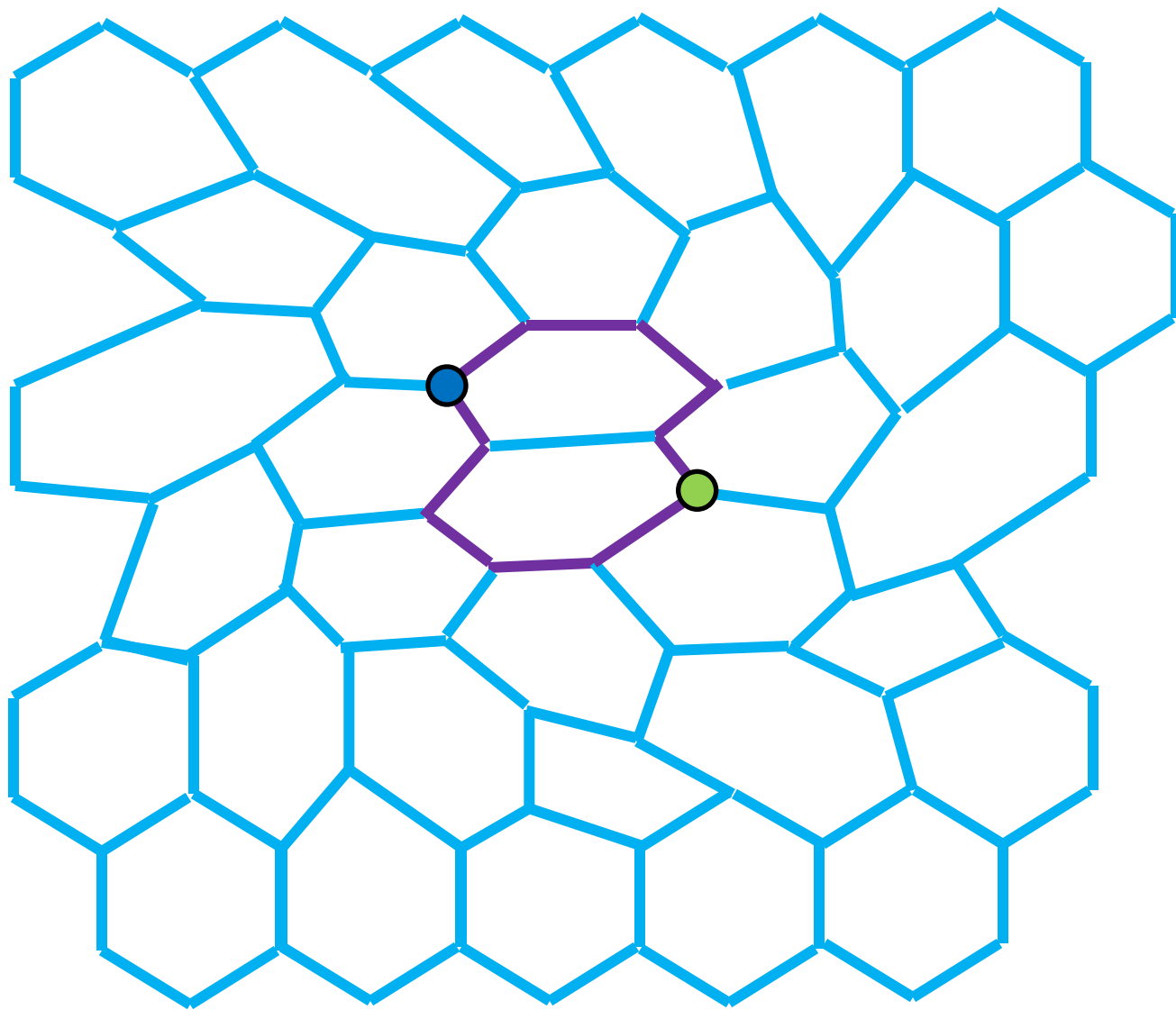
8 F-moves



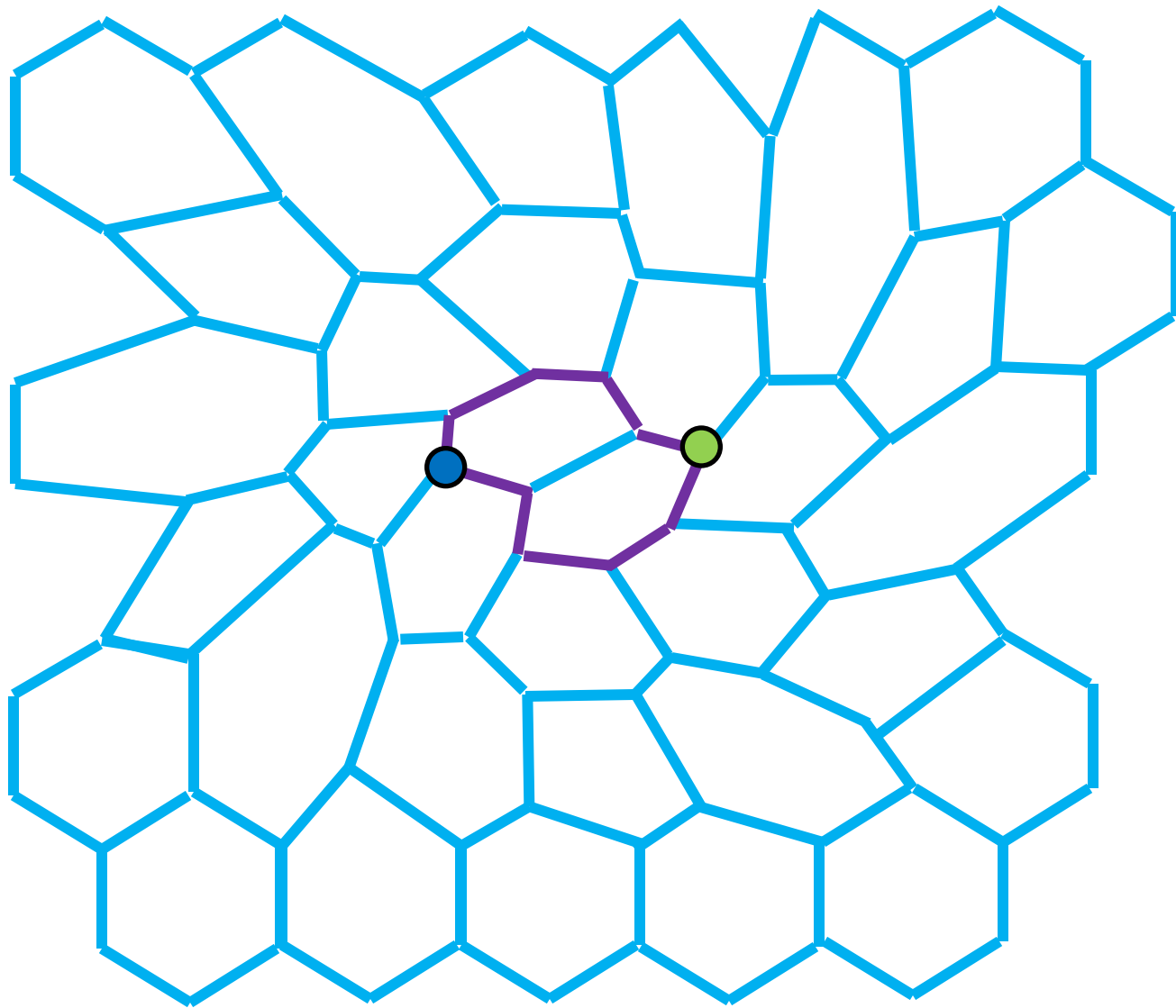
8 F-moves



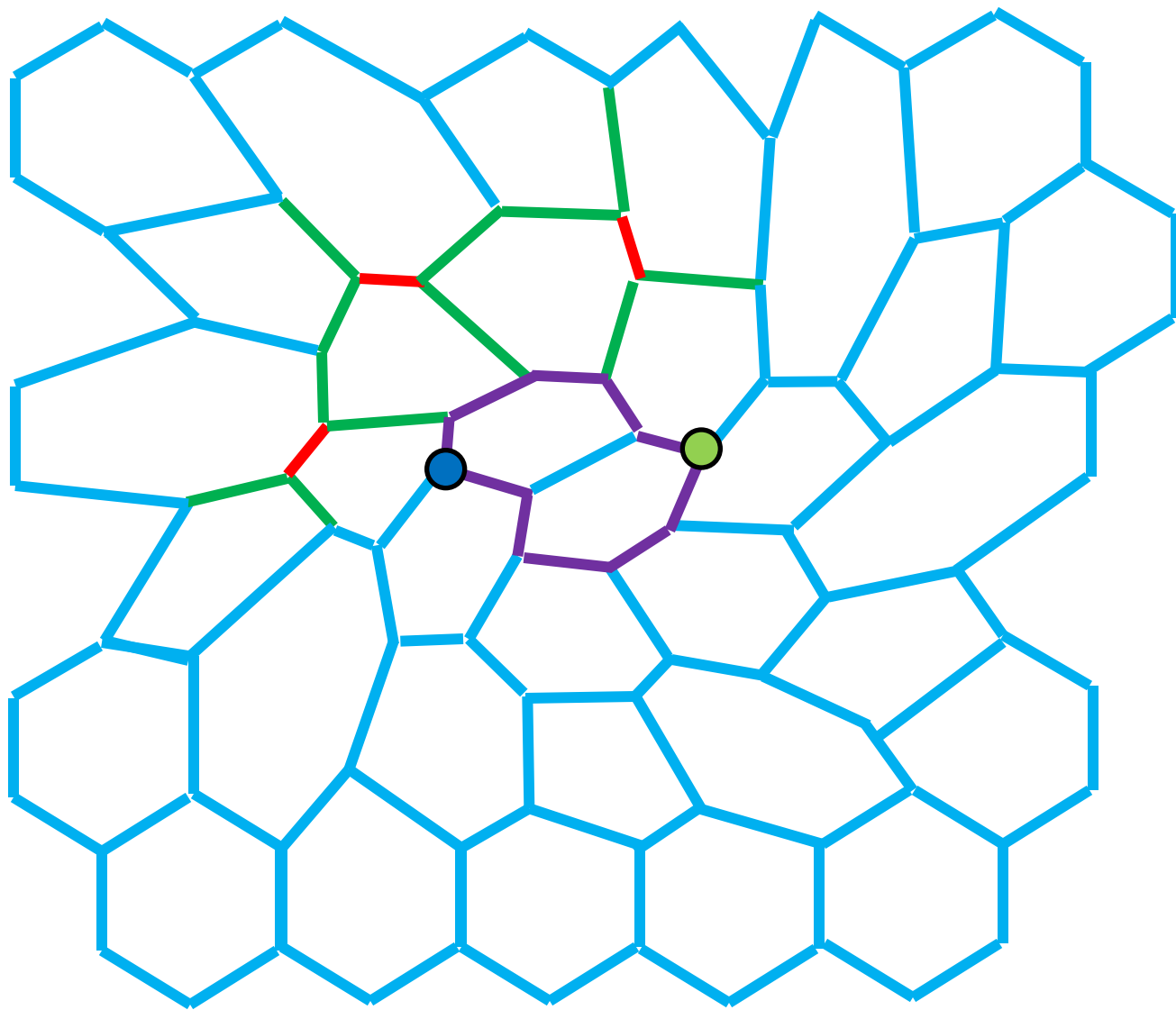
8 F-moves



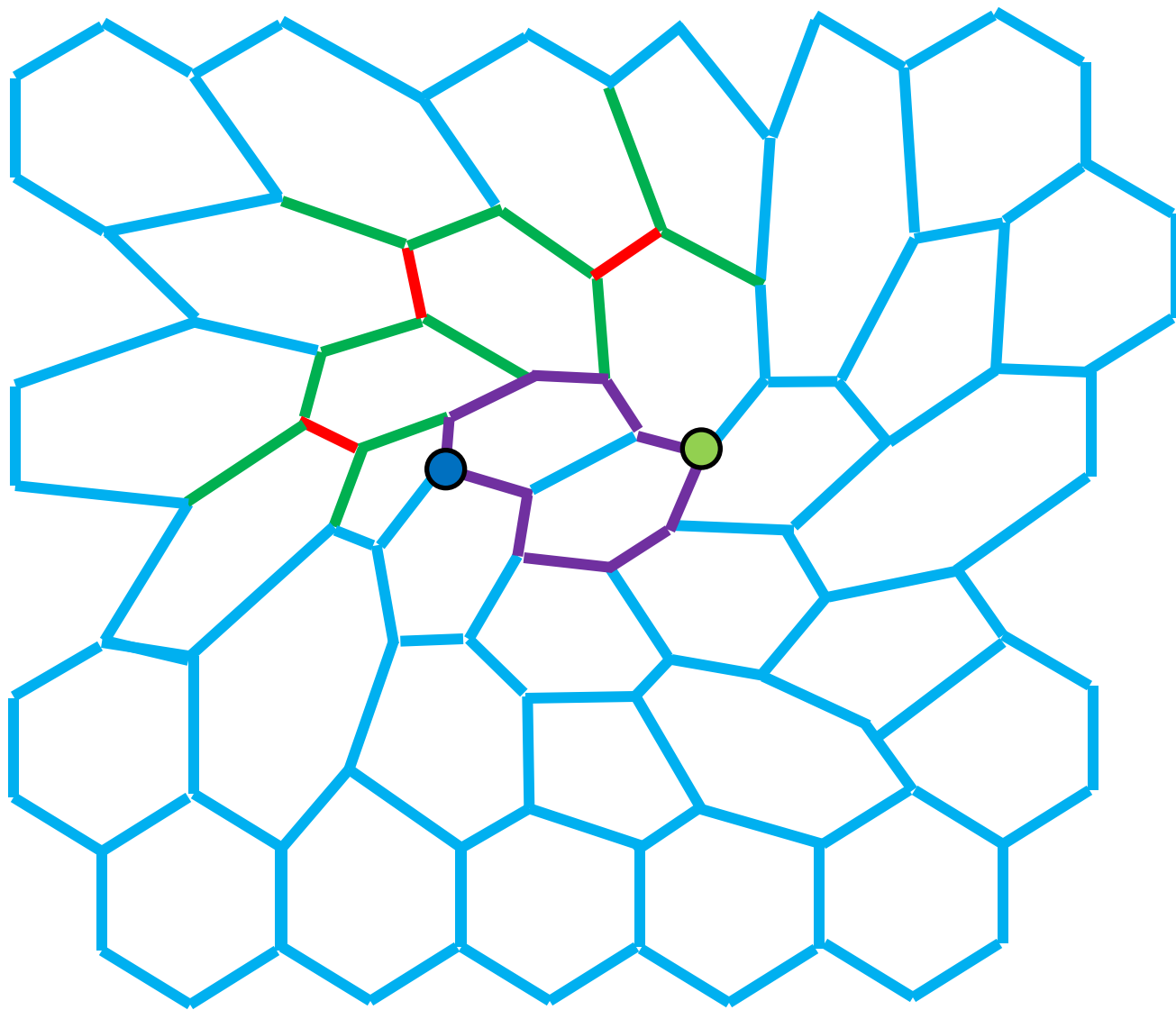
32 F-moves



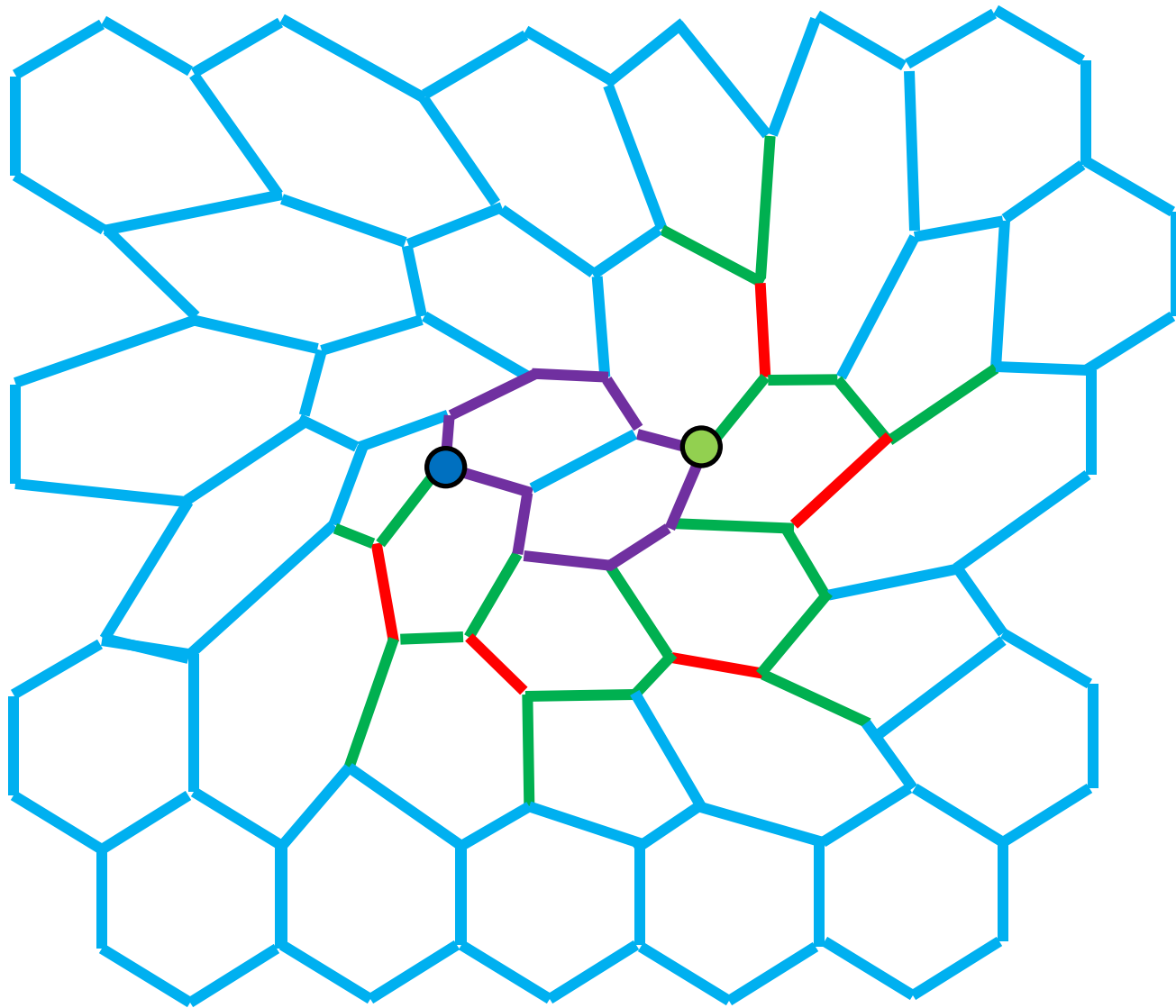
48 F-moves



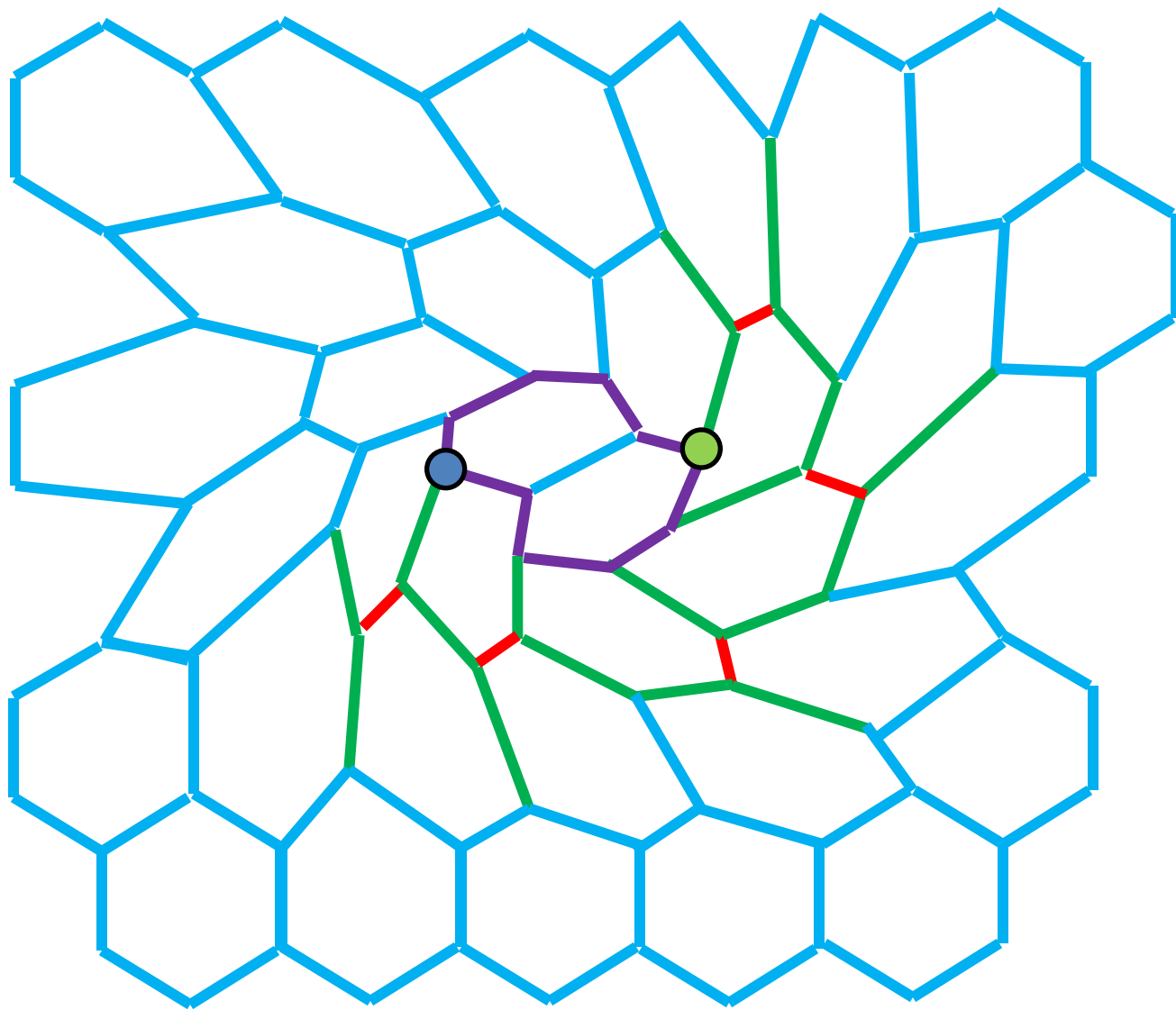
48 F-moves



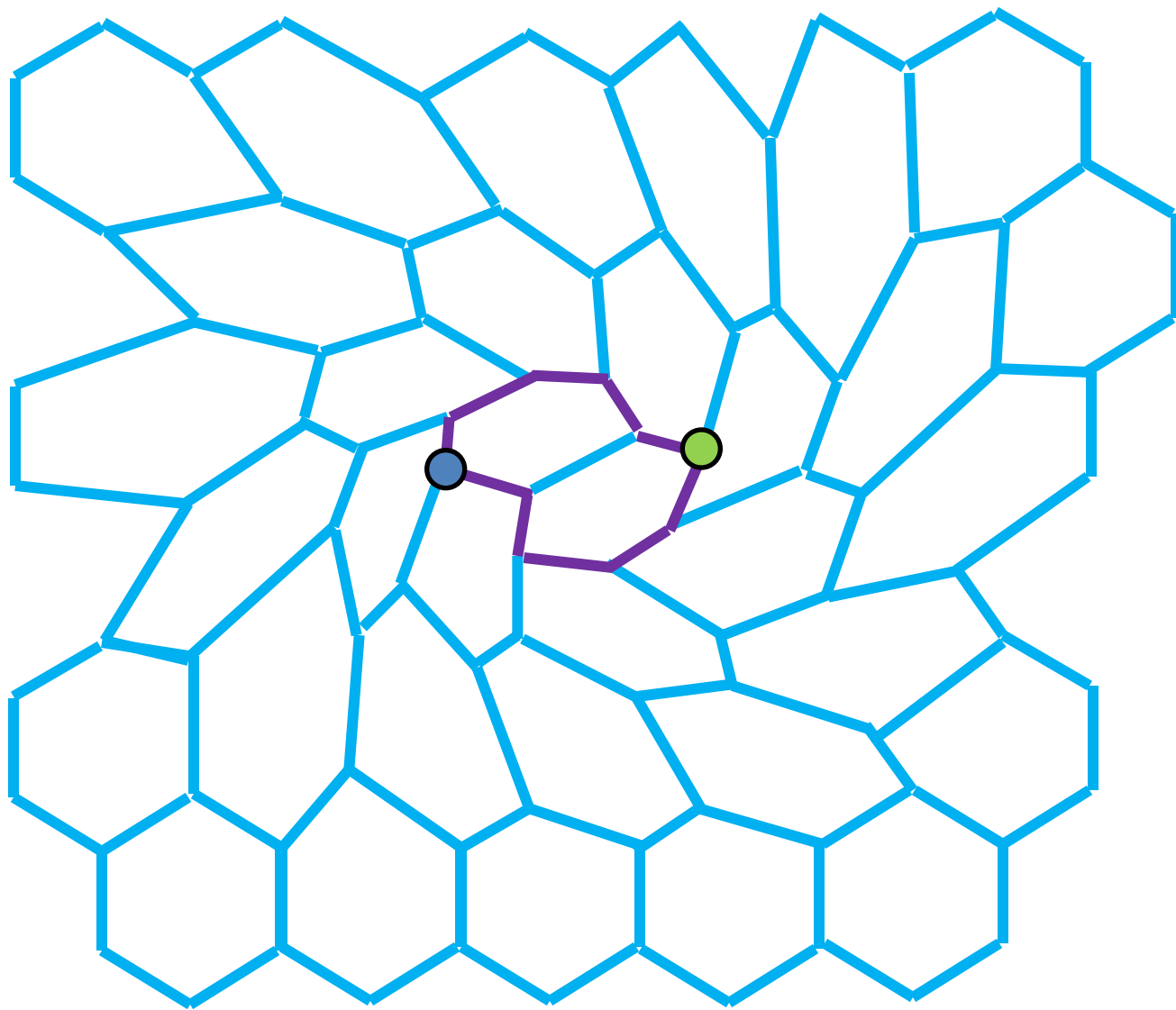
51 F-moves



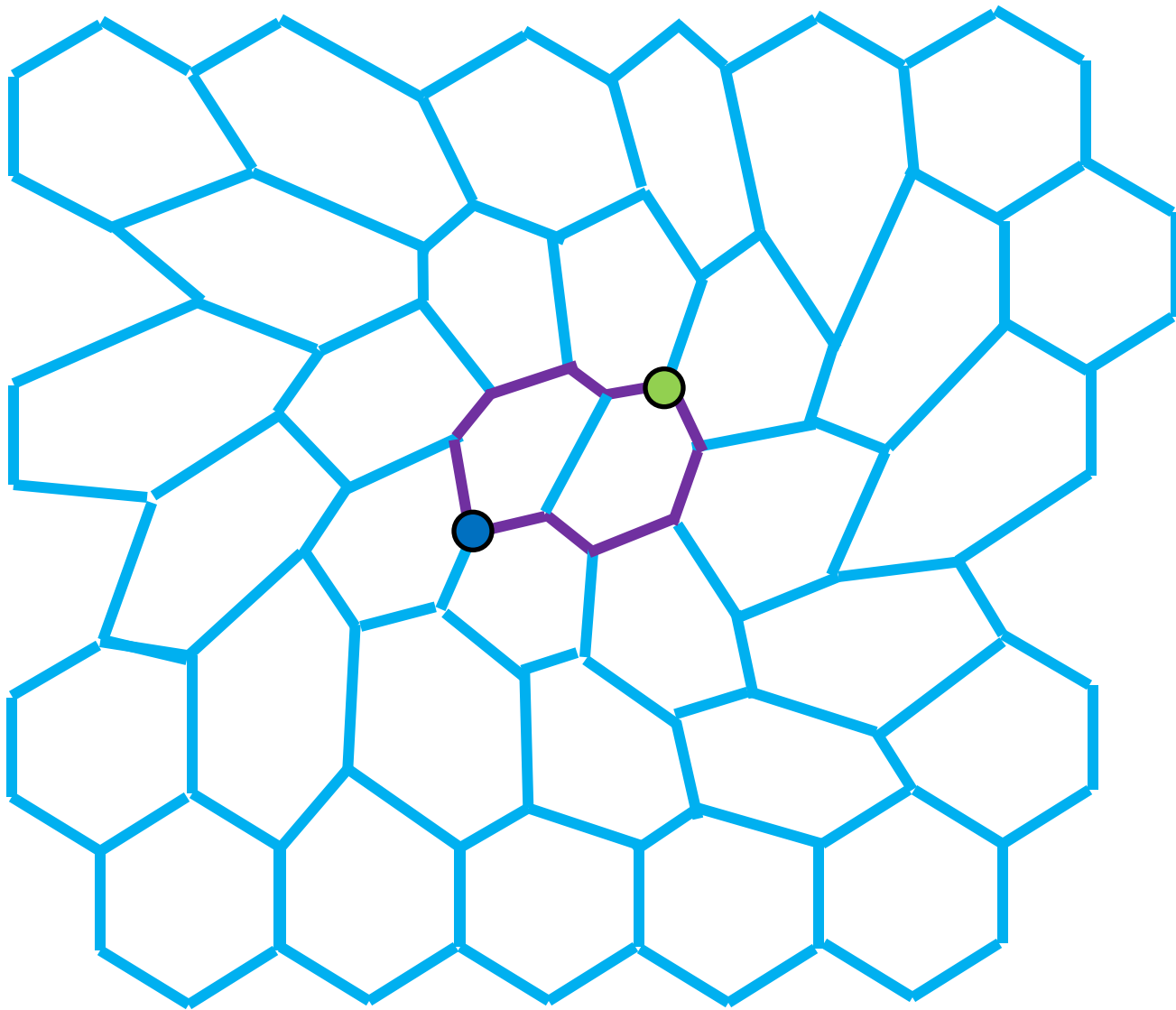
51 F-moves



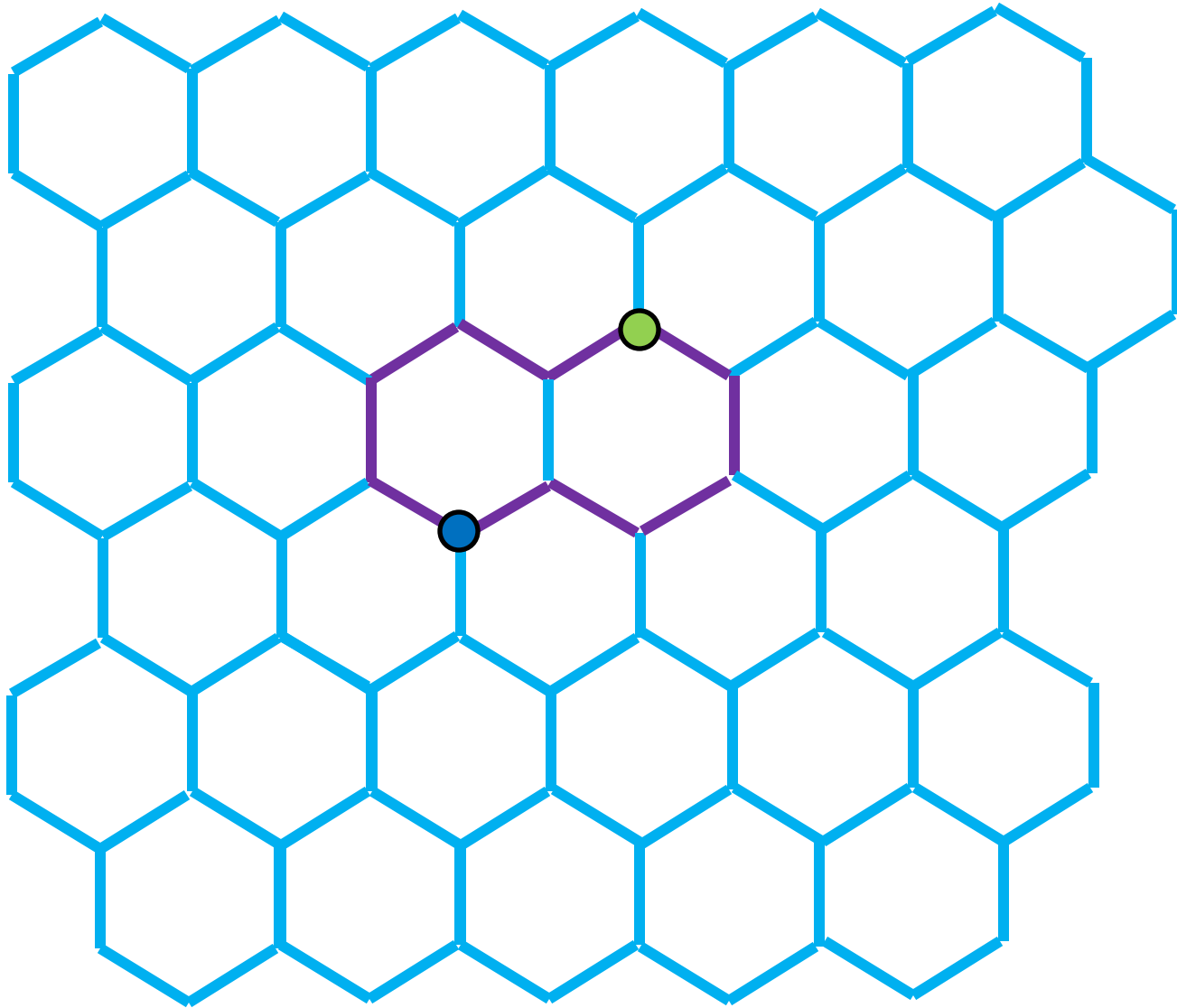
56 F-moves



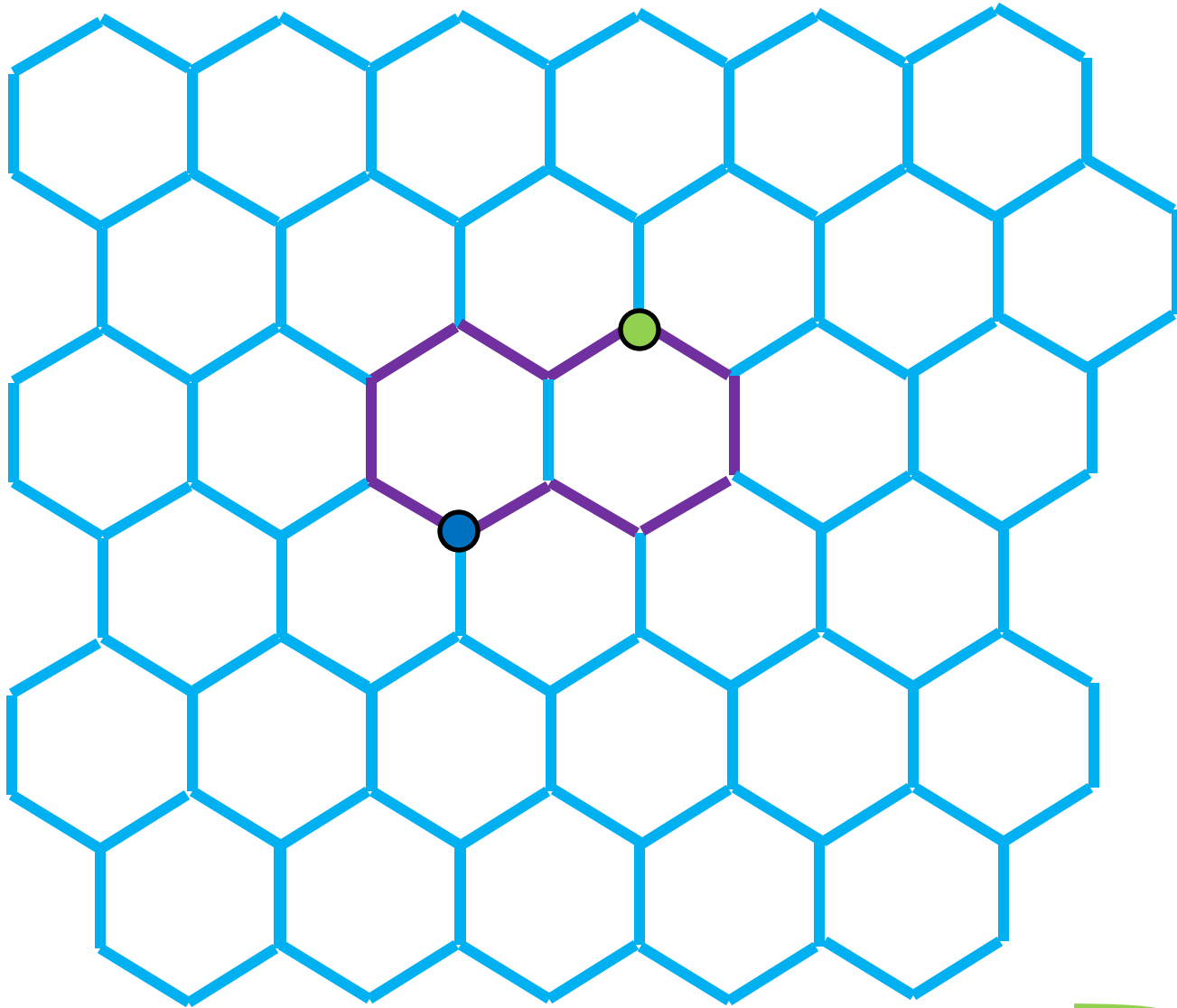
56 F-moves



56 F-moves



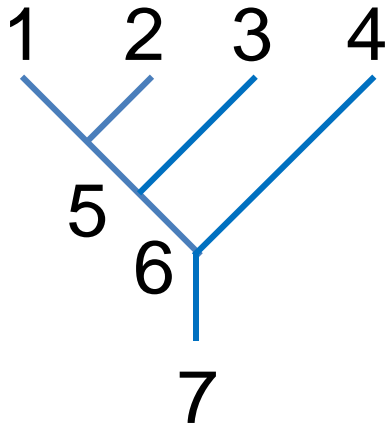
56 F-moves

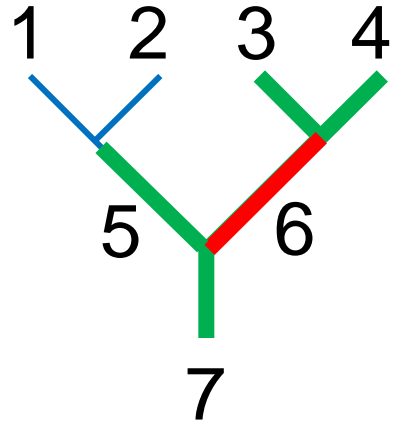
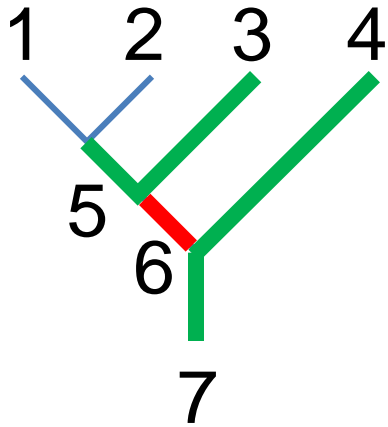


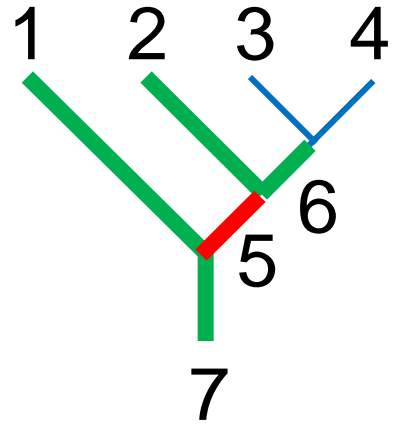
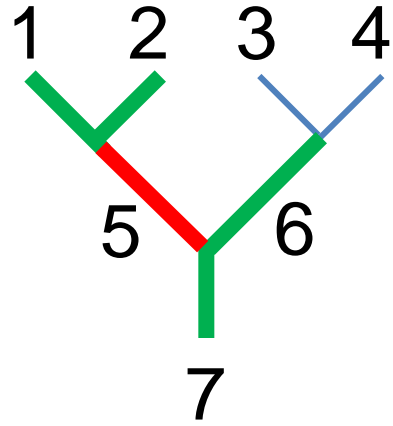
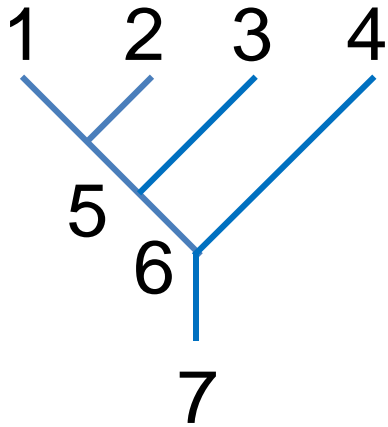
56 F-moves

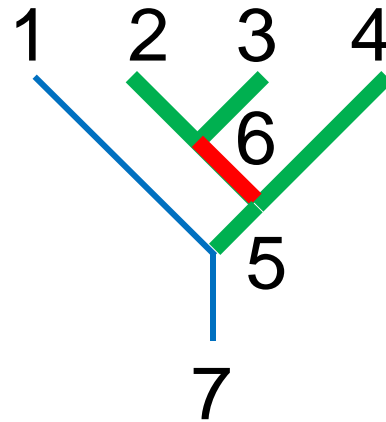
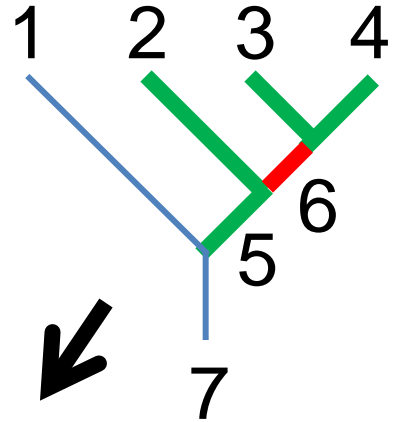
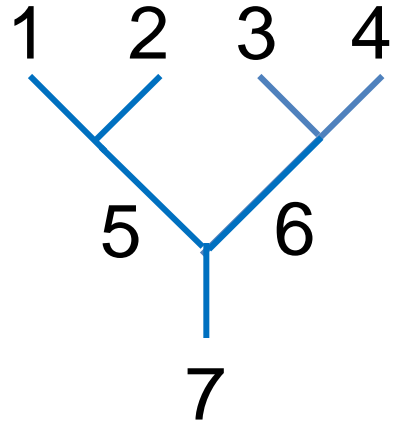
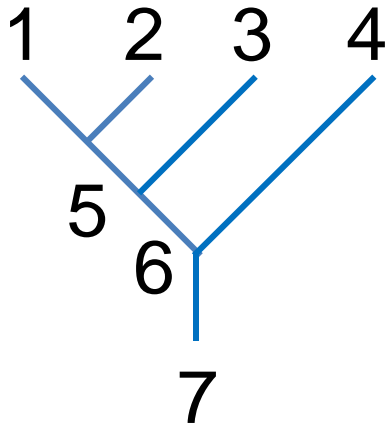
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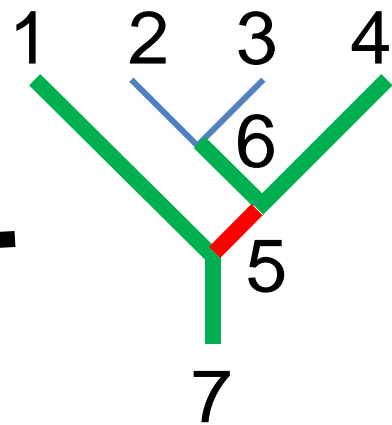
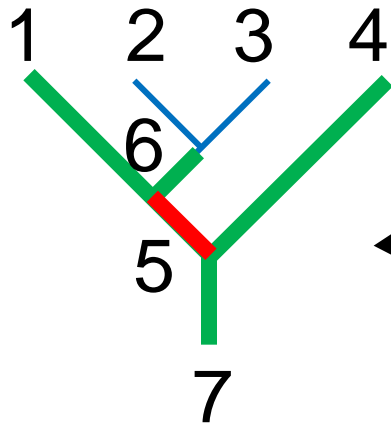
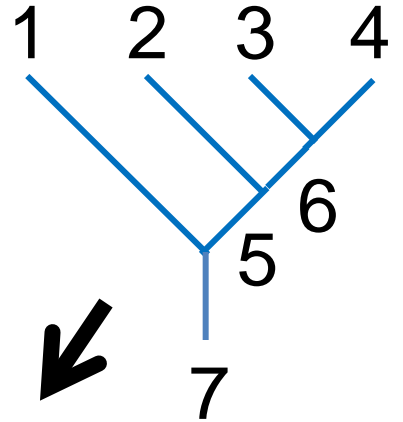
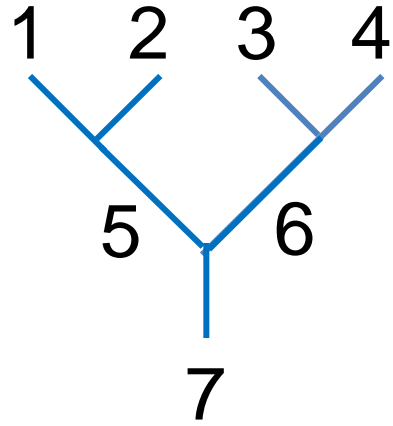
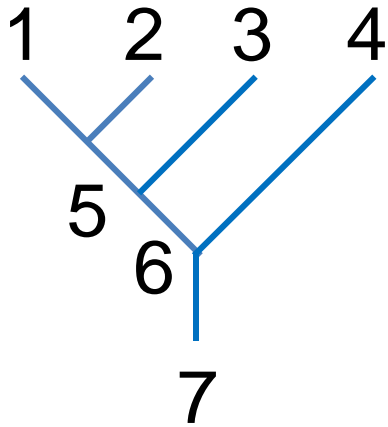


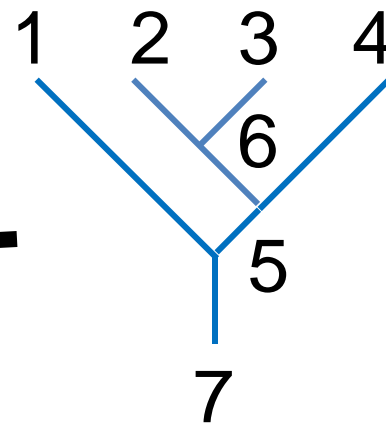
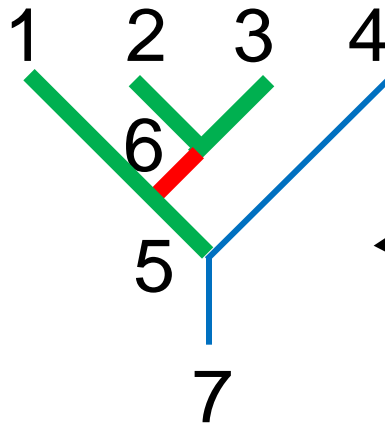
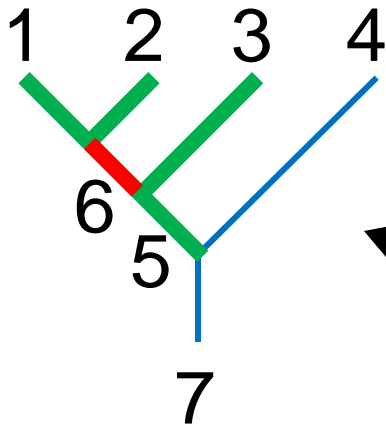
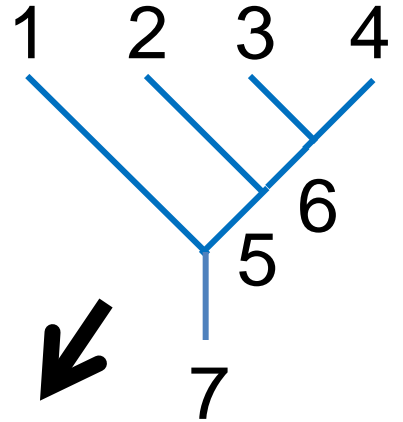
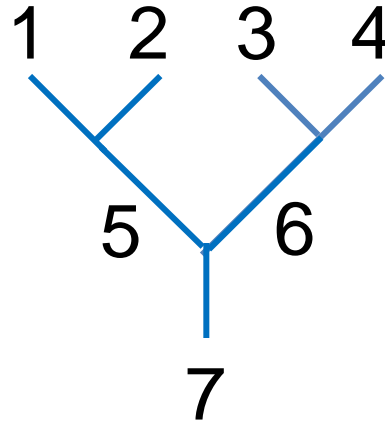
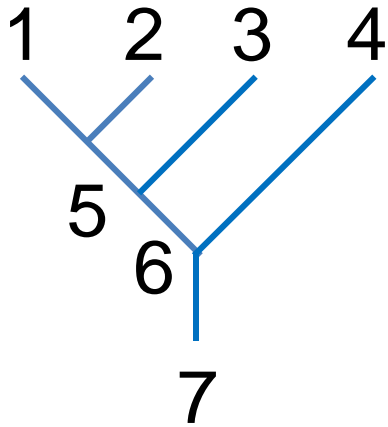


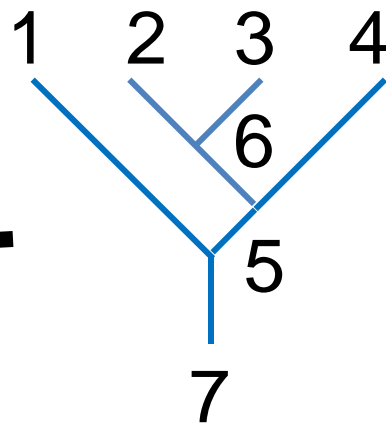
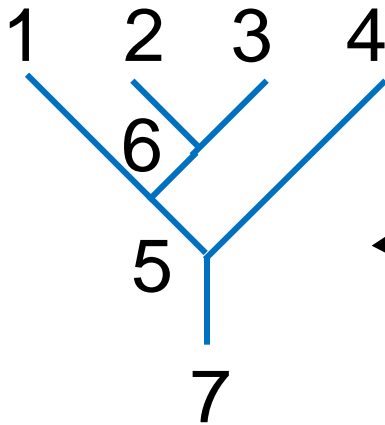
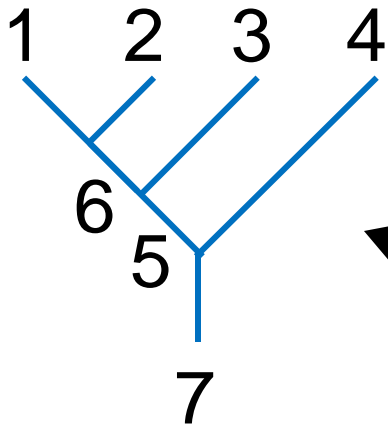
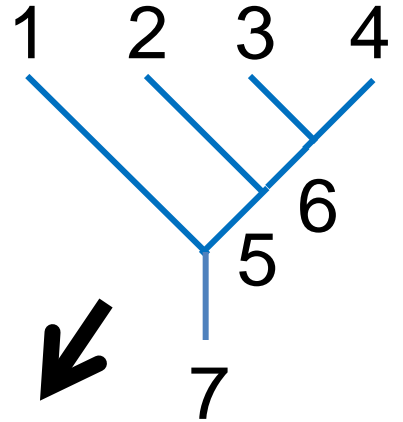
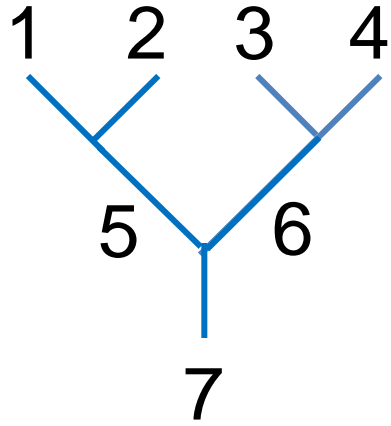
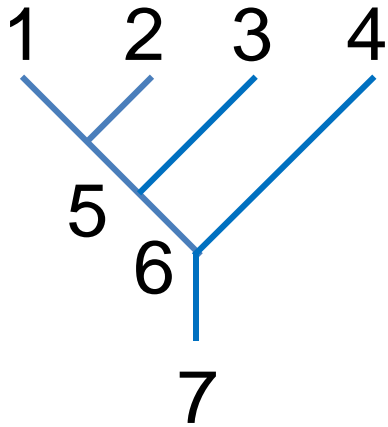


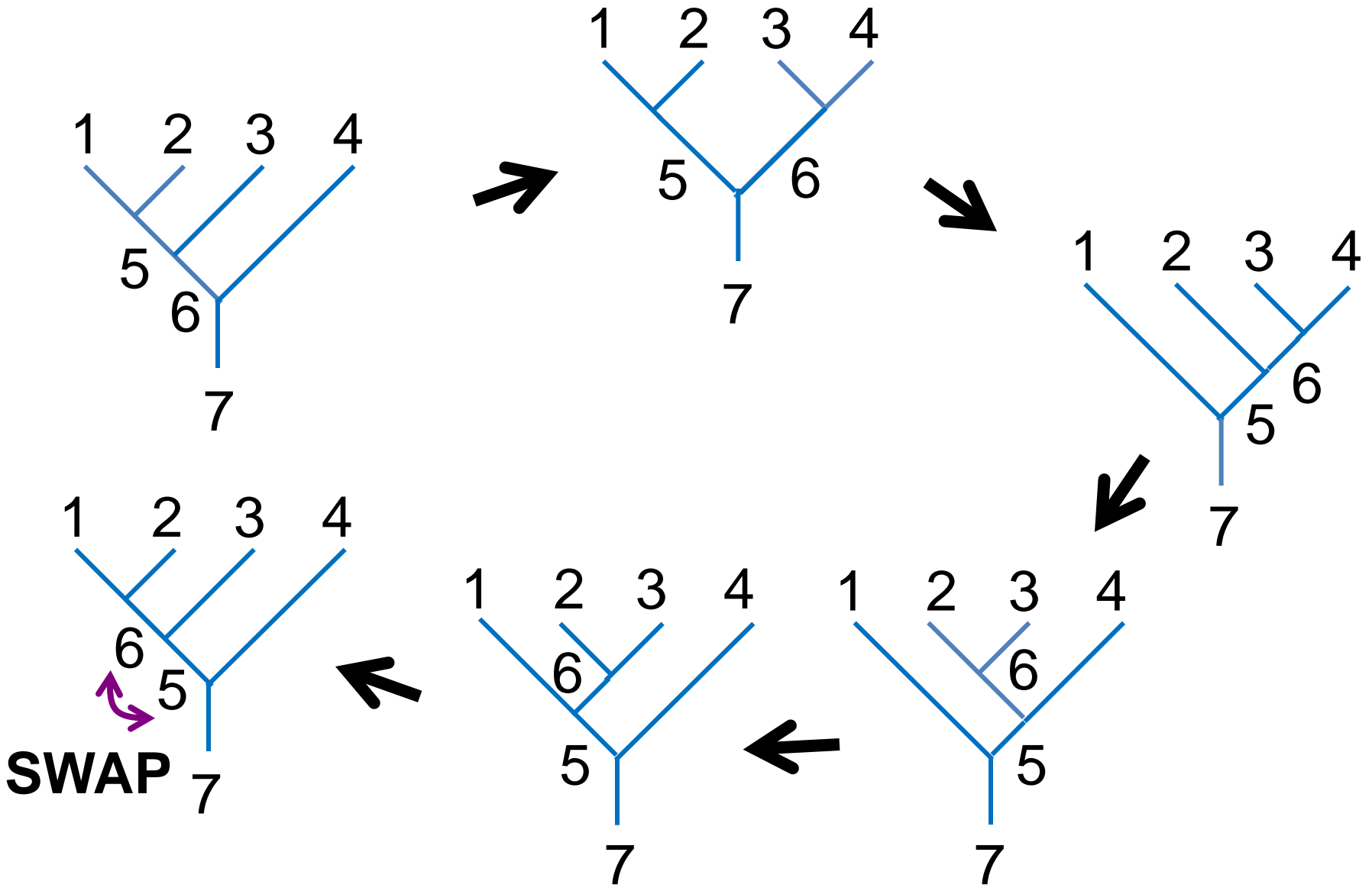




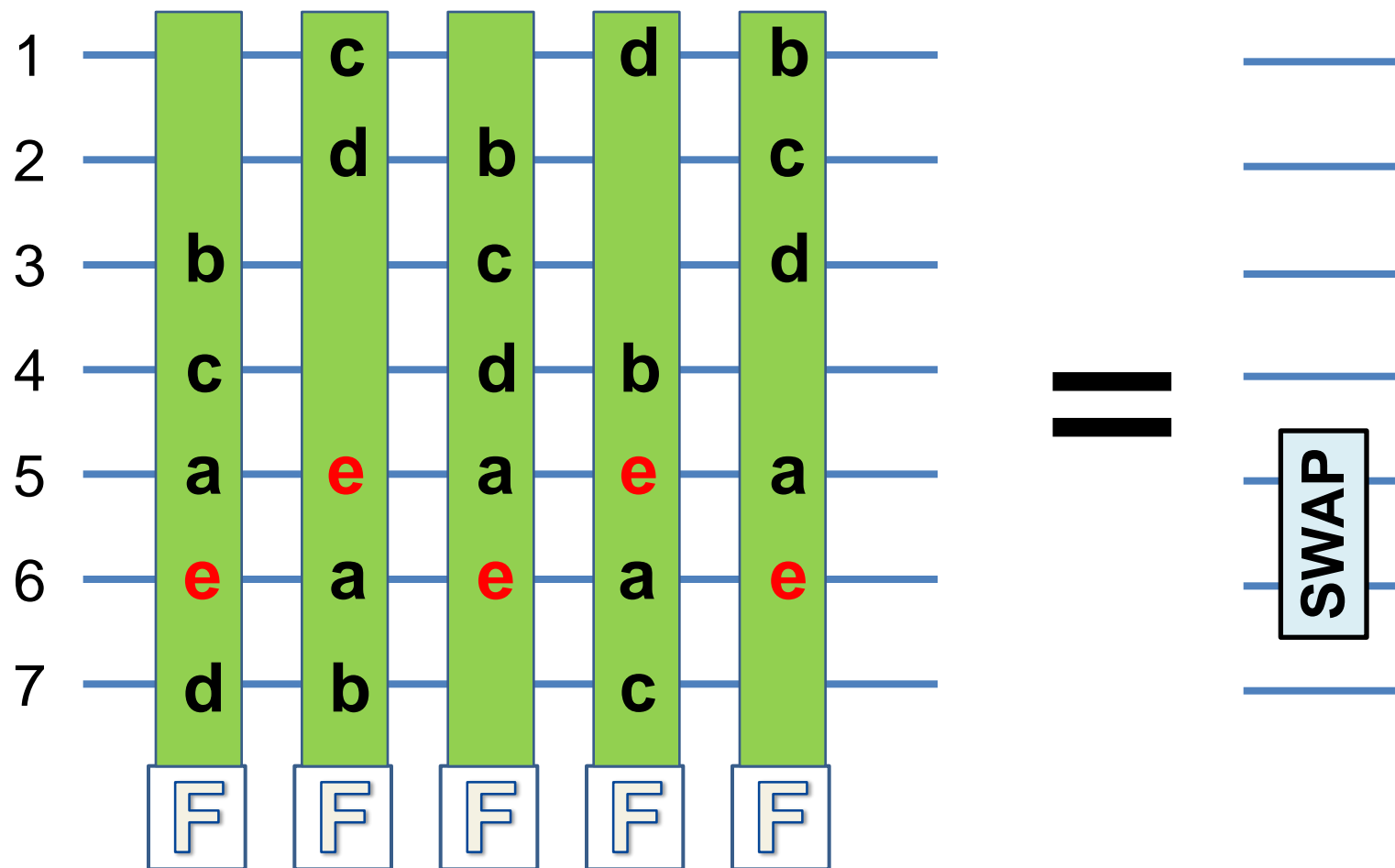




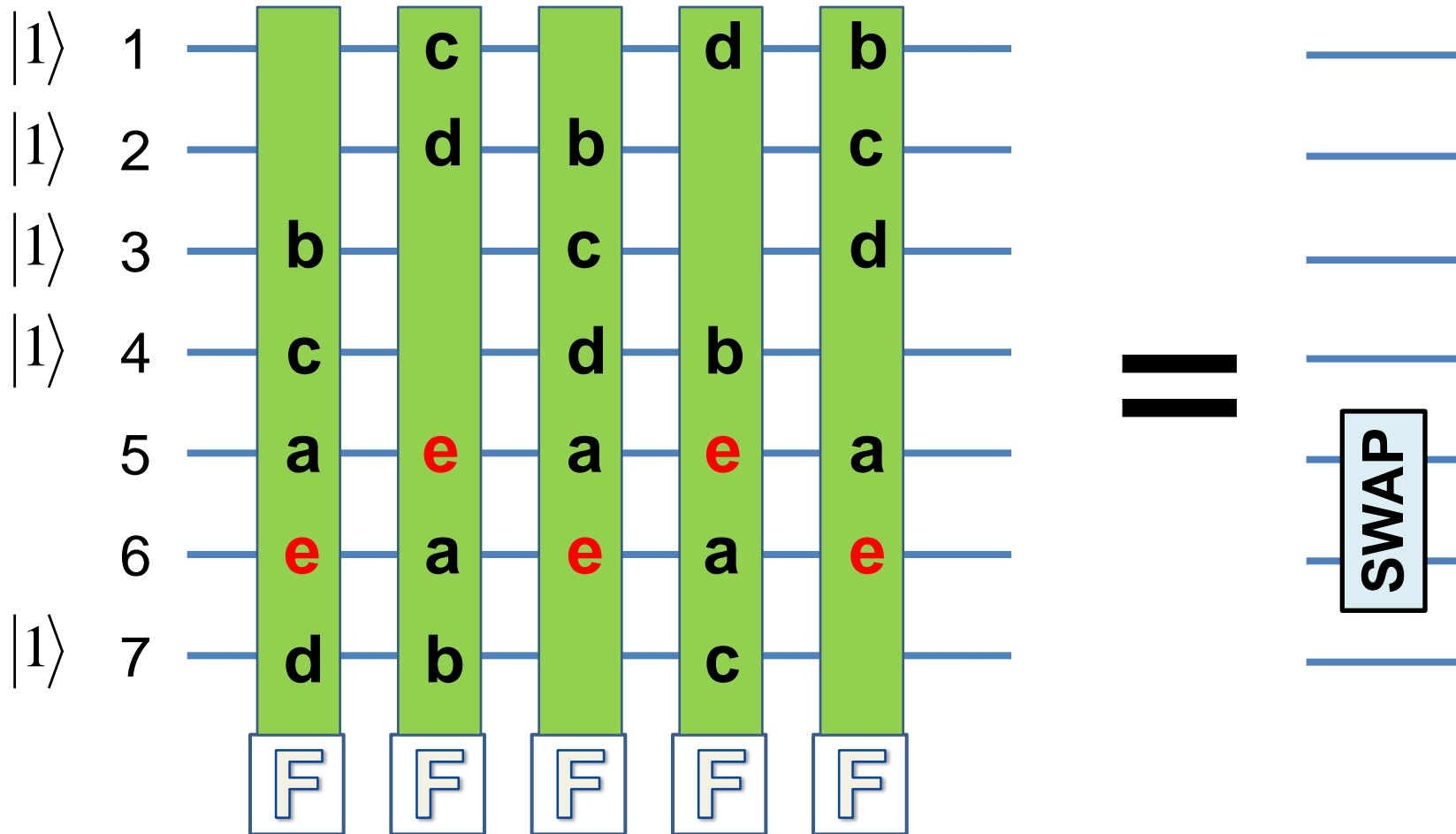




Pentagon Equation as a Quantum Circuit

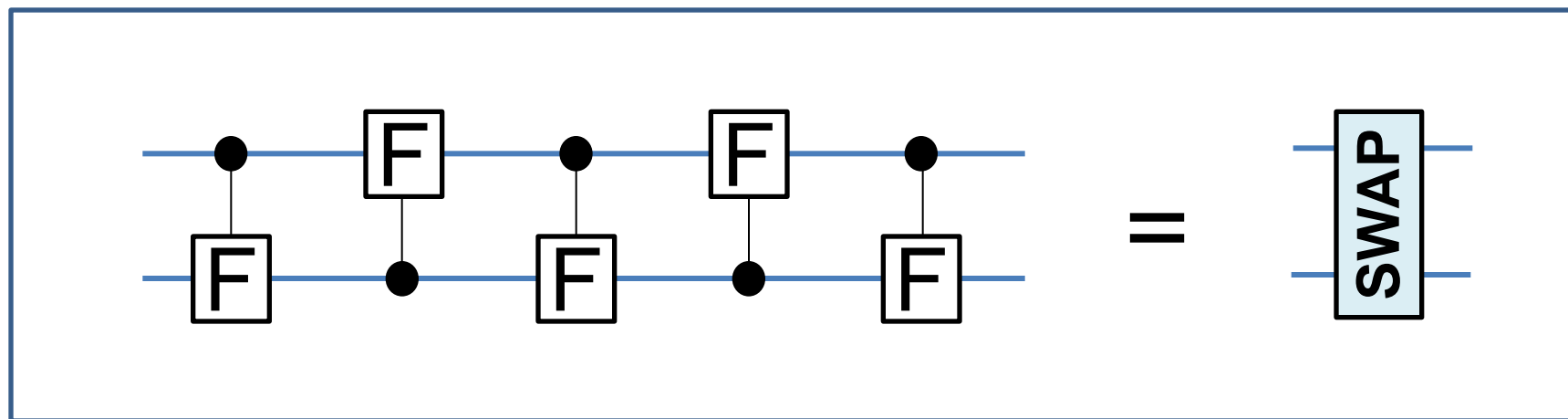


Pentagon Equation as a Quantum Circuit



Pentagon Equation as a Quantum Circuit

$$F = \begin{pmatrix} \varphi^{-1} & \varphi^{-1/2} \\ \varphi^{-1/2} & -\varphi^{-1} \end{pmatrix} \quad \varphi = \frac{\sqrt{5}+1}{2}$$



Simplified Pentagon Circuit

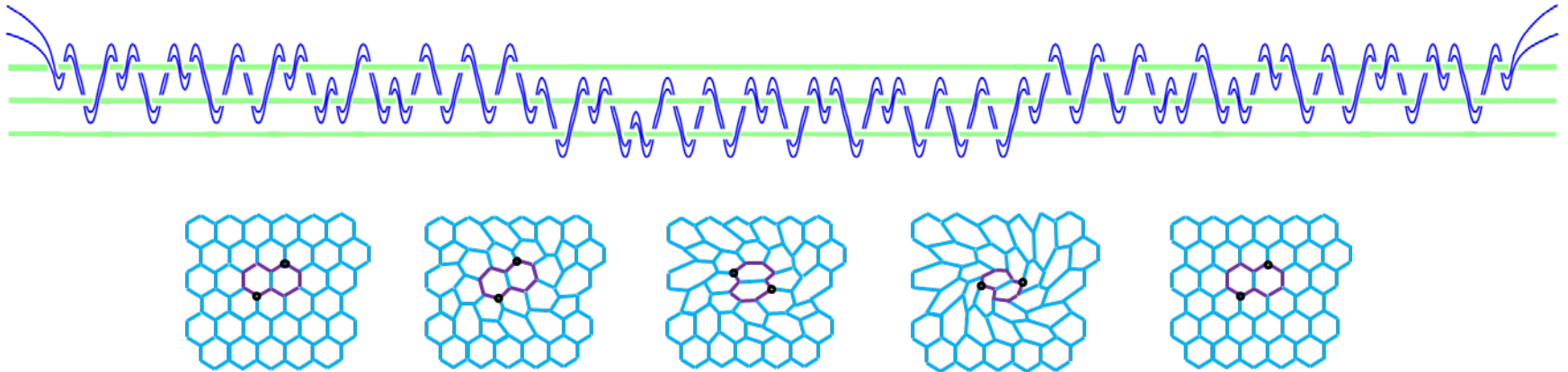
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